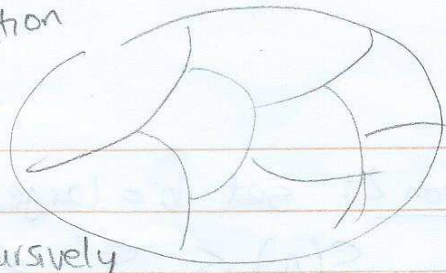


Build DS for R

Build DS inside each region recursively

after retriangulation



Query: find region of R containing  $q$   
recurse inside region

$$\left\{ \begin{array}{l} S(n) \leq \sum_i S(n_i) + S_0\left(\frac{cn}{\sqrt{b}}\right) + O(n) \\ \text{with } \sum n_i \leq n, \quad \max n_i \leq b \\ Q(n) \leq Q(b) + Q_0\left(\frac{cn}{\sqrt{b}}\right) + O(1) \end{array} \right.$$

Option 1: start with  $S_0(n) = O(n^2)$ ,  $Q_0(n) = O(\log n)$  (Method 0)

$$\text{Set } b = n^{1-\delta}: \quad S(n) \leq \sum_i S(n_i) + O\left(n^{\frac{1+\delta}{2}}\right)$$

$$\Rightarrow S(n) = O(n^{1+\delta})$$

$$Q(n) \leq Q(n^{1-\delta}) + O(\log n)$$

$$\Rightarrow Q(n) = O(\log n + (1-\delta)\log n + \dots) \\ = O(\log n)$$

bootstrap with  $S_0(n) = O(n^{1+\delta})$ ,  $Q_0(n) = O(\log n)$

$$\text{Set } b = \sqrt{n}: \quad S(n) \leq \sum_i S(n_i) + O(n)$$

$$\Rightarrow S(n) = O(n \log \log n)$$

$$Q(n) = O(\log n)$$

bootstrap again with  $S_0(n) = O(n \log \log n)$ ,  $Q_0(n) = O(\log n)$

$$\text{Set } b = (\log \log n)^2: \quad S(n) \leq \sum_i n_i + O(n)$$

& don't recurse

$$\Rightarrow S(n) = \boxed{O(n)}$$

$$Q(n) \leq Q(b + \log n)$$

$$\Rightarrow Q(n) = \boxed{O(\log n)}$$

Rmk: this gives general space reduction that preserves query time

↙

Option 2: set  $b = \text{large const}$ , recurse in  $R$  instead

$$S(n) \leq S\left(\frac{cn}{\sqrt{b}}\right) + O(n) \leq S(n/2) + O(n)$$

$$\Rightarrow O(n)$$

$$Q(n) \leq Q\left(\frac{cn}{\sqrt{b}}\right) + O(1) \leq Q(n/2) + O(1)$$

$$\Rightarrow O(\log n)$$

Method b: Kirkpatrick's Hierarchy ('83)

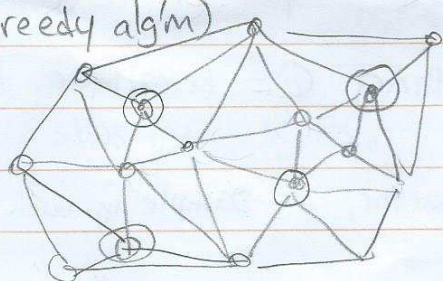
Simplification of option 2 without separators!

Fact In any planar graph with  $n$  vertices,

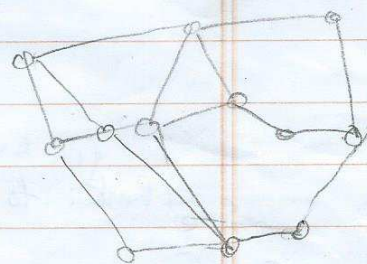
$\exists$  indep set  $I$  of size  $\geq \frac{n}{24}$ .

Furthermore, each vertex in  $I$  has  $\text{deg} \leq 11$ .

(PF: by greedy algm)



$\Rightarrow$



remove  $I \Rightarrow$  subdivision  $R$  with  $\leq \frac{23}{24}n$  vertices

each region has size  $\leq 11$ .

$$S(n) \leq S\left(\frac{23}{24}n\right) + O(n) \Rightarrow O(n)$$

$$Q(n) \leq Q\left(\frac{23}{24}n\right) + O(1) \Rightarrow O(\log n)$$

*Link: Dobkin-Kirkpatrick extends to 3D convex polyhedron*