

The role of compatibility in the diffusion of technologies through social networks

Immorlica et al.

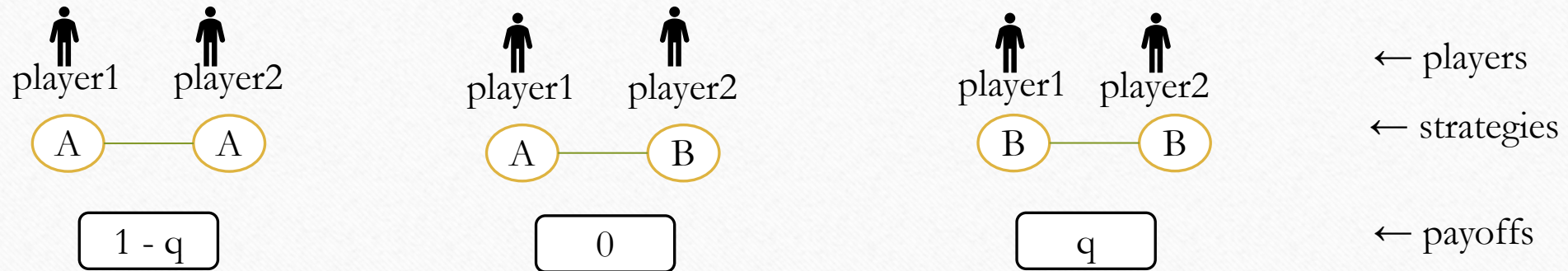
By:
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Background

- Diffusion in social network
 - Process in which new ideas, behaviors or practices diffuse through populations
 - Emergence of social norms
 - Adoption of new technologies

Background

- Coordination game in social networks

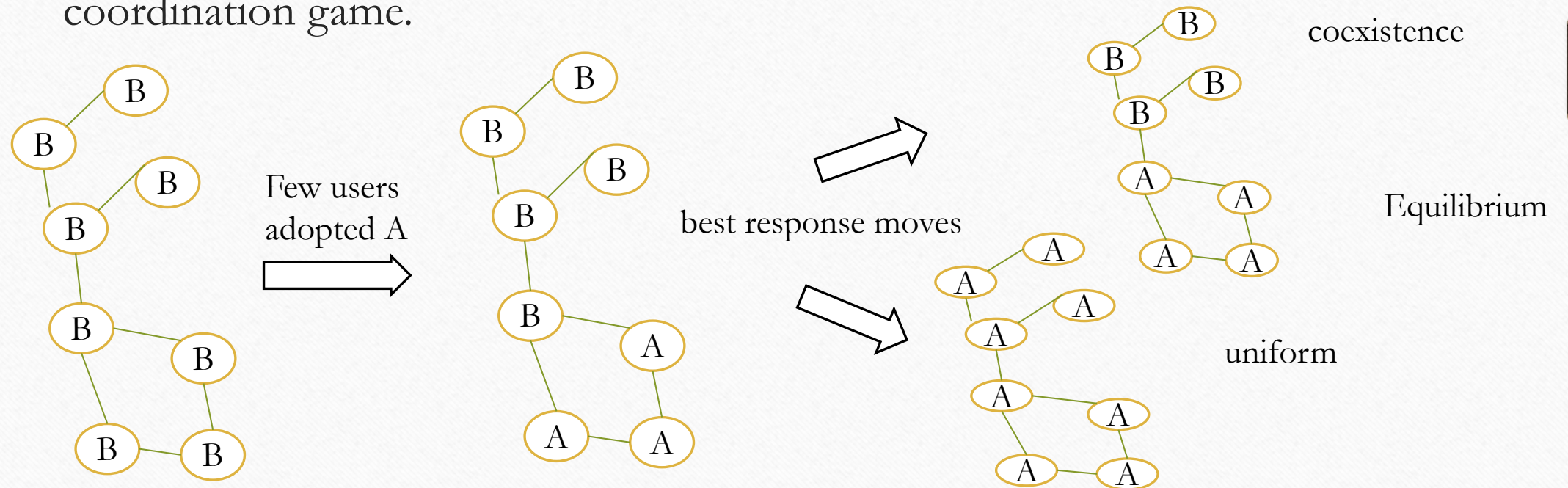


lack of interoperability

If $q < \frac{1}{2}$ means that A is better technology than B since A-A payoff $>$ B-B payoff

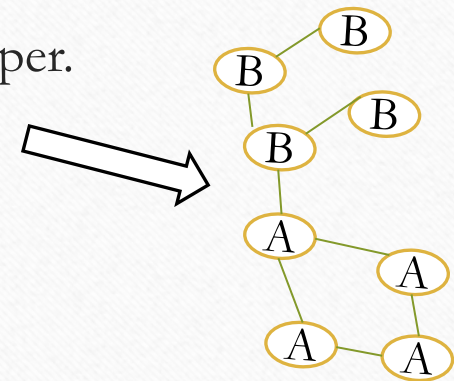
This work

- This work focuses on analyzing the dynamics of reaching equilibrium in the coordination game.

