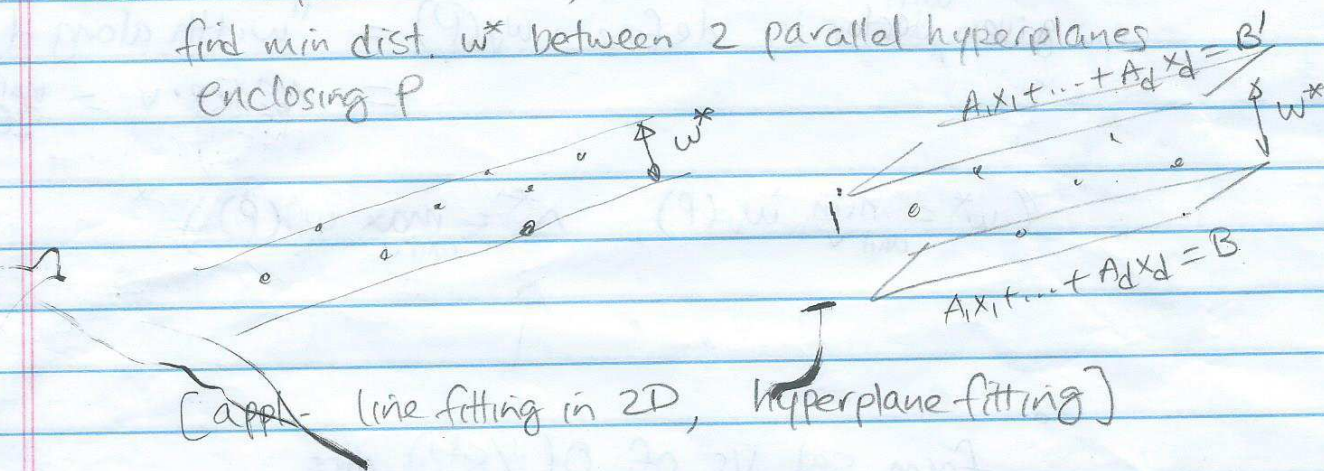


The Width Problem

Given n pts $P \subset \mathbb{R}^d$,

find min dist. w^* between 2 parallel hyperplanes enclosing P



[app. line fitting in 2D, hyperplane fitting]

Known exact algms:

$d=2 \quad O(n \log n)$

$d=3 \quad \tilde{O}(n^{3/2})$ (complicated!)

$d \geq 4 \quad O(n^{\lceil d/2 \rceil})$

(can make $B' = B + 1$)

$$w^* = \min \frac{B' - B}{\sqrt{A_1^2 + \dots + A_d^2}}$$

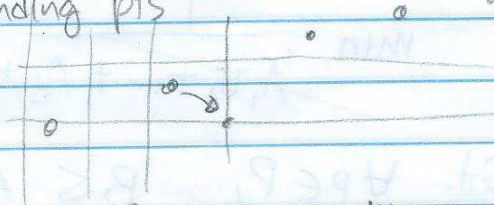
s.t. $\forall p = (p_1, \dots, p_d) \in P,$

$$B \leq A_1 p_1 + \dots + A_d p_d \leq B'$$

over vars A_1, \dots, A_d, B, B'

intersection of $O(n)$ halfspaces in $d+1$ vars
(a convex polytope)

ideas rounding pts



fail!

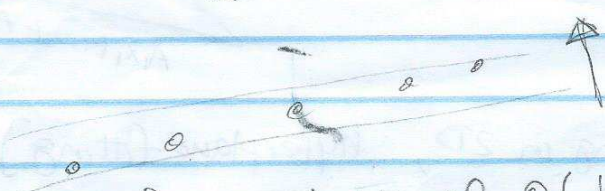
error $O(\epsilon \Delta^*)$, may be bigger than $O(\epsilon w^*)$

- rounding dirs

given ^{unit} vector v , define $w_v(P) =$ "width along dir v "
 $= \max_{p \in P} p \cdot v - \min_{q \in P} q \cdot v$

(also called "extent")

$$(w^* = \min_{\text{unit } v} w_v(P), \quad \Delta^* = \max_{\text{unit } v} w_v(P))$$



form set V_δ of $O(1/\delta^{d-1})$ dirs

try all $v \in V_\delta \Rightarrow$ error $O(\delta \Delta^*)$, not $O(\delta w^*)$
 fail!

