

# Discrete Event Simulation

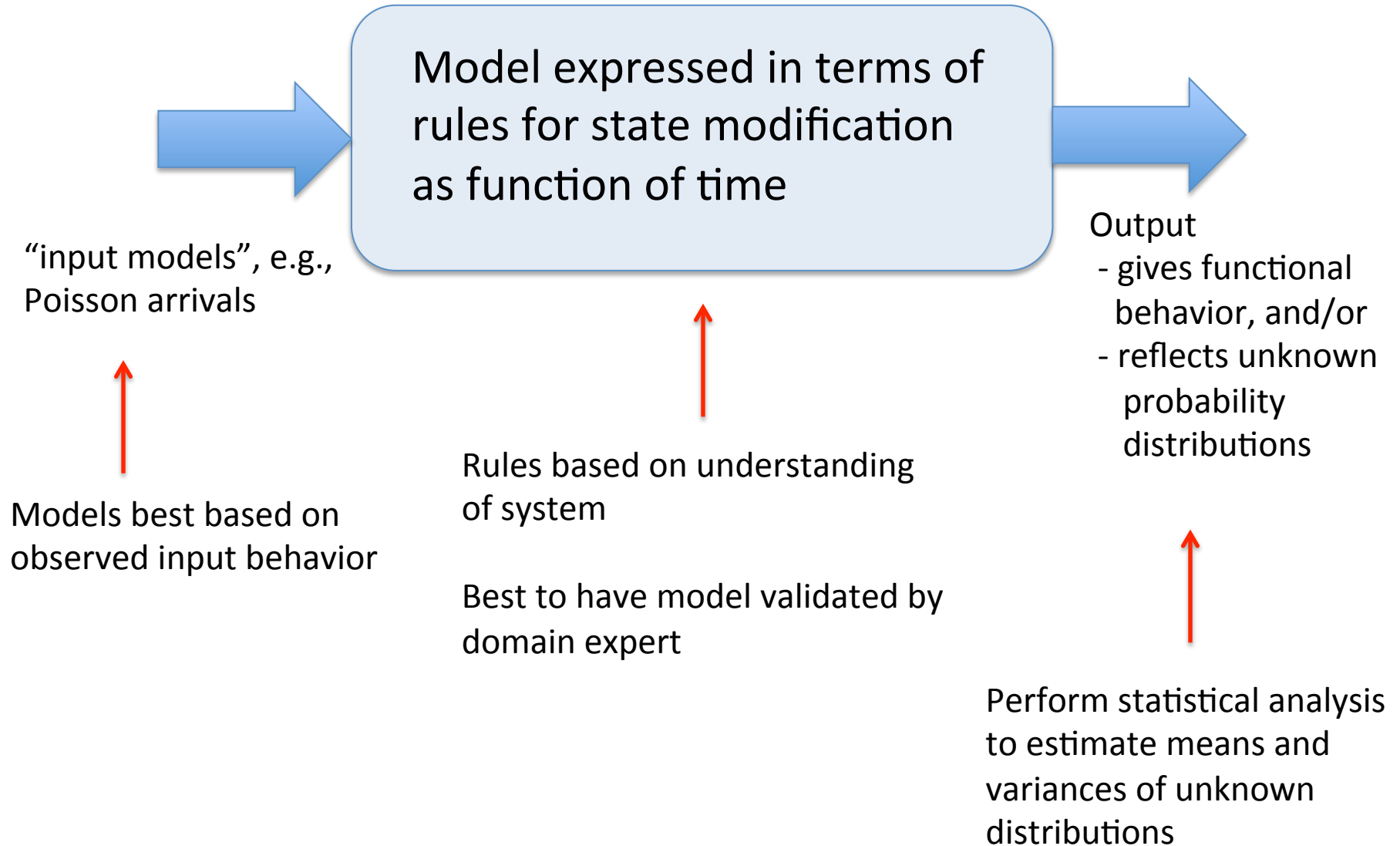
# Motivation

- Markovian models can be analyzed and have their state-spaces completely explored when simple/small enough
- When state spaces become too large, discrete event simulation is often a viable alternative.
- Discrete-event simulation can be used to solve models with arbitrarily large state spaces, as long as the desired measure is not based on a “rare event.”
- When “rare events” are present, variance reduction techniques can sometimes be used.

