

# Distributed Systems

CS425/ECE428

April 23 2021

*Instructor: Radhika Mittal*

*Acknowledgements for some of the materials: Indy Gupta, Nikita Borisov*

# Our agenda for the next 3-4 classes

- Brief overview of key-value stores
- Distributed Hash Tables
  - Peer-to-peer protocol for efficient insertion and retrieval of key-value pairs.
- Key-value stores in the cloud
  - How to run large-scale distributed computations over key-value stores?
    - Map-Reduce Programming Abstraction
  - How to design a large-scale distributed key-value store?
    - Case-study: Facebook's Cassandra

# Features of cloud

## I. Massive scale.

- Tens of thousands of servers and cloud tenants, and hundreds of thousands of VMs.

## II. On-demand access:

- Pay-as-you-go, no upfront commitment, access to anyone.

## III. Data-intensive nature:

- What was MBs has now become TBs, PBs and XBs.
  - Daily logs, forensics, Web data, etc.

# Must deal with immense complexity!

- Fault-tolerance and failure-handling
- Replication and consensus
- Cluster scheduling
  
- How would a cloud user deal with such complexity?
  - **Powerful abstractions and frameworks**
  - Provide **easy-to-use** API to users.
  - Deal with the complexity of distributed computing under the hood.

**MapReduce**  
is one such powerful  
abstraction.

# MapReduce Abstraction

- Map/Reduce
  - Programming model inspired from LISP (and other functional languages).
- Expressive: many problems can be phrased as map/reduce.
- Easy to distribute across nodes.
  - High-level job divided into multiple independent “map” tasks, followed by multiple independent “reduce” tasks.
- Nice retry/failure semantics.

# MapReduce Architecture

- *MapReduce programming abstraction:*
  - Easy to program distributed computing tasks.
- MapReduce programming abstraction offered by multiple open-source *application frameworks*:
  - Handle creation of “map” and “reduce” tasks.
  - e.g. *Hadoop: one of the earliest map-reduce frameworks.*
  - e.g. *Spark: easier API and performance optimizations.*
- Application frameworks use *resource managers*.
  - Deal with the hassle of distributed cluster management.
  - e.g. *Kubernetes, YARN, Mesos, etc.*

# MapReduce Architecture

- *Map/Reduce abstraction:*
  - Easy to program distributed computing tasks.
- MapReduce features:
  - Automatic parallelization & distribution
  - Fault tolerance
  - Scheduling
  - Monitoring & status updates
- Application frameworks use *resource managers*.
  - Deal with the hassle of distributed cluster management.
  - e.g. *Kubernetes, YARN, Mesos, etc.*



# Map/Reduce in LISP

Sum of squares:

- `(map square '(1 2 3 4))`
  - Output: `(1 4 9 16)`
  
- `(reduce + 0 '(1 4 9 16))`
  - `(+ 16 (+ 9 (+ 4 (+ 1 + 0) ) ) )`
  - Output: 30

# Map/Reduce in LISP

Sum of squares:

- `(map square '(1 2 3 4))` Unary operator
  - Output: (1 4 9 16)
  - [processes each record sequentially and independently]
- `(reduce + 0 '(1 4 9 16))`
  - `(+ 16 (+ 9 (+ 4 (+ 1 + 0) ) ) )`
  - Output: 30

# Map/Reduce in LISP

Sum of squares:

- `(map square '(1 2 3 4))`

Unary operator

- Output: `(1 4 9 16)`

[processes each record sequentially and independently]

- `(reduce + 0 '(1 4 9 16))`

Binary operator

- `(+ 16 (+ 9 (+ 4 (+ 1 + 0) ) ) )`

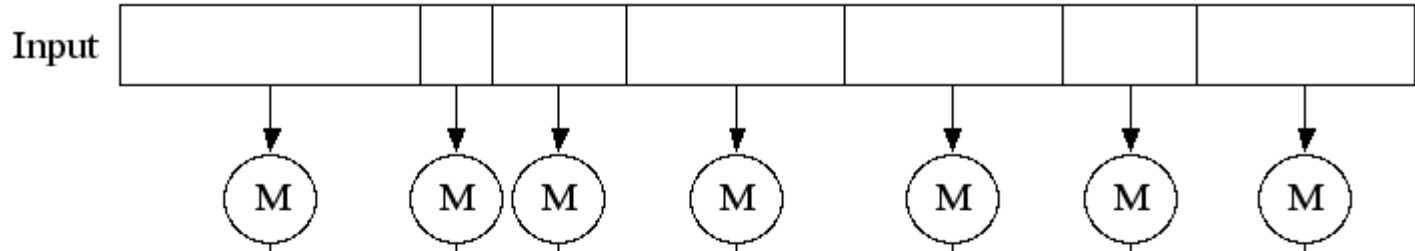
- Output: 30

[processes set of *all* records in batches]

