# CS 598. Spectral Graph Theory Problem Set 2

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# Problem 1 Alon Boppana

(5 pts.) Show that there exists a constant c, such that for every finite connected d-regular graph G with diameter D, the second eigenvalue of the adjacency matrix of G satisfies the following bound:

$$\lambda_2(G) \ge 2\sqrt{d-1}\left(1 - \frac{c}{D^2}\right)$$

Please try to do this without looking at the Alon Boppana proof, following the style of proof that we talked about in Lecture 10.

# **Problem 2** Abelian Cayley Graphs

(5 pts.)

Let G be an Abelian group on n elements, and H a Cayley graph of G of degree d with second adjacency matrix eigenvalue less than  $(1-\epsilon)d$  for some constant  $\epsilon$ . Show that  $d \geq \Omega_{\epsilon}(\log n)$ . Namely, in order to construct expanders from abelian Cayley graphs, the degree needs to be logarithmic.

# **Problem 3** Diameter and Multiplicity

(7 pts.) Let G be any graph, and let L = D - A be the Laplacian, where D is the diagonal degree matrix and A is the adjacency matrix. You will prove that, for any  $r \ge 1$ , if L has at most r+1 distinct eigenvalues, then the diameter of G is at most r. In particular, this implies that the only graph with exactly two distinct eigenvalues is the complete graph. Let D be the diameter of G.

- 1. Prove that if p is a polynomial of degree m such that p(A) has all non-zero entries, then  $D \leq m$ .
- 2. Prove the same thing for L instead of A.

3. Now, let  $0 = \lambda_1 \leq \lambda_2, \dots, \leq \lambda_n$  be the eigenvalues of L with corresponding orthogonal eigenvectors  $u_1, u_2, \dots, u_n$ . Prove that for any polynomial p such that p(0) = 1, we have

$$p(L) = \frac{J}{n} + \sum_{i=1}^{n} p(\lambda_i) u_i u_i^T$$

where J is the all one's matrix.

4. Using the fact that L has only r + 1 distinct eigenvalues, use the preceding two parts to show that  $D \leq r$ .

#### **Problem 4** Generalized Cheeger Rounding

(8 pts.)

Let G be an n-node, undirected, connected graph (you can also assume it is unweighted if you want) with adjacency matrix A. Define the normalized K-cut of a graph, into  $K \ge 2$  disjoint sets of vertices  $S_1, S_2, \dots, S_K$ , to be

$$NCUT(S_1, S_2, \cdots, S_K) = \sum_{i=1}^K \frac{Cut(S_i)}{Vol(S_i)}$$

where  $Cut(S) = \sum_{i \in S, j \in V \setminus S} A(i, j)$  and  $Vol(S) = \sum_{i \in S} d_i$ . Here  $d_i$  is the degree of vertex *i*. Let  $h_i$  be an "indicator vector" for set  $S_i$  as follows:

$$h_i(v) = \begin{cases} \frac{1}{\sqrt{Vol(S_i)}} & \text{if } v \in S_i; \\ 0 & \text{o.w.} \end{cases}$$

Let H be an n by K matrix whose columns are the  $h_i$ , i.e.  $H = [h_1, \dots, h_K]$ .

- 1. Show that  $Ncut(S_1, \dots, S_K) = Tr(H^T L H)$ , where L is the Laplacian of G.
- 2. Show that  $H^T D H = I_{K \times K}$
- 3. Show a spectral relaxation of minimum K-cut, generalizing Cheeger's rounding.