Distributed Systems

CS425/ECE428

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Acknowledgements for the materials: Indy Gupta

Logistics

- HW6 is due tomorrow (Thursday, Apr 28).
- MP3 is due next week, May 5th.
- Final exam on May 11th
  - Please register on CBTF.
  - Same format as midterms, but longer (3 hours).
  - **Comprehensive:** includes everything covered in the course.
  - ~50% weightage assigned to materials that were not covered by midterm 1 and midterm 2 syllabus (i.e. blockchains and beyond).

Grade distribution

	3-credit	4-credit
Homework	33%	16% (drop 2 worst HWs)
Midterms	33%	25%
Final	33%	25%
MPs	N/A	33%
Participation	1%	1%

Grading

- Midterm curving formula (tentative)
  - absolute: 100 * your score/ total score
  - relative: 80 + 10*(your score – avg_UG_score) / standard_dev
  - We will use max(absolute, relative) to get final score out of 100.
  - Midterm 1:
    - avg_UG_score = 55.43 (out of 70)
    - standard_dev = 8.24
  - Midterm 2:
    - avg_UG_score = 43.13 (out of 65)
    - standard_dev = 9.72
  - Multiply the final score (out of 100) for each midterm by:
    - 0.165 for 3-credit students
    - 0.125 for 4-credit students.
- Finals will be similarly curved, but has higher weightage.

Grading

- Homeworks will not be curved.
  - For 3-credit students:
    - (sum of all 6 homework scores) * 100 * 0.33 / 240
  - For 4-credit students:
    - (sum of best 4 homework scores) * 100 * 0.16 / 160
- MPs will not be curved.
  - (sum of all four MP scores) * 100 * 0.33 / 330
- Participation score: directly taken from Campuswire
  - if reported score > 100, you get full 1%
  - Else you get (reported score /100)%

Final Grades

- Tentative mapping from score to grade (*rough estimate*):
  - Cutoff for B: 80%
  - Bump up a grade for each 4% leap above 80%.
    - B+ 84%, A- 88%, A 92%, A+ 96%
  - Bump down a grade for each 4% leap below 80%
    - B- 76%, C+ 72%, …..

- This is subject to change!

Our agenda

- Brief overview of key-value stores
- Distributed Hash Tables
  - Peer-to-peer protocol for efficient insertion and retrieval of key-value pairs.
- Key-value stores in the cloud
  - How to run large-scale distributed computations over key-value stores?
    - Map-Reduce Programming Abstraction
  - How to design a large-scale distributed key-value store?
    - Case-study: Facebook's Cassandra

Today's focus

- Brief overview of key-value stores
- Distributed Hash Tables
    - Peer-to-peer protocol for efficient insertion and retrieval of key-value pairs.
- Key-value stores in the cloud
    - How to run large-scale distributed computations over key-value stores?
        - Map-Reduce Programming Abstraction
    - How to design a large-scale distributed key-value store?
        - Case-study: Facebook's Cassandra

Distributed datastores

- Distributed datastores
  - Service for managing distributed storage.
- Distributed NoSQL key-value stores
  - BigTable by Google
  - HBase open-sourced by Yahoo and used by Hadoop.
  - DynamoDB by Amazon
  - Cassandra by Facebook
  - Voldemort by LinkedIn
  - MongoDB,
  - …
- *Spanner is not a NoSQL datastore. It's more like a distributed relational database.*

Key-value/NoSQL Data Model

- NoSQL = "Not Only SQL"
- Necessary API operations: get(key) and put(key, value)
  - And some extended operations, e.g., "CQL" in Cassandra key-value store
- Tables
  - Like RDBMS tables, but …
  - May be unstructured: May not have schemas
    - Some columns may be missing from some rows
  - Don't always support joins or have foreign keys
  - Can have index tables, just like RDBMSs

How to design a distributed key-value datastore?

Design Requirements

- High performance, low cost, and scalability.
  - Speed (high throughput and low latency for read/write)
  - Low TCO (total cost of operation)
  - Fewer system administrators
  - Incremental scalability
    - Scale out: add more machines.
    - Scale up: upgrade to powerful machines.
    - *Cheaper to scale out than to scale up.*