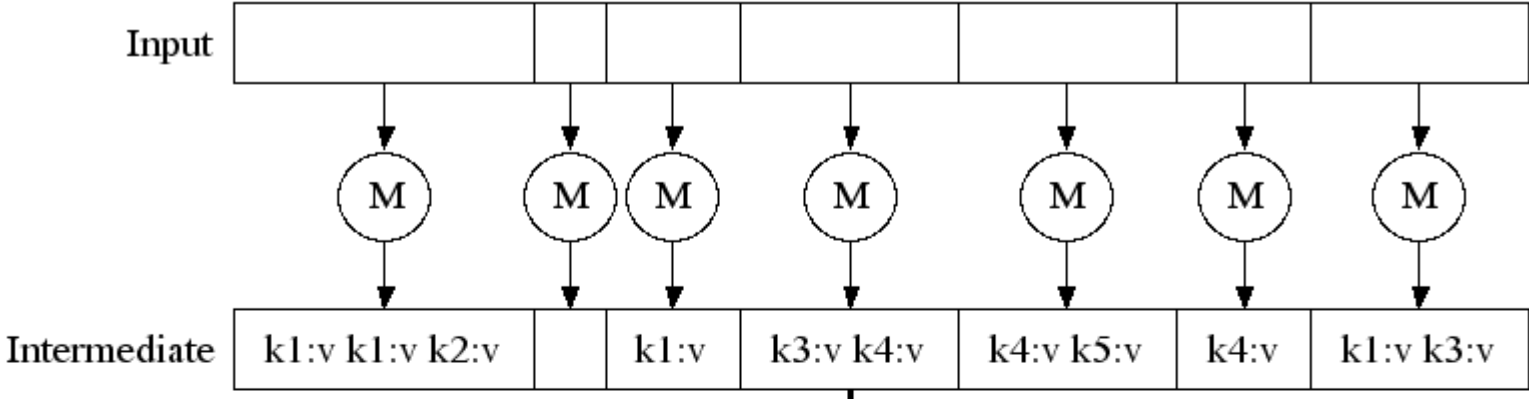
# MapReduce Overview

- Input: a set of key/value pairs
- User supplies two functions:
  - $\text{map}(k,v) \rightarrow \text{list}(k1,v1)$
  - $\text{reduce}(k1, \text{list}(v1)) \rightarrow v2$
- $(k1,v1)$  is an intermediate key/value pair.
- Output is the set of  $(k1,v2)$  pairs.