# Simulation as Model Experimentation

- State-based methods enumerate all possible states a system can be in, and then apply a numerical solution method on the generated state space.
- Simulation generates one or more trajectories (possible behaviors from the high-level model), and collects statistics from these trajectories to estimate the desired performance/dependability measures.
- Just how this trajectory is generated depends on the:
  - nature of the notion of state (continuous or discrete)
  - type of stochastic process (e.g., ergodic, reducible)
  - nature of the measure desired (transient or steady-state)
  - types of delay distributions considered (exponential or general)

# View of Simulation Model



# Types of Simulation

*Continuous-state simulation* is applicable to systems where the notion of state is continuous and typically involves solving (numerically) systems of differential equations. Circuit-level simulators are an example of continuous-state simulation.

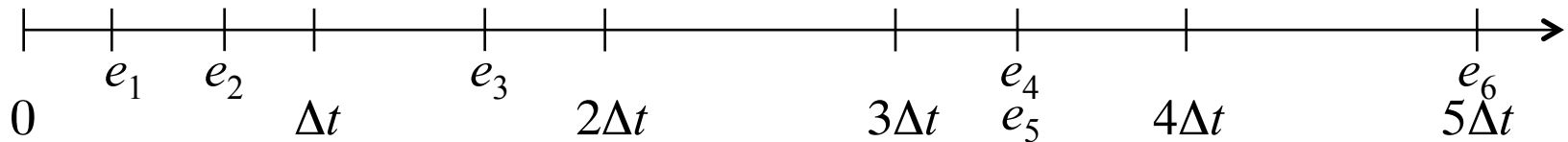
*Discrete-event simulation* is applicable to systems in which the state of the system changes at discrete instants of time, with a finite number of changes occurring in any finite interval of time.

There are two types of discrete-event simulation execution algorithms:

- Fixed-time-stamp advance
- Variable-time-stamp advance

# Fixed-Time-Stamp Advance Simulation

- Simulation clock is incremented a fixed time  $\Delta t$  at each step of the simulation.
- After each time increment, each event type is checked to see if it should have completed during the time of the last increment.
- All event types that should have completed are completed and a new state of the model is generated.
- Rules must be given to determine the ordering of events that occur in each interval of time.
- Example:

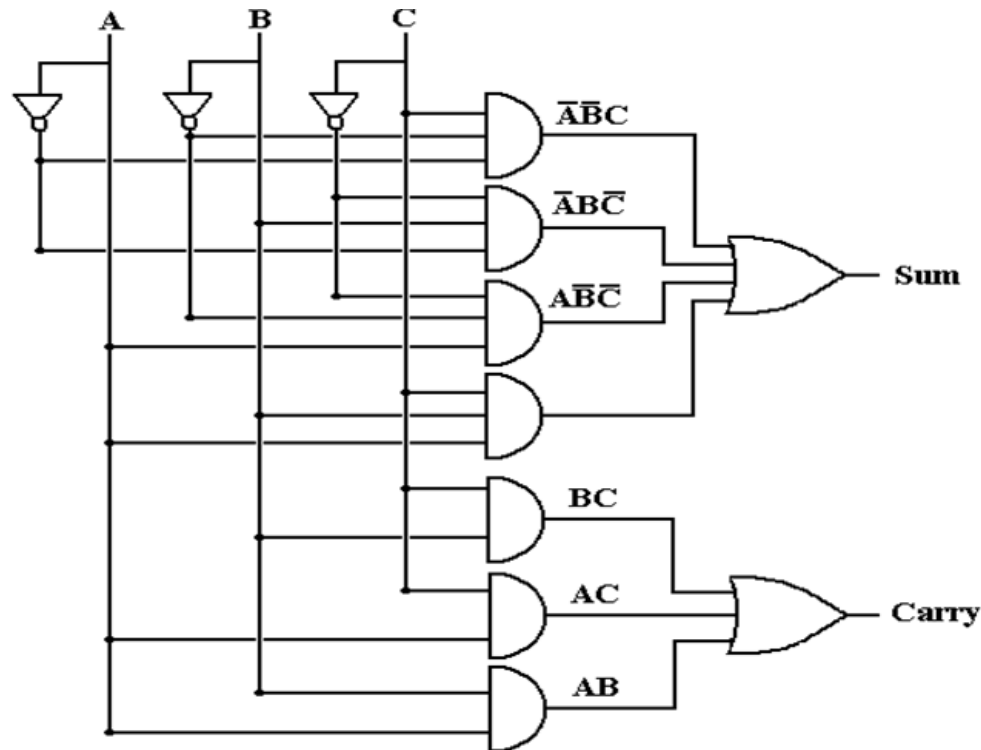


- Good for all models where most events happen at fixed increments of time (e.g., gate-level simulations).
- Has the advantage that no “future event list” needs to be maintained.
- Can be inefficient if events occur in a bursty manner, relative to time-step used.