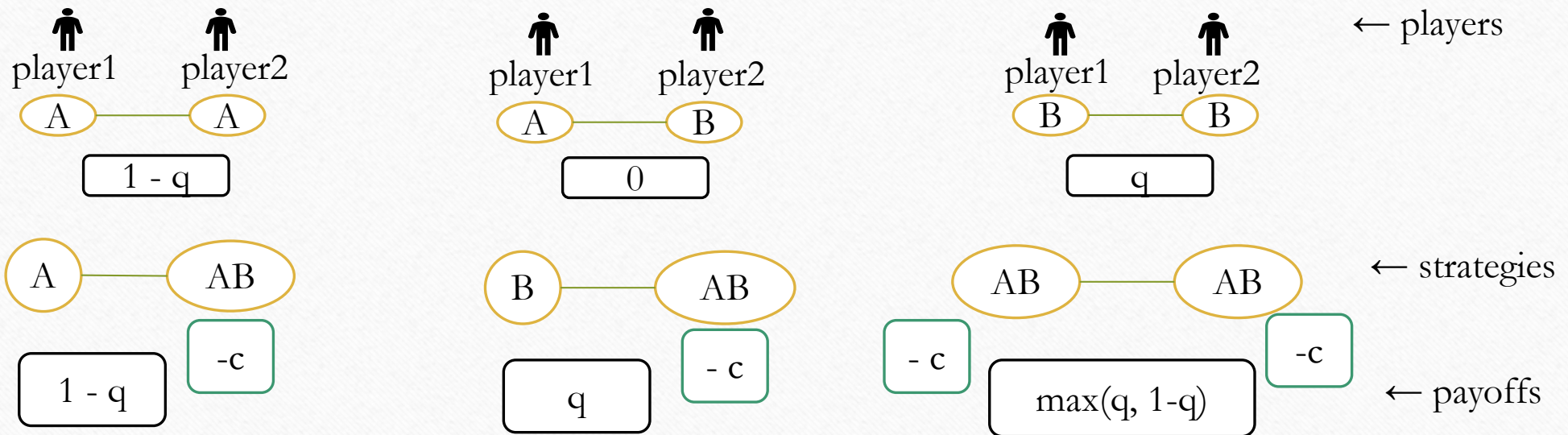


Key observation:

- coexistence is the typical outcome
 - detailed analysis of the coexistent boundary is the focus of the paper.
- individuals can become bilingual
 - people speak multiple languages
 - people have accounts on multiple online systems



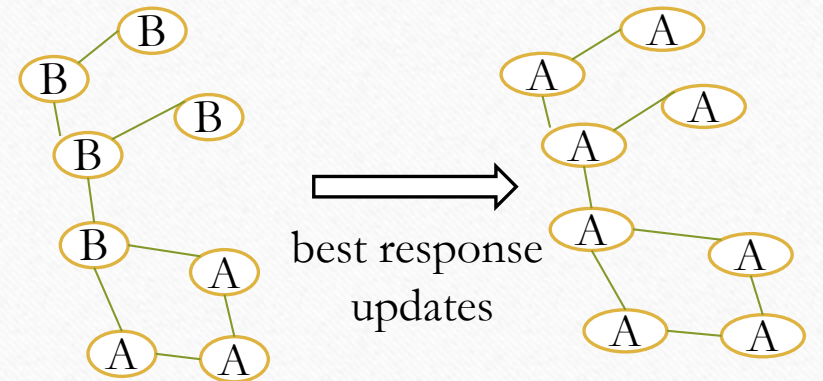
Game with bilingual behavior (AB strategy)



c : fixed cost penalty paid by adopter of AB
 two parameters: q, c

Definitions

- Assume underlying graph is infinite and each node has degree Δ
- A can become *epidemic*
- $r = c / \Delta$; penalty per edge cost