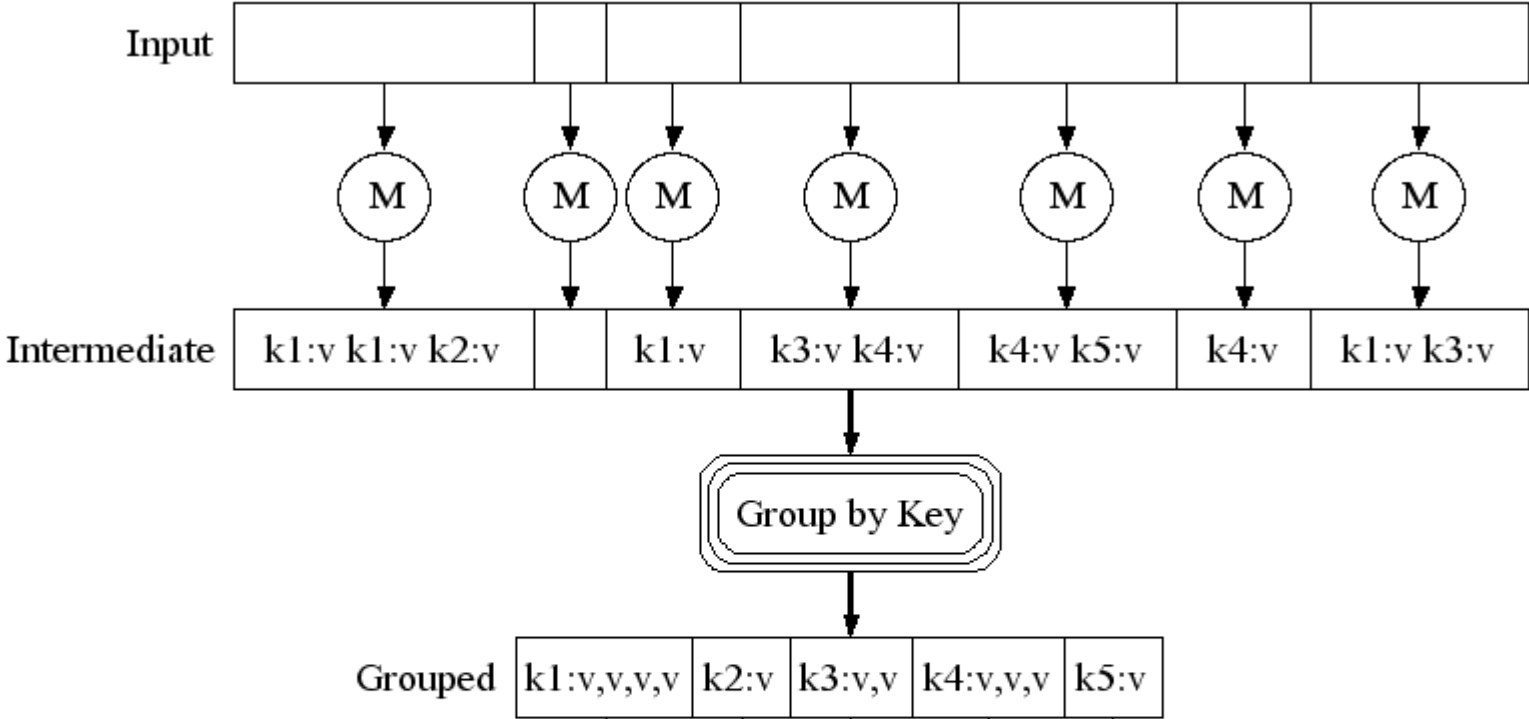
# MapReduce Overview



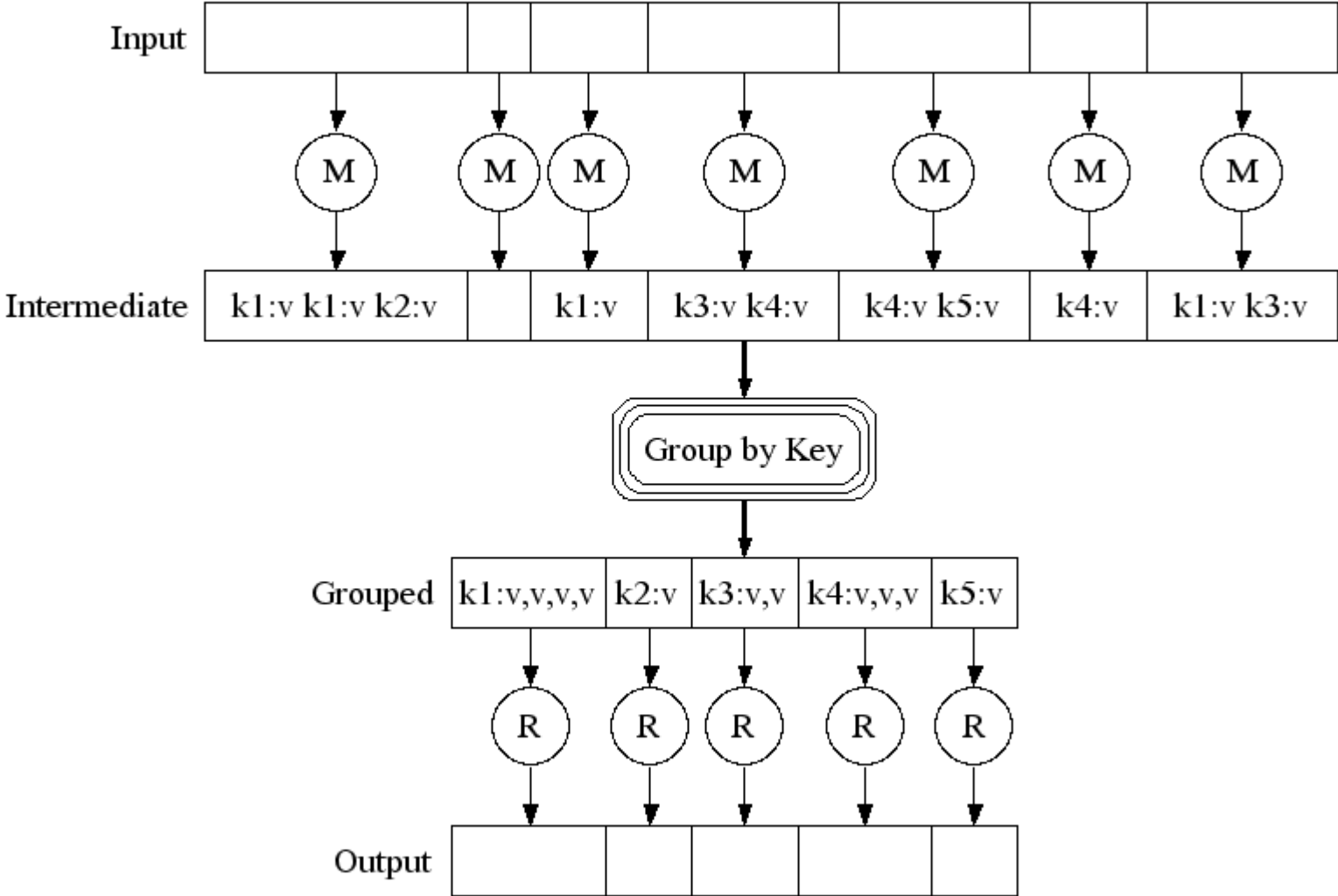
# MapReduce Overview



# MapReduce Overview



# MapReduce Overview





# Typical Example: Word Count

- We have a large file of words containing multiple lines (or records).
- Count the number of times each distinct word appears in the file.
- *Sample application:* analyze web server logs to find popular URLs.

# Map

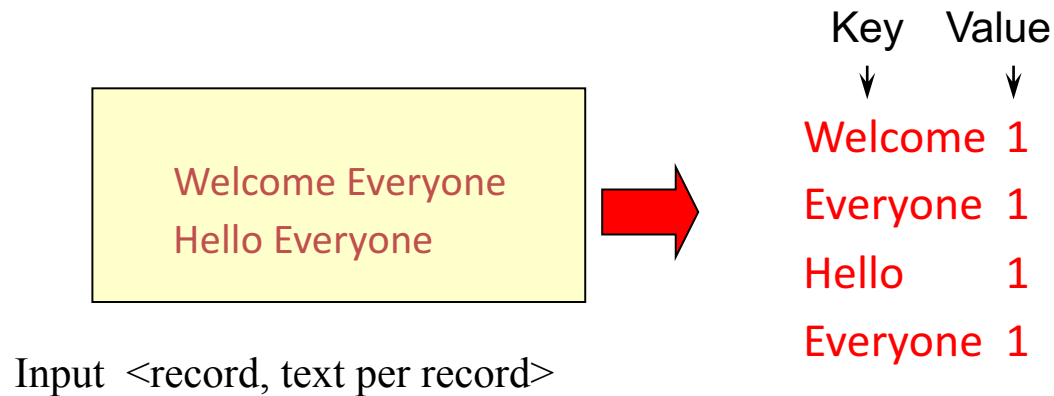
- Process individual records to generate *intermediate key/value pairs*.