# Example of Fixed Time Step Simulation

## Gate level simulator

- State updates everywhere each  $\Delta$  unit of time
  - Simple logic simulator, each  $\Delta$  a gate delay time
    - Each  $\Delta$  set gate outputs as function of inputs

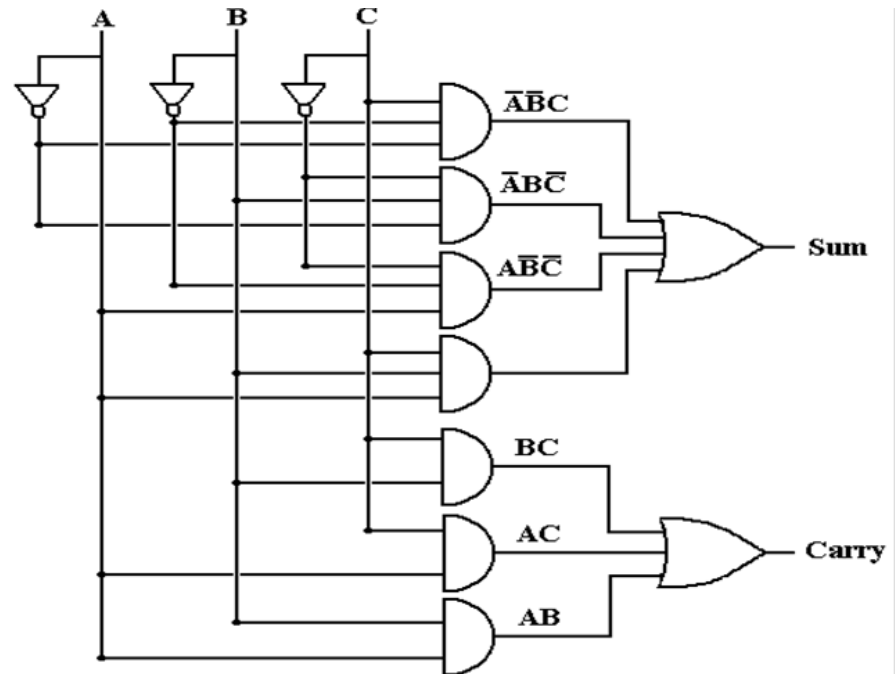


# Example of Fixed Time Step Simulation

Focus state changes on instants in time and location in model where state change happens

- Means we *schedule* localized state updates
- This is an “event”
  - Specifies action, location in model, and time

Event = (set value, gate input, t)





# Example of Fixed Time Step Simulation

Focus state changes on instants in time and location in model where state change happens

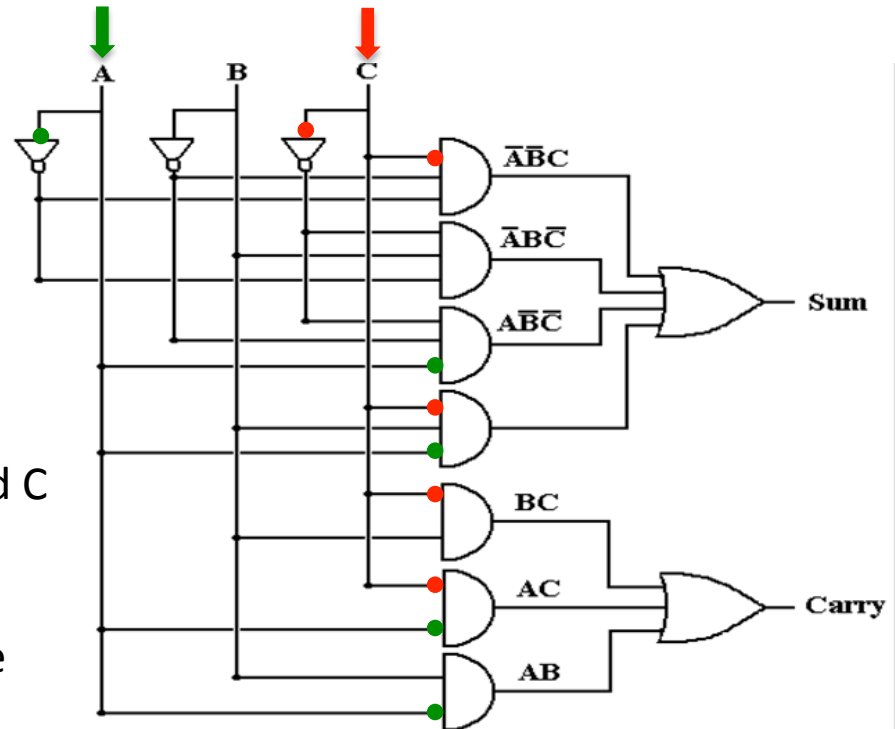
- Means we *schedule* localized state updates
- This is an “event”
  - Specifies action, location in model, and time

