

### Communication Networks

- Communication networks consist of links over which messages can pass between nodes (or terminals or hosts)
- Messages to distant nodes have to pass over multiple links and multiple nodes
- The intermediate links can fail
- The intermediate nodes can fail
- What is the probability that two nodes can communicate?

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### Links only; we don't do nodes

- A node that fails makes all the links connected to it inoperable
- All paths between other nodes that pass through the failed node are not available
- In the simple analyses presented in this course, we assume that nodes do not fail
- The only failures that we are concerned with are the link failures

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### Communication over parallel links

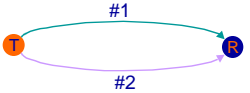
- The diagram illustrates two parallel links connecting a transmitter and a receiver
- The transmitter can send messages as long as at least one link is **viable**
- Both links must **fail** in order for communication to be impossible



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### Some notation

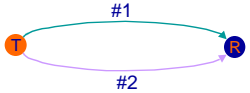
- Let  $V_i$  denote the event that link #i is viable that is, link #i has not failed
- The complementary event is  $F_i$
- $V$  ( $F = V^c$ ) denotes the event that there is (is not) a communication path from T to R



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### As easy as rolling off a log...

- $V = V_1 \cap V_2$       •  $F = F_1 \cup F_2$
- $P(F) = P(F_1 \cup F_2) = P(F_1) + P(F_2) - P(F_1 F_2)$
- **if the events  $F_1$  and  $F_2$  can be assumed to be independent**
- Similarly,  $P(V) = P(V_1) + P(V_2) - P(V_1 V_2)$   
 $= P(V_1) + P(V_2) - P(V_1)P(V_2)$



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### Let's be sure we understand...

- $V = V_1 \cap V_2$       •  $F = F_1 \cup F_2$
- $P(F) = P(F_1 \cup F_2) = P(F_1) + P(F_2) - P(F_1 F_2)$
- **if the events  $F_1$  and  $F_2$  can be assumed to be independent**
- $P(V) = P(V_1) + P(V_2) - P(V_1 V_2)$   
 $= P(V_1) + P(V_2) - P(V_1)P(V_2)$   
 $= q_1 + q_2 - q_1 q_2$   
 $= (1 - p_1) + (1 - p_2) - (1 - p_1)(1 - p_2)$   
 $= 1 - p_1 p_2$

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### Did matters improve?

- $P(F) = P(F_1F_2) = P(F_1)P(F_2) = p_1p_2$  for independent failures
- $P(F) \ll \min\{p_1, p_2\}$
- Example: If  $p_1 = p_2 = 10^{-3}$ , then  $P(F) = 10^{-6}$
- Improvement in  $P(V)$  is not as “dramatic”

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### Path over serial links

- The diagram illustrates two serial links connecting a transmitter and a receiver
- The transmitter can send messages as long as both links are **viable**
- Communication is impossible if either link **fails**

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### Even easier than rolling off a log...

- $V = V_1V_2$       •  $F = F_1 \cup F_2$
- $P(V) = P(V_1V_2) = P(V_1)P(V_2) = q_1q_2$   
**if the events  $V_1$  and  $V_2$  can be assumed to be independent**
- Similarly,  $P(F) = P(F_1) + P(F_2) - P(F_1F_2)$   
 $= p_1 + p_2 - p_1p_2$

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### Be sure to understand this too...

- $V = V_1V_2$       •  $F = F_1 \cup F_2$
- $P(V) = P(V_1V_2) = P(V_1)P(V_2) = q_1q_2$   
**if the events  $V_1$  and  $V_2$  can be assumed to be independent**
- $P(F) = P(F_1) + P(F_2) - P(F_1F_2)$   
 $= P(F_1) + P(F_2) - P(F_1)P(F_2)$   
 $= p_1 + p_2 - p_1p_2$   
 $= (1 - q_1) + (1 - q_2) - (1 - q_1)(1 - q_2)$   
 $= 1 - q_1q_2$

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### Matters did not improve

- $P(V) = P(V_1V_2) = P(V_1)P(V_2) = q_1q_2$  for independent failures
- $P(V) < \min\{q_1, q_2\}$
- $P(F) = p_1 + p_2 - p_1p_2 > \max\{p_1, p_2\}$
- If  $p_1 = p_2 = 10^{-3}$ , then  
 $P(F) = 10^{-3} + 10^{-3} - 10^{-6} = 2 \times 10^{-3}$

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### Generalization

- For  $n$  parallel links with failure probabilities  $p_1, p_2, \dots, p_n$ , the **probability of failure** (that is, not being able to communicate) is  
 $p_1p_2 \dots p_n$
- For  $n$  serial links with failure probabilities  $p_1, p_2, \dots, p_n$ , the **probability of failure** (that is, not being able to communicate) is  
 $1 - q_1q_2 \dots q_n$

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### More complicated networks

- In some cases, complicated networks can be simplified by replacing sets of serial (or parallel) links by an **equivalent link**

- $p_5 = p_1 p_2$     $p_6 = p_3 p_4$
- $p = p_5 + p_6 - p_5 p_6 = p_1 p_2 + p_3 p_4 - p_1 p_2 p_3 p_4$

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### Another example

- $q_5 = q_1 q_3$     $q_6 = q_2 q_4$     $q = q_5 + q_6 - q_5 q_6$

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### Even more complicated networks

- Simplification is not always possible

- In dealing with such problems, it is convenient to use **conditional probabilities** and the **theorem of total probability**

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### Assume Link #5 has not failed

- The network is equivalent to the following

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### When Link #5 has not failed...