Input <record, text per record>

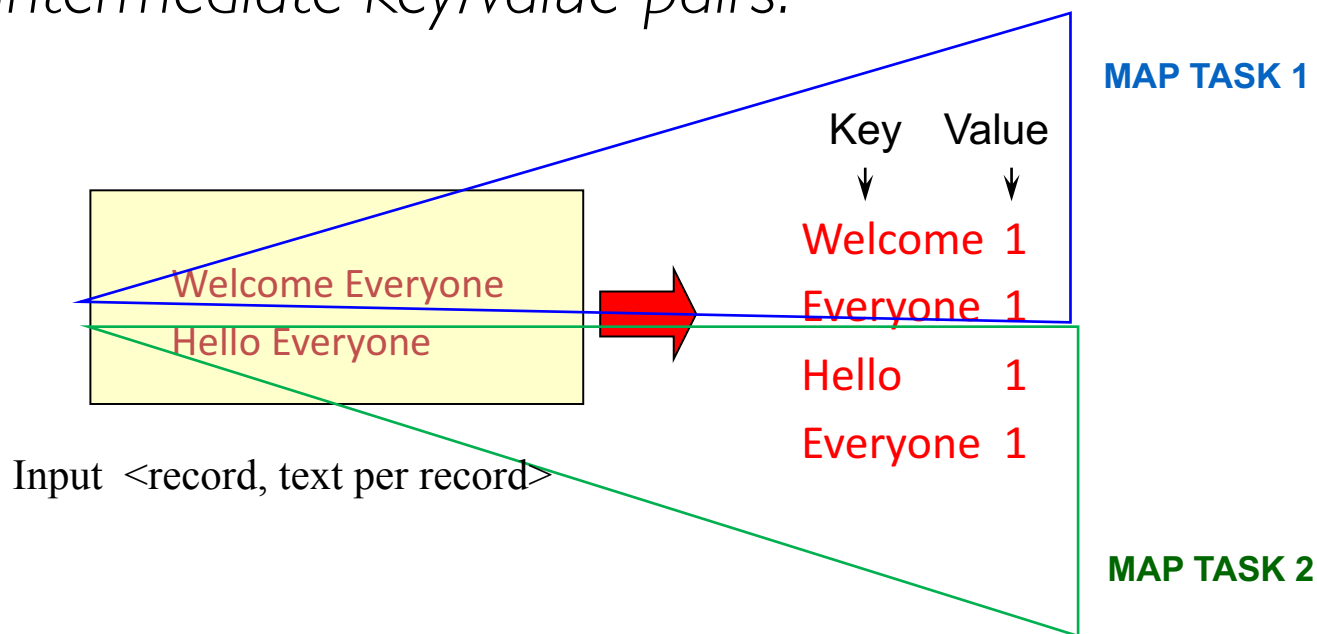
# Map

- Process individual records to generate *intermediate key/value pairs*.



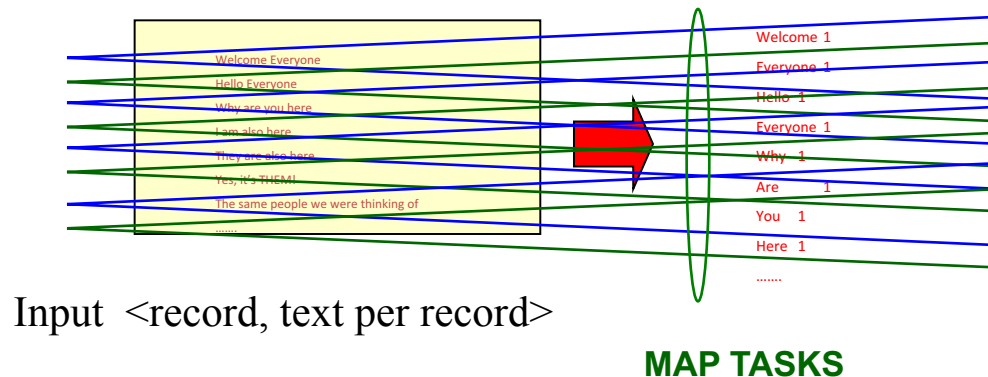
# Map

- **Parallely** process individual records to generate *intermediate key/value pairs*.



# Map

- **Parallely** process **a large number** of individual records to generate intermediate key/value pairs.



# Reduce

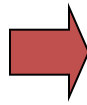
- Processes and merges all intermediate values associated per key.