Event = (set value, gate input, t)

T = 1000 gate delays

Observe input changes  
as result of changes on A and C

Schedule only one gate eval  
when multiple inputs change



# Example of Fixed Time Step Simulation

Focus state changes on instants in time and location in model where state change happens

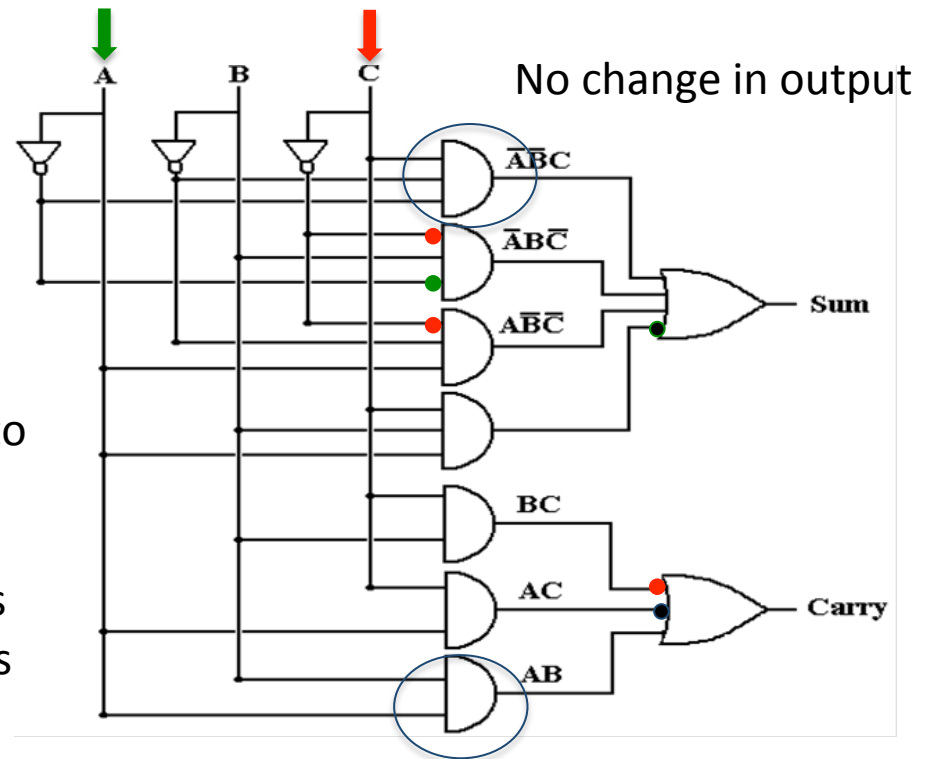
- Means we *schedule* localized state updates
- This is an “event”
  - Specifies action, location in model, and time

Event = (set value, gate input, t)

T = 1001 gate delays

Observe input changes due to Inverter output changes

Observe some input changes  
Do not cause output changes



# Example of Fixed Time Step Simulation

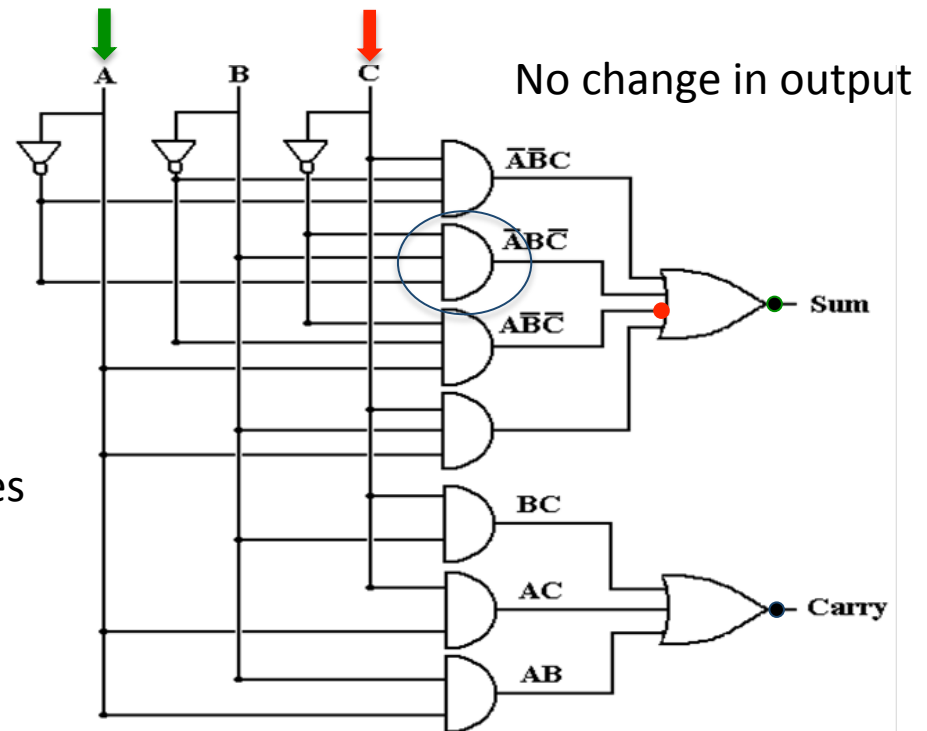
Focus state changes on instants in time and location in model where state change happens