- The network is equivalent to the following

- $P(F | \bar{V}_5) = p_1 p_2 + p_3 p_4 - p_1 p_2 p_3 p_4$
- Note that all this works because we are assuming that the failures (and non-failures) are independent

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### Assume Link #5 has failed

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### When Link #5 has failed...

- The network is equivalent to the following

- $P(V|F_5) = q_1q_3 + q_2q_4 - q_1q_2q_3q_4$
- $P(F|F_5) = 1 - (q_1q_3 + q_2q_4 - q_1q_2q_3q_4)$

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### Now put it all together ...

- $P(F) = P(F|V_5)P(V_5) + P(F|F_5)P(F_5)$   
 $= (p_1p_2 + p_3p_4 - p_1p_2p_3p_4)(1 - p_5)$   
 $+ (1 - (q_1q_3 + q_2q_4 - q_1q_2q_3q_4))p_5$
- Thus, we see that using conditional probabilities of failure together with the theorem of total probability is a useful trick

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### It is not just for breakfast anymore

- The same method could have been used in even simpler cases

- $P(F) = P(F|V_1)P(V_1) + P(F|F_1)P(F_1)$   
 $= 0 + p_2p_1$   
 $= p_1p_2$  just as before!

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### What is communication capacity?

- An important question in communication networks is how much traffic (e.g. calls) can be carried from one node to another
- The amount of traffic that can be carried is called the communication capacity
- Distinguish between the traffic that can be carried, and the traffic that is being carried
- The latter usually depends on conditions outside the network

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### More detailed answers, please!

- Failure analysis reveals only that “Yes, communication is possible” or “No, communication is not possible”
- Failure analysis crudely quantifies the answer to the question of communication capacity into two values: zero and nonzero
- We seek a more detailed description
- Because of link failures, the capacity can be considered to be a random variable

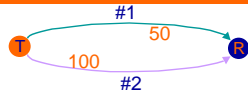
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### Link capacity vs. system capacity

- Link #1 has capacity 50 and Link #2 has capacity 100 (e.g. capacity is the number of telephone calls that can be carried)
- If both links are viable, system capacity is 150, but if links have failed, the system capacity can be 100, or 50, or 0

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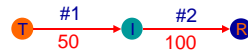
### System capacity = random variable



- Let  $X$  denote the system capacity
- $P\{X = 150\} = q_1q_2$
- $P\{X = 100\} = p_1q_2$
- $P\{X = 50\} = q_1p_2$
- $P\{X = 0\} = p_1p_2$

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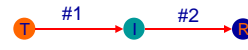
### Link capacity vs. system capacity



- Link #1 has capacity 50 and Link #2 has capacity 100 (e.g. capacity is the number of telephone calls that can be carried)
- If both links are viable, **system capacity** is  $50 = \min\{50, 100\}$
- If either (or both!) link has failed, system capacity is 0

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### System capacity = random variable

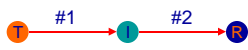


- Let  $X$  denote the system capacity
- $P\{X = 50\} = q_1q_2$
- $P\{X = 0\} = 1 - q_1q_2$

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### The network paradigm

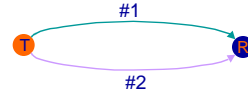
- Communication networks and their failures are a paradigm for the study of failures of systems in general
- If a system has two components and fails if either component fails, then an equivalent network model is



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### More on the network paradigm

- If a system has two components and fails only when both components fail, then an equivalent network model is



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### But, complications arise...

- In modeling some systems as a communication network, it is sometimes necessary to have a link appear in more than one place
- When the corresponding component fails, **all** these links fail

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### Example

- Consider the TMR system shown

- System fails when two or more circuits fail
- Majority gate does not fail

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### Network Model for TMR system

- The network model for the TMR system is as shown below

- Note that links are duplicated
- If link #3 fails, it disappears from 2 places

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### Network Model for TMR system

- Note that links are duplicated
- If link #3 fails, it disappears from 2 places
- But a path still exists between T and R because links #1 and #2 are viable

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### General TMR system

- In a general TMR system, the majority gates are also replicated
- This greatly improves the reliability

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### Network model: TMR system

- $P\{\text{two or more TMR outputs in error}\}$   
 $P\{2 \text{ or more circuits fail}\}$   
 $+ P\{2 \text{ or more majority gates fail}\}$   
 $P\{2 \text{ or more circuits fail}\}$   
 $= 3 \times 10^{-8} \text{ if } P\{\text{circuit fail}\} = 10^{-4}$

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### Summary

- We discussed communication networks and how the link failure probabilities can be used to compute the system failure probabilities
- We discussed how general reliability problems can be modeled in terms of communication network problems

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