Welcome 1

Everyone 1

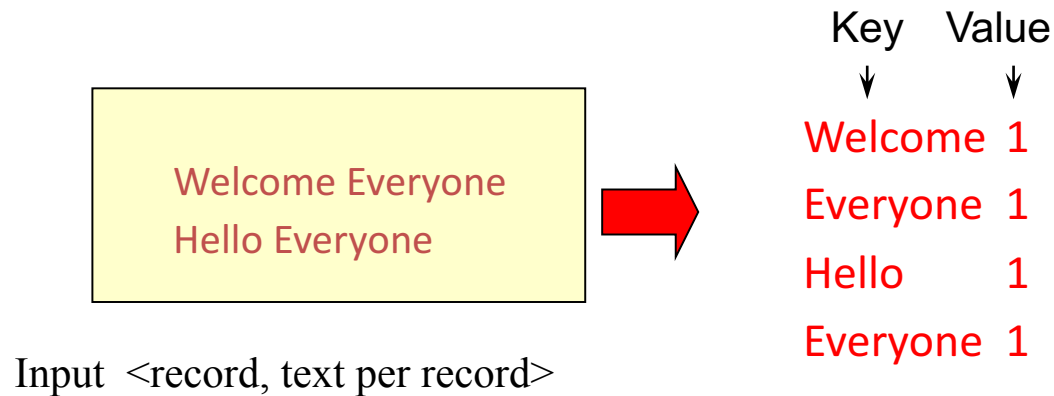
Hello 1

Everyone 1



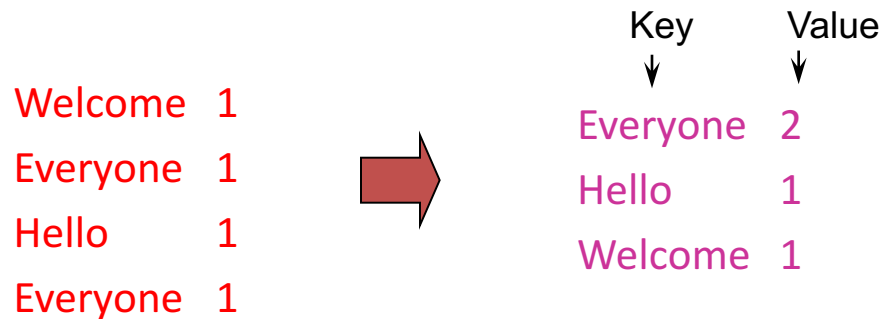
# Map

- Process individual records to generate *intermediate key/value pairs*.



# Reduce

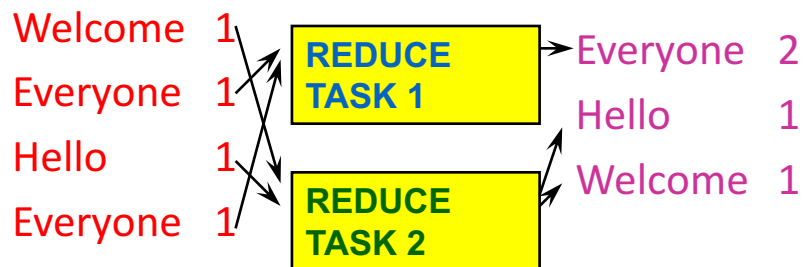
- Processes and merges all intermediate values associated per key.





# Reduce

- Each key assigned to one Reduce task.
- **Parallely** processes and merges all intermediate values partitioned per key.



- Popular: *Hash partitioning*, i.e., key is assigned to
  - $\text{reduce \#} = \text{hash}(\text{key}) \% \text{number of reduce tasks}$

# MapReduce Overview

- Input: a set of key/value pairs
- User supplies two functions:
  - $\text{map}(k,v) \rightarrow \text{list}(k1,v1)$
  - $\text{reduce}(k1, \text{list}(v1)) \rightarrow v2$
- $(k1,v1)$  is an intermediate key/value pair.
- Output is the set of  $(k1,v2)$  pairs.

# MapReduce Overview

- Input: a set of key/value pairs (record, list of words)
- User supplies two functions:
  - $\text{map}(k, v) \rightarrow \text{list}(k1, v1)$
  - $\text{reduce}(k1, \text{list}(v1)) \rightarrow v2$
- $(k1, v1)$  is an intermediate key/value pair. (word, 1)
- Output is the set of  $(k1, v2)$  pairs. (word, count)

# Word Count using MapReduce

```
map(key, value):
```

```
// key: record (line no.); value: list of words in the record
```

```
  for each word w in value:
```

```
    emit(w, 1)
```

```
reduce(key, values):
```

```
// key: a word; values: an iterator over counts
```

```
  result = 0
```

```
  for each count v in values:
```

```
    result += v
```

```
  emit(key, result)
```

# Hadoop Code - Map

```
public static class MapClass extends MapReduceBase

    implements Mapper<LongWritable, Text, Text, IntWritable> {

    private final static IntWritable one =

        new IntWritable(1);

    private Text word = new Text();

    public void map( LongWritable key, Text value,

        OutputCollector<Text, IntWritable> output, Reporter reporter) // key is empty, value is the line

        throws IOException {

        String line = value.toString();

        StringTokenizer itr = new StringTokenizer(line);

        while (itr.hasMoreTokens()) {

            word.set(itr.nextToken());

            output.collect(word, one);

        }

    }

} // Source: http://developer.yahoo.com/hadoop/tutorial/module4.html#wordcount
```