- Means we *schedule* localized state updates
- This is an “event”
  - Specifies action, location in model, and time

Event = (set value, gate input, t)

T = 1002 gate delays

Observe continued selective propagation of signal changes



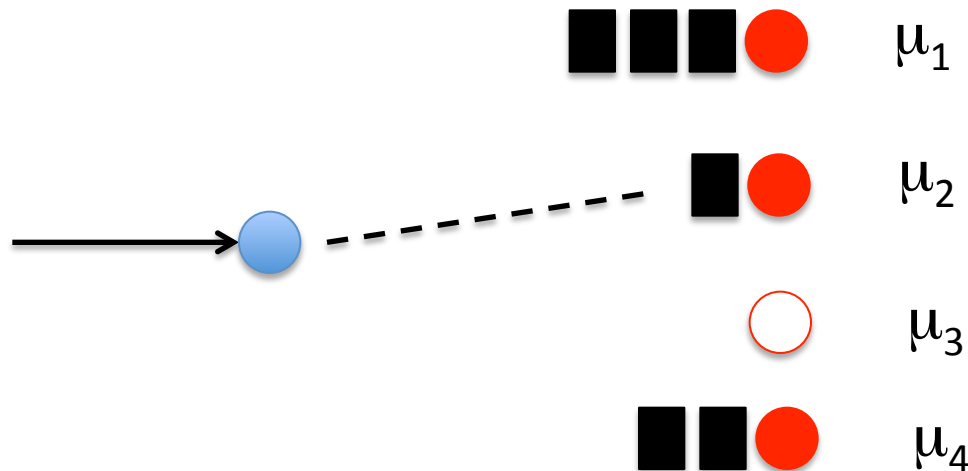
# Variable-Time Step Advance Simulation

- Simulation clock advanced a variable amount of time each step of the simulation, to time of next event.
- If event times are general (have memory) then “future event list” is needed.
- Has the advantage (over fixed-time-stamp increment) that periods of inactivity are skipped over, and models with a bursty occurrence of events are not inefficient.

# Example of Variable Time Step Simulation

## Random inter-arrivals to multi-queue system

- Arrival sent to queue with shortest expected finishing time (number in system\* (mean service))
- Service distributions different queues vary



# Queueing Example #1

1. First schedule an arrival to the system
  - (arrival, system, t)
2. Choose “next event” in simulation time
  - Implicitly, model state remains untouched in time since most recent event

## **Processing of Arrival to System (arrival, system, t)**

- a) Find queue Q with least **expected** finishing time if job->Q  
(assuming FCFS queue management, service not yet known)  
If Q not empty then add job to tail of Q  
else  
    sample service time  $s_0$   
    schedule departure event (departure, Q,  $t+s_0$ )
- b) schedule next arrival to system  
    sample inter-arrival time  $a_0$   
    schedule (arrival, system,  $t+a_0$ )

# Queueing Example #1

3. Next event is either (departure,  $Q$ ,  $t+s_0$ ) or (arrival, system,  $t+a_0$ )

Suppose  $t+s_0 < t+a_0$

## **Processing of (departure, $Q$ , $t'$ )**

compute statistics associated with departing job

remove job from  $Q$

if  $Q$  not empty

    schedule next departure

        sample service time  $s_1$

        schedule (departure,  $Q$ ,  $t'+s_1$ )

# Event View of DES

## Key concepts

- Event specifying what/where/when to do
- Objects where activities occur
- Event processing logic
  - Changes state
  - Schedules future events
- Event list
  - Data structure supporting efficient
    - » Insertion of new events
    - » Removal of event with least time-stamp
    - » Deletion of events



# Queueing Example #2

Same as previous example, except

- Jobs arrive to queue with pre-declared service requirements
- Queueing policy at each queue is to give service to the job with least remaining service time (shortest-job-first preemptive-resume)
  - Implies that a job in service may be suspended