Alg'm 1 (Duncan-Goodrich-Ramos '97)

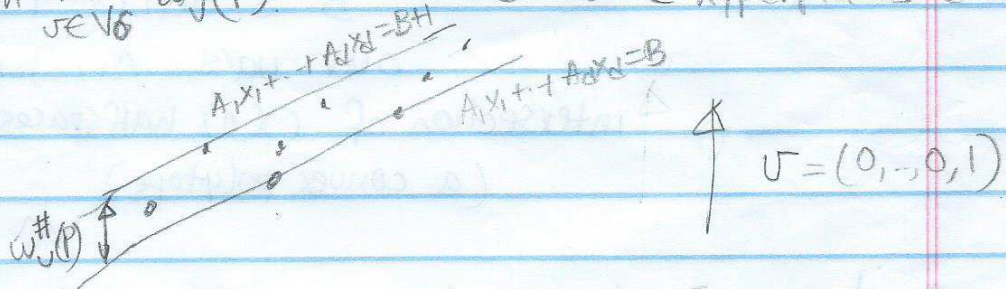
idea - round dirs, but more cleverly

for each $v \in V_\delta$

compute $w_v^\#(P) \equiv$ "skewed width along dir v "
 $=$ min distance measured along v
 between 2 hyperplanes enclos. P .

return $\min_{v \in V_\delta} w_v^\#(P)$

eg.



$$w_v^\#(P) = \min \frac{1}{A_1 v_1 + \dots + A_d v_d} \max_{p \in P} A_1 p_1 + \dots + A_d p_d$$

st. $\forall p \in P, B \leq A_1 p_1 + \dots + A_d p_d \leq B+1$.

this is linear programming in $d+1$ vars

$\Rightarrow O(n)$ time by known alg'ms

$$\text{Total time } O(|V_\delta|n) = O\left(\frac{1}{\delta^{d-1}}n\right) = \boxed{O\left(\frac{1}{\epsilon^{\frac{d-1}{2}}}\right)n}$$

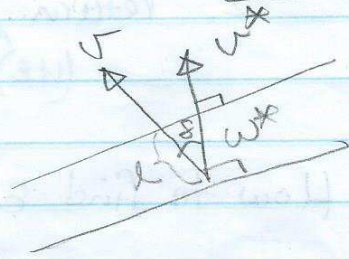
Approx factor:

let v^* be optimal dir

let $v \in V_\delta$ s.t. $\angle(v, v^*) \leq \delta$

$$\Rightarrow w_v^\#(P) \leq l \leq \frac{w^*}{\cos \delta}$$

$$\Rightarrow \text{factor } \frac{1}{\cos \delta} = 1 + \epsilon \text{ by setting } \delta = \theta(\sqrt{\epsilon})$$



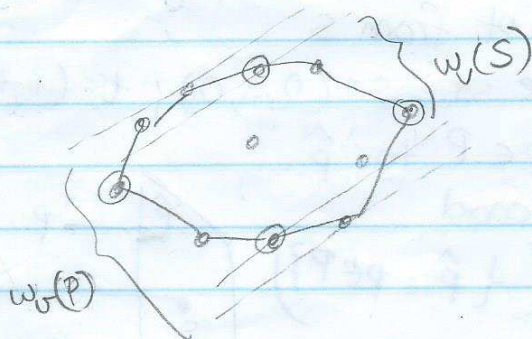
Alg'm 2: (Agarwal-Har-Peled-Urabkarajan '04)

solves more general problem

of approximating width along all dirs

Def An ϵ -kernel (also called ϵ -coreset for directional width) is a subset $S \subseteq P$ s.t.

$$\forall v \in \mathbb{R}^d, w_v(S) \geq (1-\epsilon) w_v(P)$$



$w_v(S) \geq (1-\epsilon) w_v(P)$
related to
"approx convex hull"



Main Thm \exists ϵ -kernel of size $O\left(\frac{1}{\epsilon^{(d-1)/2}}\right)$

\uparrow
[optimal]