# Hadoop Code - Reduce

```
public static class ReduceClass extends MapReduceBase           implements
Reducer<Text, IntWritable, Text, IntWritable> {

    public void reduce(

        Text key,

        Iterator<IntWritable> values,

        OutputCollector<Text, IntWritable> output,

        Reporter reporter)

            throws IOException {

        // key is word, values is a list of 1's

        int sum = 0;

        while (values.hasNext()) {

            sum += values.next().get();

        }

        output.collect(key, new IntWritable(sum));

    }

} // Source: http://developer.yahoo.com/hadoop/tutorial/module4.html#wordcount
```

# Hadoop Code - Driver

```
// Tells Hadoop how to run your Map-Reduce job

public void run (String inputPath, String outputPath)
    throws Exception {

    // The job. WordCount contains MapClass and Reduce.
    JobConf conf = new JobConf(WordCount.class);

    conf.setJobName("mywordcount");

    // The keys are words
    (strings) conf.setOutputKeyClass(Text.class);

    // The values are counts (ints)
    conf.setOutputValueClass(IntWritable.class);

    conf.setMapperClass (MapClass.class);

    conf.setReducerClass (ReduceClass.class);

    FileInputFormat.addInputPath(
        conf, newPath(inputPath));

    FileOutputFormat.setOutputPath(
        conf, new Path(outputPath));

    JobClient.runJob(conf);
} // Source: http://developer.yahoo.com/hadoop/tutorial/module4.html#wordcount
```

# Spark Code

Python:

```
text_file = sc.textFile("hdfs://...")
counts = text_file.flatMap \
    (lambda line: line.split(" ")) \
    .map(lambda word: (word, 1)) \
    .reduceByKey(lambda a, b: a + b)
counts.saveAsTextFile("hdfs://...")
```

// Source: <http://spark.apache.org/examples.html>



# More examples: Host size

- Suppose we have a large web corpus
- Metadata file
  - Lines of the form (URL, size, date, ...)
- For each host, find the total number of bytes
  - i.e., the sum of the sizes for all pages from a given host/URL

```
map(key, value):  
// key: metadata record#;  
//value: (URL, size, ...) :  
    for each (URL, size) in value:  
        emit(URL, size)
```

```
reduce(key, values):  
// key: URL, values: iterator over sizes:  
result = 0  
for each size s in values:  
    result += s  
emit(key, result)
```

# More examples: Distributed Grep

- Input: large set of files
- Output: unique lines that match pattern

```
map(key, value):
```

```
// key: file, value: list of lines
```

```
  for each line in value:
```

```
    if "pattern" in line:
```

```
      emit(line, 1)
```

```
reduce(key, values):
```

```
// key: line that matches pattern; values: 1's
```

```
  emit(key, 1)
```

# More examples: Graph reversal

- Input: Web graph: tuples (a, b) where (page a → page b)
- Output: For each page, list of pages that link to it

map(key, value):

// key: source page,

//value: target page

emit(value, key)

reduce(key, values):

// key: target; values: list of pages that link to it.

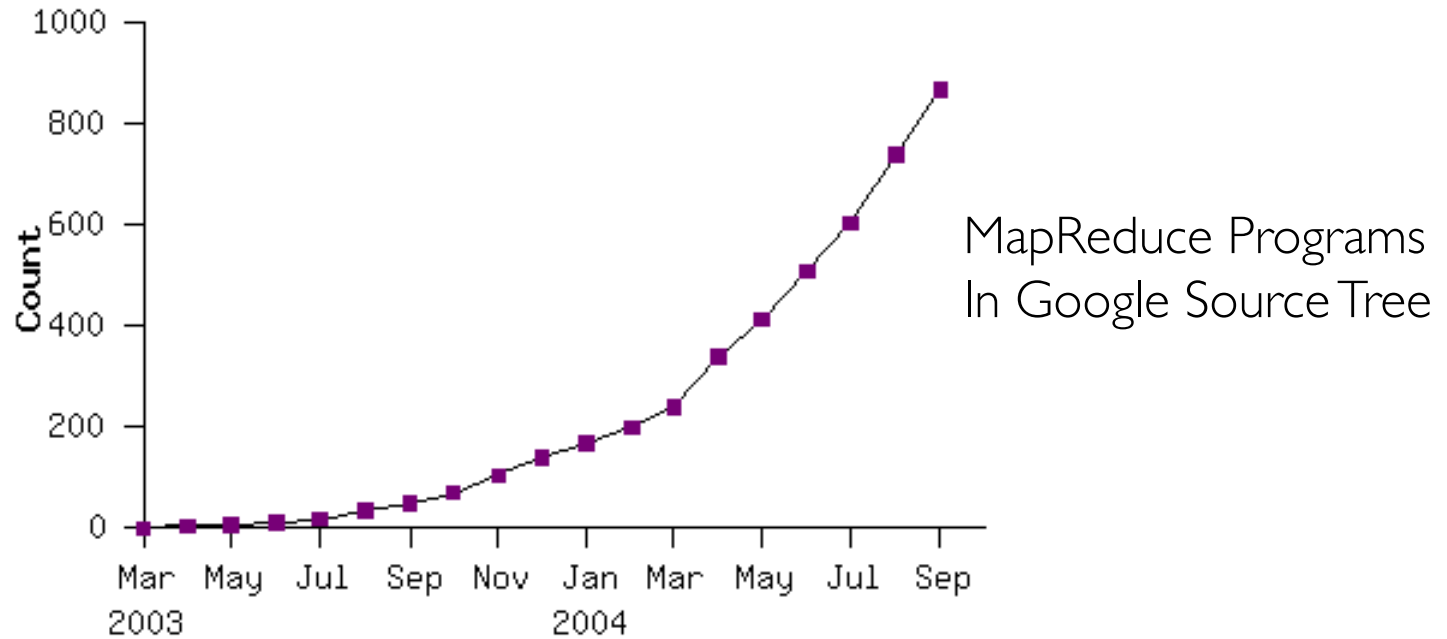
result = concatenate(values)