Changes to previous logic

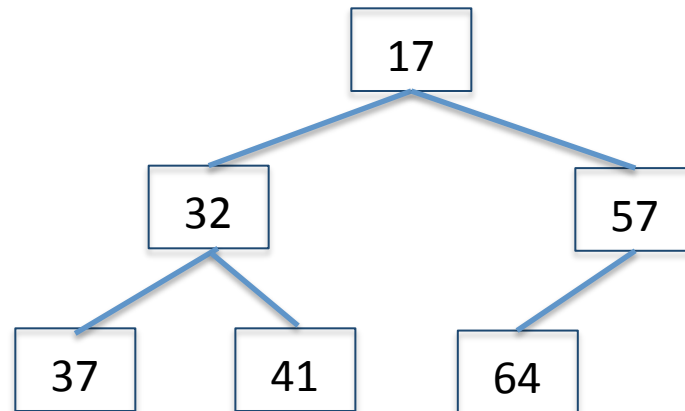
- Event data structure should identify job with known (residual) service time  
e.g., (arrival, system, job0, t), (departure, Q, jobk, t')
- Computation of expected finishing time (and selection of Q) based on queueing policy
- Processing of job arrival to system needs to allow for interrupting a job in service
  - Means **cancelling** a scheduled departure event (departure, Q, jobk, T)

# Event List data structures

- Sorted list of events
  - Get minimum event in  $O(1)$  time, insertion is  $O(N)$
  - Splay-tree
    - All operations are  $O(\log N)$
  - Calendar queue
    - With right parameters,  $O(1)$  operations
      - Degrades with wrong parameters
  - Binary heap