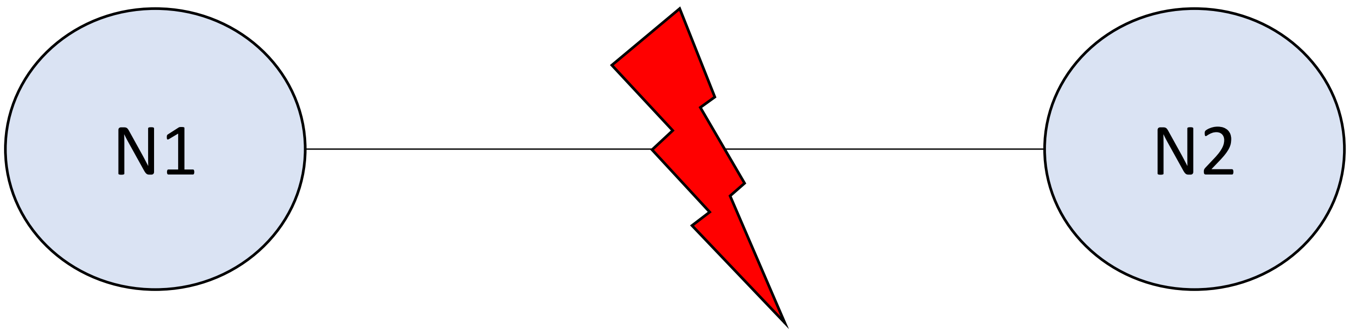
Design Requirements

- High performance, low cost, and scalability.
- Avoid single-point of failure
  - Replication across multiple nodes.
- Consistency: reads return latest written value by any client (all nodes see same data at any time).
  - *Different from the C of ACID properties for transaction semantics!*
- Availability: every request received by a non-failing node in the system must result in a response (quickly).
  - Follows from requirement for high performance.
- Partition-tolerance: the system continues to work in spite of network partitions.

CAP Theorem

- **C**onsistency**:** reads return latest written value by any client (all nodes see same data at any time).
- **A**vailability**:** every request received by a non-failing node in the system must result in a response (quickly).
- **P**artition-tolerance**:** the system continues to work in spite of network partitions.
- **In a distributed system you can only guarantee at most 2 out of the above 3 properties.**
  - Proposed by Eric Brewer (UC Berkeley)
  - Subsequently proved by Gilbert and Lynch (NUS and MIT)

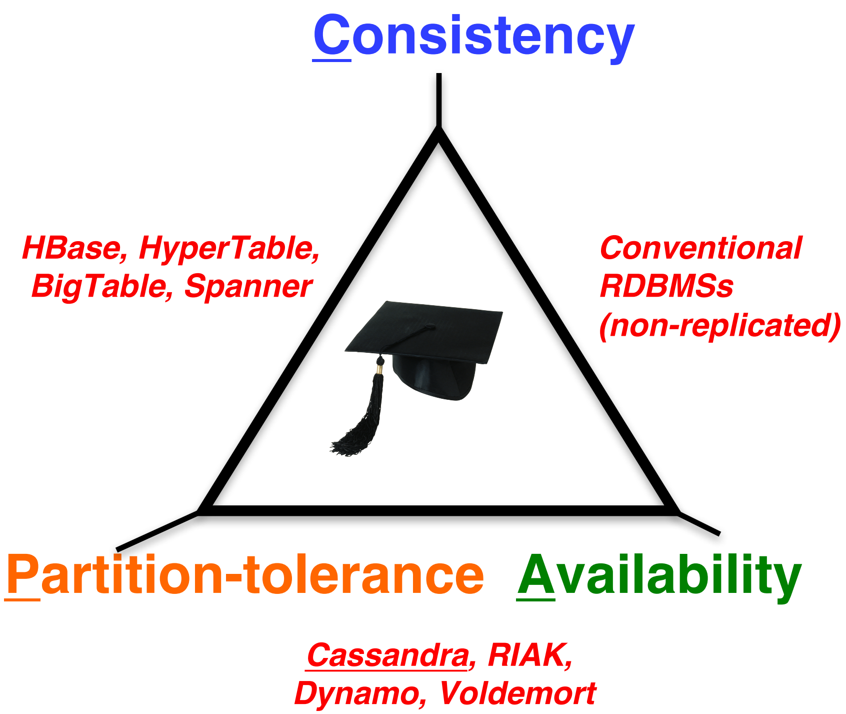
CAP Theorem



- Data replicated across both N1 and N2.
- If network is partitioned, N1 can no longer talk to N2.
- Consistency + availability require N1 and N2 must talk.
  - no partition-tolerance.
- Partition-tolerance + consistency:
  - only respond to requests received at N1 (no availability).
- Partition-tolerance + availability:
  - write at N1 will not be captured by a read at N2 (no consistency).

CAP Tradeoff

- Starting point for NoSQL Revolution
- A distributed storage system can achieve at most two of C, A, and P.
- When partition-tolerance is important, you have to choose between consistency and availability



Modern key-value stores vs. RDBMS

- While RDBMS provide ACID
  - Atomicity
  - Consistency
  - Isolation
  - Durability

- Many modern key-value stores provide BASE
  - Basically Available Soft-state Eventual Consistency
  - Prefers Availability over Consistency