emit(key, result)

# MapReduce Chains

- map1 -> reduce1 -> map2 -> reduce2
- E.g., output the most common words by frequency
  - Map1: emit ("word", 1)
  - Reduce1: emit ("word", count)
  - Map2: emit (count, "word")
  - Reduce2: identity, i.e. emit(count, list of words)

# MapReduce is popular and widely applicable



## Example uses:

distributed grep

term-vector / host

document clustering

...

distributed sort

web access log stats

machine learning

...

web link-graph reversal

inverted index construction

statistical machine  
translation

...

# MapReduce Execution

Externally: For user

1. Write a Map program (short), write a Reduce program (short)
2. Specify number of Maps and Reduces (parallelism level)
3. Submit job; wait for result
4. Need to know very little about parallel/distributed programming!

# MapReduce Execution

Internally: For the framework and resource manager in the cloud

1. Parallelize Map
2. Transfer data from Map to Reduce (**shuffle data**)
3. Parallelize Reduce
4. Implement Storage for Map input, Map output, Reduce input, and Reduce output

(Ensure that no Reduce starts before all Maps are finished. That is, ensure the **barrier** between the Map phase and Reduce phase)

# MapReduce Execution

Internally: For the framework and resource manager in the cloud

1. Parallelize Map (*easy!*)
  - Each map task is independent of the other!
2. Transfer data from Map to Reduce (*shuffle data*)
3. Parallelize Reduce
4. Implement Storage for Map input, Map output, Reduce input, and Reduce output

(Ensure that no Reduce starts before all Maps are finished. That is, ensure the *barrier* between the Map phase and Reduce phase)



# MapReduce Execution

Internally: For the framework and resource manager in the cloud

1. Parallelize Map (*easy!*)
2. Transfer data from Map to Reduce (**shuffle data**)
  - All Map output records with same key assigned to same Reduce
  - Use *partitioning function, e.g.,  $\text{hash}(\text{key})\% \text{number of reducers}$*
3. Parallelize Reduce
4. Implement Storage for Map input, Map output, Reduce input, and Reduce output

(Ensure that no Reduce starts before all Maps are finished. That is, ensure the **barrier** between the Map phase and Reduce phase)

# MapReduce Execution

Internally: For the framework and resource manager in the cloud

1. Parallelize Map (*easy!*)
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4. Implement Storage for Map input, Map output, Reduce input, and Reduce output