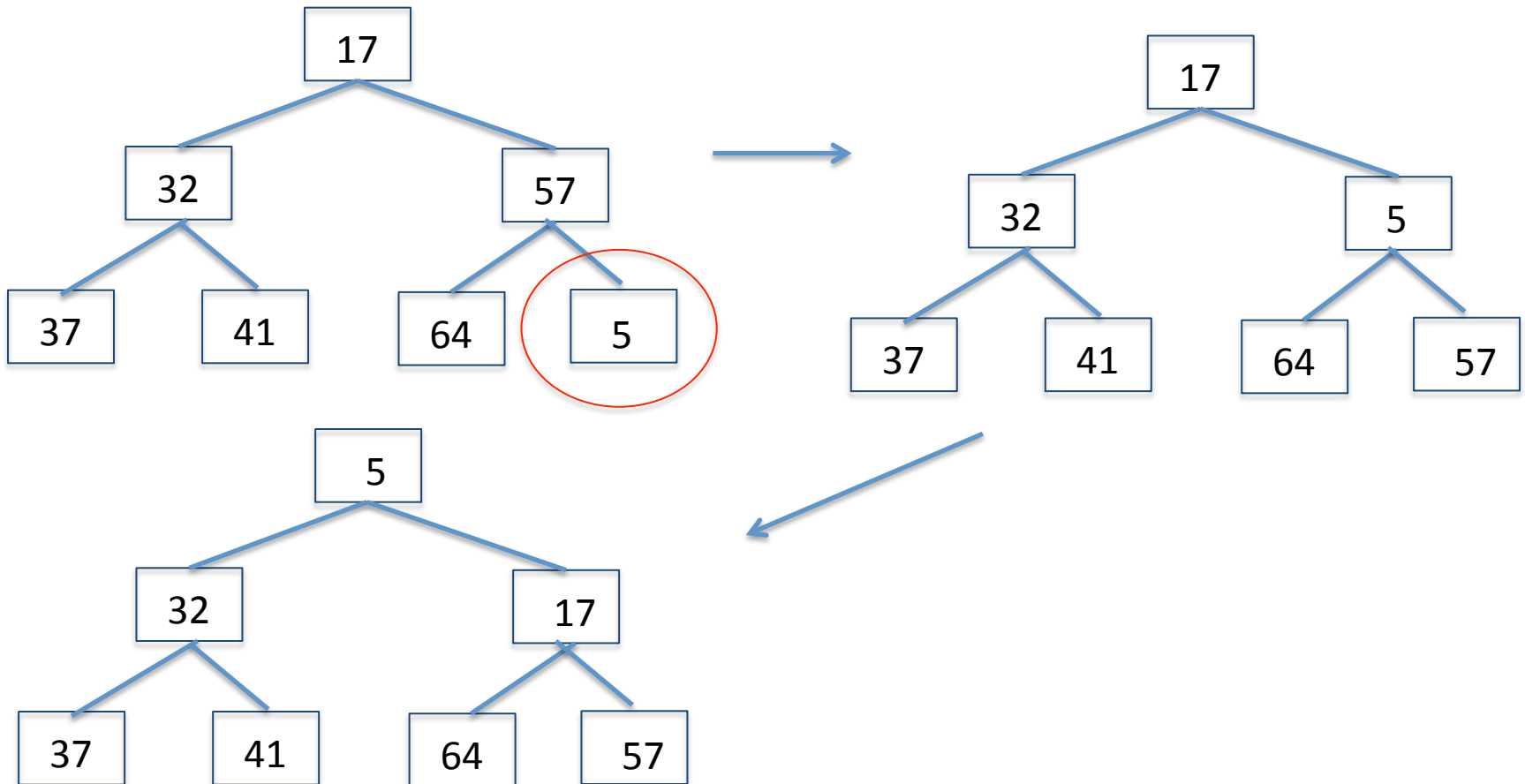
# Event List data structure in python

- Basic idea is binary min-heap, implemented in python with `heapq`
- Constant time identification of event with least time-stamp
- Logarithmic time cost for inserting event
- Maintaining 'heap property'
  - Key of parent node always no smaller than key of child node



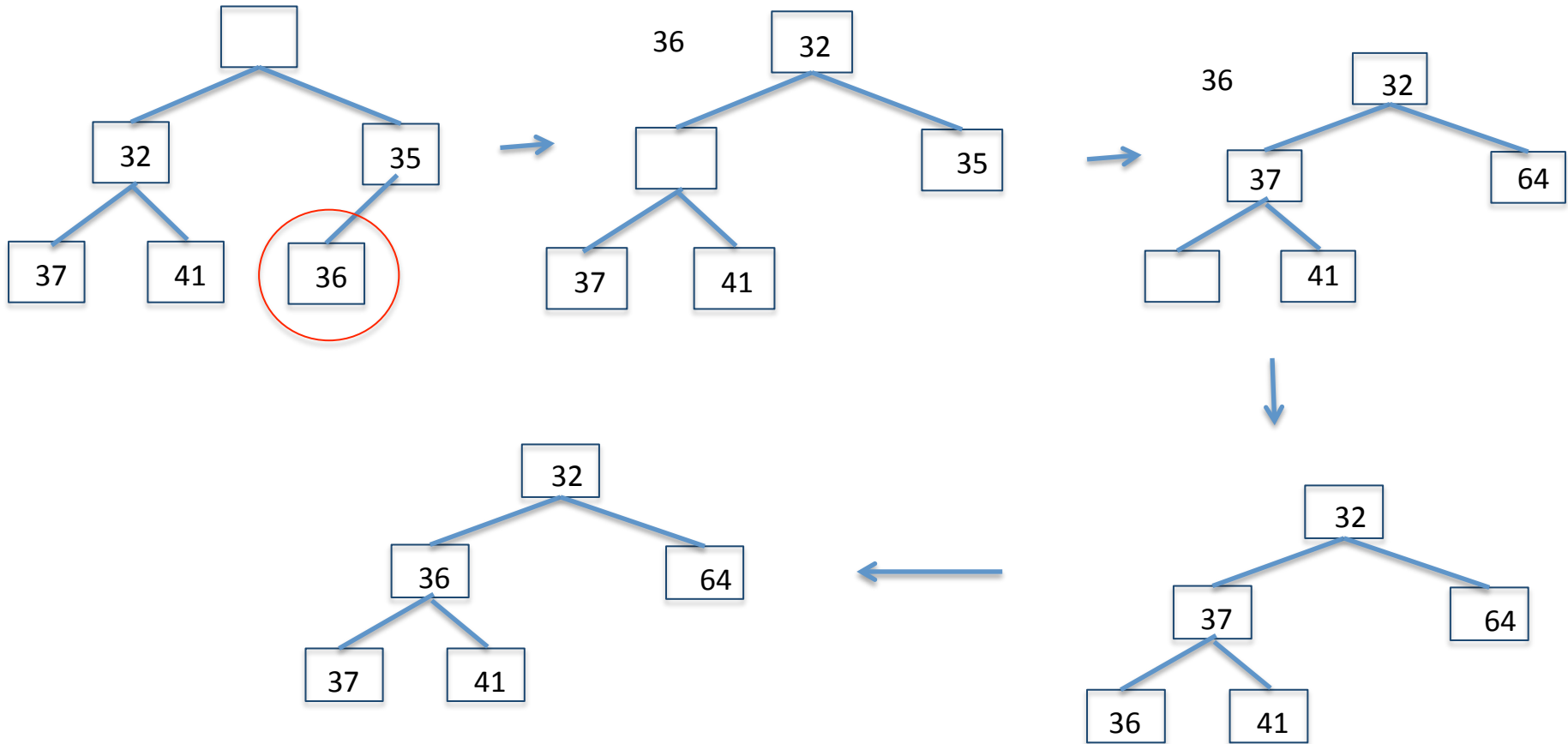
# Event List data structure in python

- Adding element
  - Put new element in 'last place', then swap child and parent towards root to maintain heap property