Case Study: Cassandra

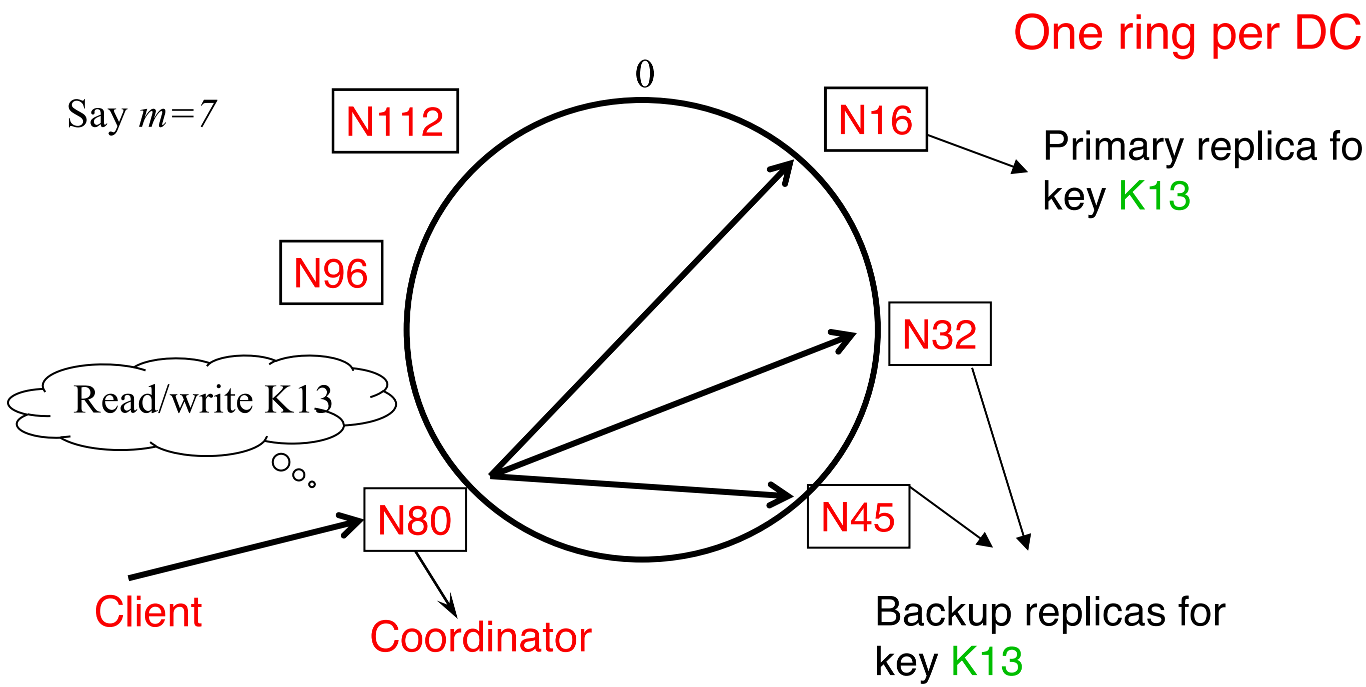
Cassandra

- A distributed key-value store.
- Intended to run in a datacenter (and also across DCs).
- Originally designed at Facebook.
- Open-sourced later, today an Apache project.
- Some of the companies that use Cassandra in their production clusters.
  - IBM, Adobe, HP, eBay, Ericsson, Symantec
  - Twitter, Spotify
  - PBS Kids
  - Netflix: uses Cassandra to keep track of your current position in the video you're watching

Data Partitioning: Key to Server Mapping

- How do you decide which server(s) a key-value resides on?

Cassandra uses a ring-based DHT but without finger or routing tables.



Partitioner

- Component responsible for key to server mapping (hash function).

- Two types:
  - *Chord-like hash partitioning*
    - *Murmer3Partitioner* (default): uses *murmer3* hash function.
    - *RandomPartitioner*: uses MD5 hash function.
  - *ByteOrderedPartitioner*: Assigns ranges of keys to servers.
    - Easier for *range queries* (e.g., get me all twitter users starting with [a-b])

- Determines the primary replica for a key.

Replication Policies

Two options for replication strategy:

1. SimpleStrategy:
    - First replica placed based on the partitioner.
    - Remaining replicas clockwise in relation to the primary replica.
2. NetworkTopologyStrategy: for multi-DC deployments
    - Two or three replicas per DC.
    - Per DC
        - First replica placed according to Partitioner.
        - Then go clockwise around ring until you hit a different rack.

Writes

- Need to be lock-free and fast (no reads or disk seeks).
- Client sends write to one coordinator node in Cassandra cluster.
  - Coordinator may be per-key, or per-client, or per-query.
- Coordinator uses Partitioner to send query to all replica nodes responsible for key.
- When X replicas respond, coordinator returns an acknowledgement to the client
  - X = any one, majority, all….(consistency spectrum)
  - More details later!

Writes: Hinted Handoff

- Always writable: Hinted Handoff mechanism
  - If any replica is down, the coordinator writes to all other replicas, and keeps the write locally until down replica comes back up.
  - When all replicas are down, the Coordinator (front end) buffers writes (for up to a few hours).

Writes at a replica node

On receiving a write

1. Log it in disk commit log (for failure recovery)
2. Make changes to appropriate memtables
   - **Memtable** = In-memory representation of multiple key-value pairs
   - Cache that can be searched by key
   - Write-back cache as opposed to write-through
3. Later, when memtable is full or old, flush to disk
   - Data File: An **SSTable** (Sorted String Table) – list of key-value pairs, sorted by key
   - Index file: An SSTable of (key, position in data sstable) pairs
   - And a Bloom filter (for efficient search) – next slide.

To be continued in next class

- Wrap up writes.
- Reads.
- Cluster membership.
- Eventual consistency model.