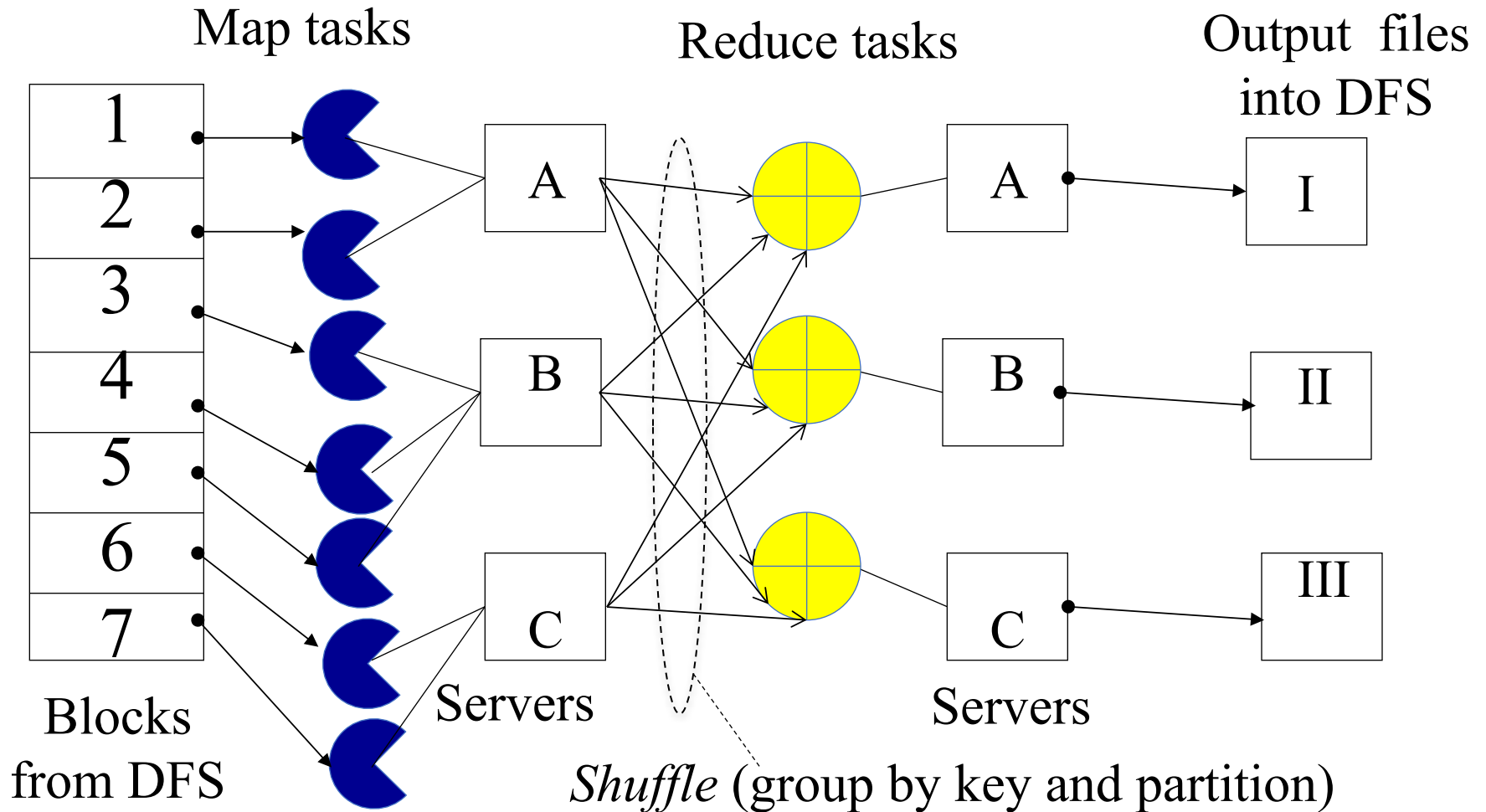
(Ensure that no Reduce starts before all Maps are finished. That is, ensure the **barrier** between the Map phase and Reduce phase)

# MapReduce Execution

Internally: For the framework and resource manager in the cloud

1. Parallelize Map
2. Transfer data from Map to Reduce (**shuffle data**)
3. Parallelize Reduce
4. Implement Storage for Map input, Map output, Reduce input, and Reduce output
  - Map input: from **distributed file system/data store**
  - Map output: to local disk (at Map node); uses **local file system**
  - Reduce input: from (multiple) remote disks; uses **local file systems**
  - Reduce output: to **distributed file system/data store**
    - local file system** (e.g. Linux FS)
    - distributed file system** (e.g. Google File System, Hadoop Distributed File System)
    - distributed data store** (e.g. Cassandra, BigTable, Spanner, DynamoDB)

# MapReduce Execution



Resource Manager (assigns map and reduce tasks to servers)

# Resource Manager

- Examples:
  - *YARN* (Yet Another Resource Negotiator), used underneath Hadoop 2.x +
  - *Kubernetes, Borg, Mesos*, etc.
- Treats each server as a collection of *containers*
  - Container = fixed CPU + fixed memory (e.g. *Docker*)
  - Each tasks runs in a container.
- Has 3 main components
  - *Global Resource Manager (RM)*: Cluster Scheduling
  - *Per-server Node Manager (NM)*: Daemon and server-specific functions
  - *Per-application (job) Application Master (AM)*
    - Container negotiation with RM and NMs.
    - Handling task failures of that job.

# Fault Tolerance

- NM heartbeats to RM
  - If server fails: RM times out waiting for next heartbeat, RM lets all affected AMs know, and AMs take appropriate action.
- NM keeps track of each task running at its server
  - If task fails while in-progress, mark the task as idle and restart it.
- AM heartbeats to RM
  - On failure, RM restarts AM, which then syncs it up with its running tasks.
- RM Failure
  - Use old checkpoints and bring up secondary RM.

# Slow Servers

Slow tasks are called **Stragglers**.

- The slowest task slows the entire job down (why?)
- Due to bad disk, network bandwidth, CPU, or memory
- Keep track of “progress” of each task (% done)
- Perform proactive backup (replicated) execution of some straggler tasks
  - A task considered done when its first replica complete (other replicas can then be killed).
  - Approach called **Speculative Execution**.
- Straggler mitigation has been a very active area of research.

Barrier at the end  
of Map phase!

# Task Scheduling

- Favour data locality:
  - attempts to schedule a map task on a machine that contains a replica of corresponding input data.
  - *if that's not possible*, on the same rack as a machine containing the input.
  - *if that's not possible*, anywhere.
- What does “*if that's not possible*” mean?
  - No more resources available on the machine.
  - Might be worth waiting a while for resources to become available.
    - Delay scheduling in Spark!
- Cluster scheduling is also a very active area of research.



# Summary

- Cloud provides distributed computing infrastructure as a service.
- Running a distributed job on the cloud cluster can be very complex:
  - Must deal with parallelization, scheduling, fault-tolerance, etc.
- MapReduce is a powerful abstraction to hide this complexity.
  - User programming via easy-to-use API.
  - Distributed computing complexity handled by underlying frameworks and resource managers
- Plenty of ongoing research work in scheduling, fault-tolerance, and straggler mitigation for MapReduce.