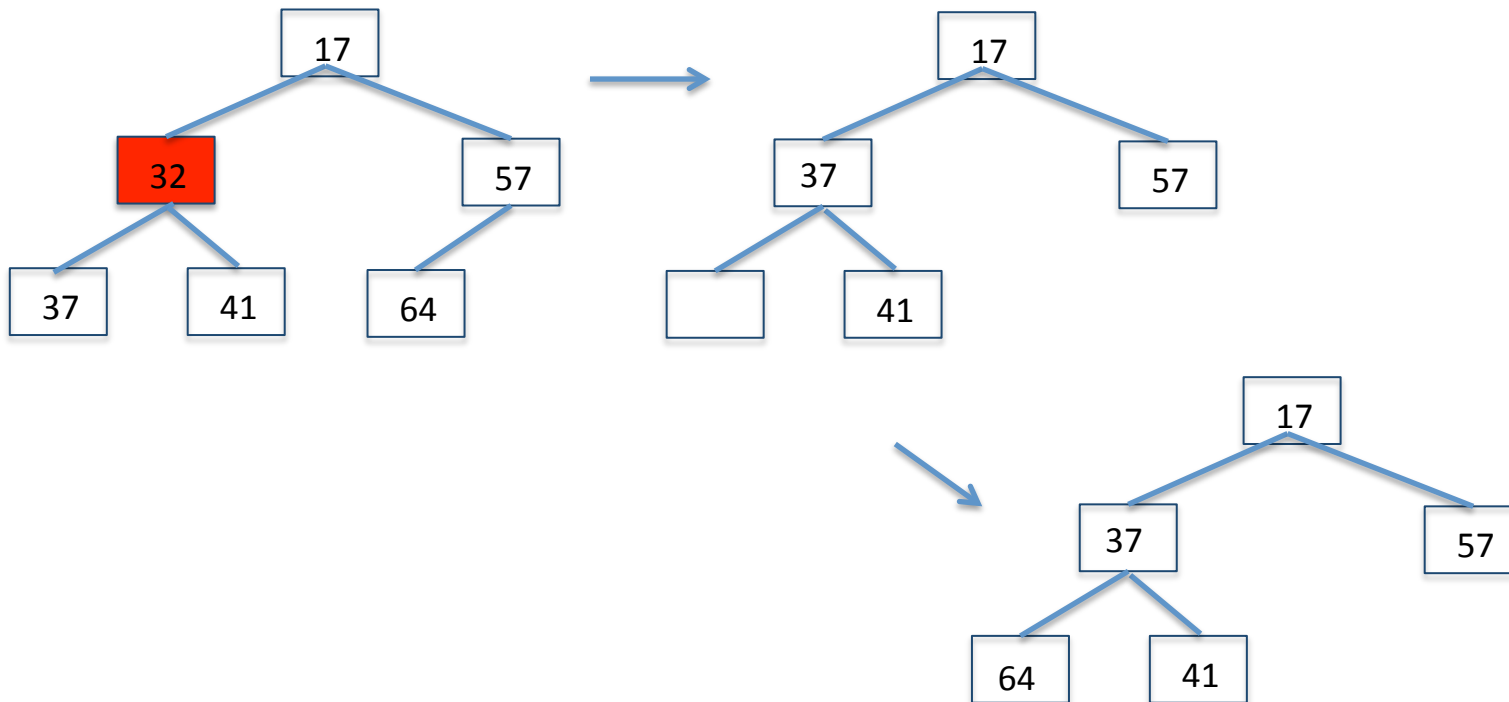
# Event List data structure in python

- Removing front element
  - Take last item from heap, move smaller children up to maintain heap property
  - Insert formerly last item from heap



# Event List data structure in python

- Removing interior element
  - Same operation as removing root, but removing root of subtree whose root is the deleted location



# Implementing Event Cancellation

An event in a min-heap will move position as the simulation progresses

- Need a means of quickly identifying where it is in the heap

`sim_evt.py` has code for

- Creating an event list
- Adding an event to the event list
- Removing least-time event from the event list
- Cancelling a previously added event from the event list
- Modified `heapq` to support cancellation

For event cancellation the API returns a code for a scheduled event, which is offered when that event is to be cancelled

# In the weeds

Go through `sim_evt.py`

Application to a model

Go through `balanced_queue.py`