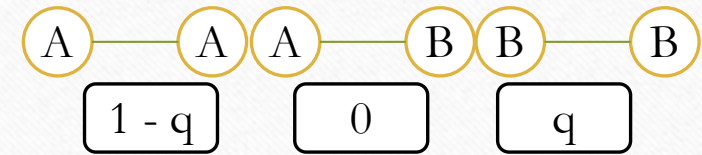


Start from
state in which
finite set S
adopt A

A has become
epidemic

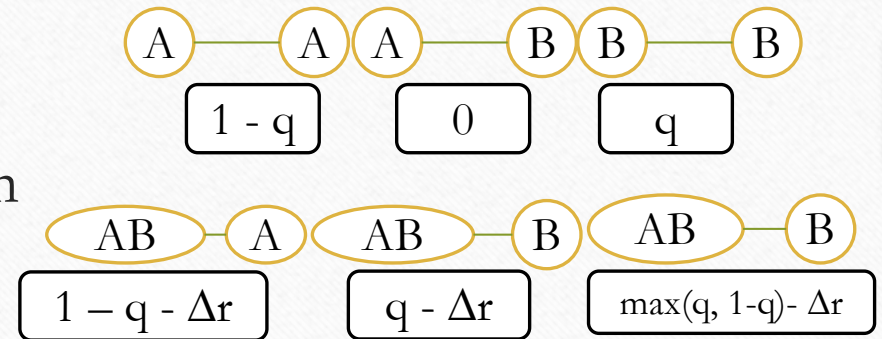
Contagion Threshold - non-bilingual model

- supremum of q for which A can become epidemic in G .
- $\sup q^*(G) = 1/2$
- Alternatively, for $q > 1/2$, there is no graph G in which A can become epidemic.
- If $q > 1/2$, the already existing B-B edges has better payoff and no individual will switch to A.



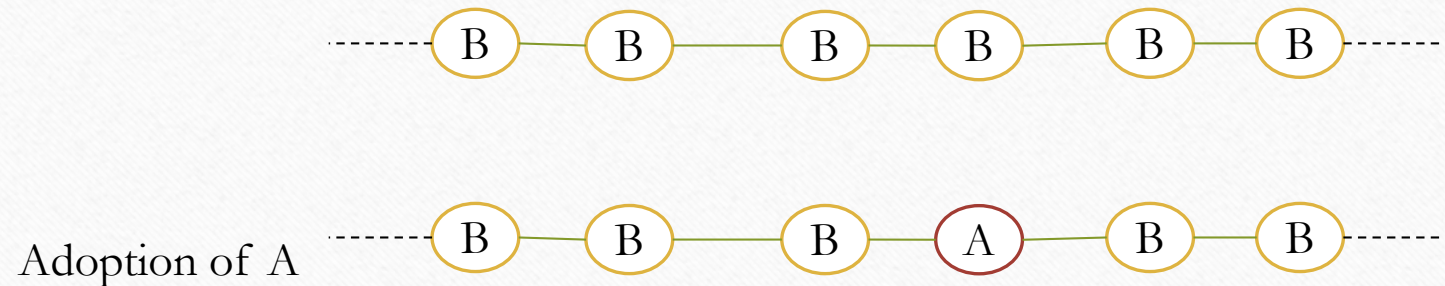
Contagion Threshold - for bilingual strategy

- There are two parameters in this game: q, r
- So, instead of a contagion threshold we have an epidemic region in a two dimensional space.



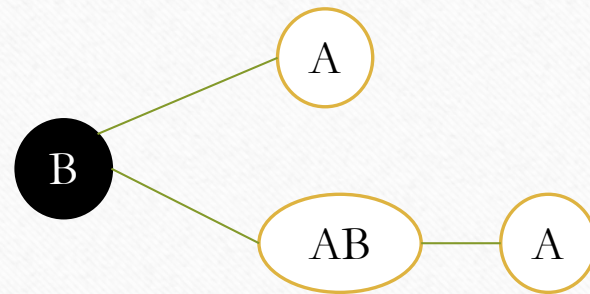
Epidemic Region for Infinite line

- Infinite Line:
 - $\Delta=2$; all nodes have degree 2

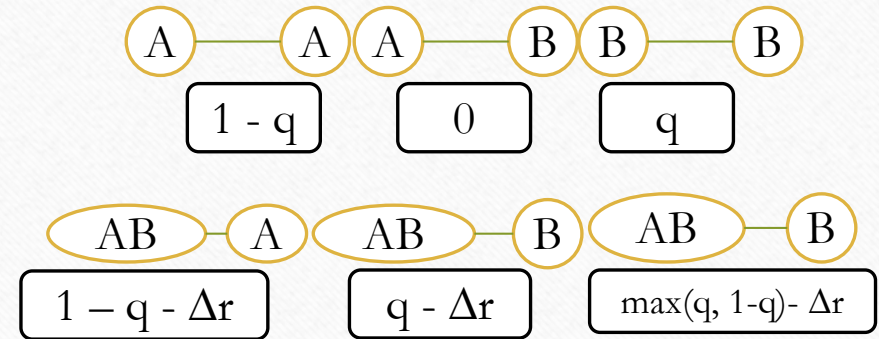


Epidemic Region for Infinite line

- Two ways for A to be epidemic:
- **Case1:** B nodes directly switch to A
- **Case2:** B nodes switch to AB then A



Epidemic Region for Infinite line

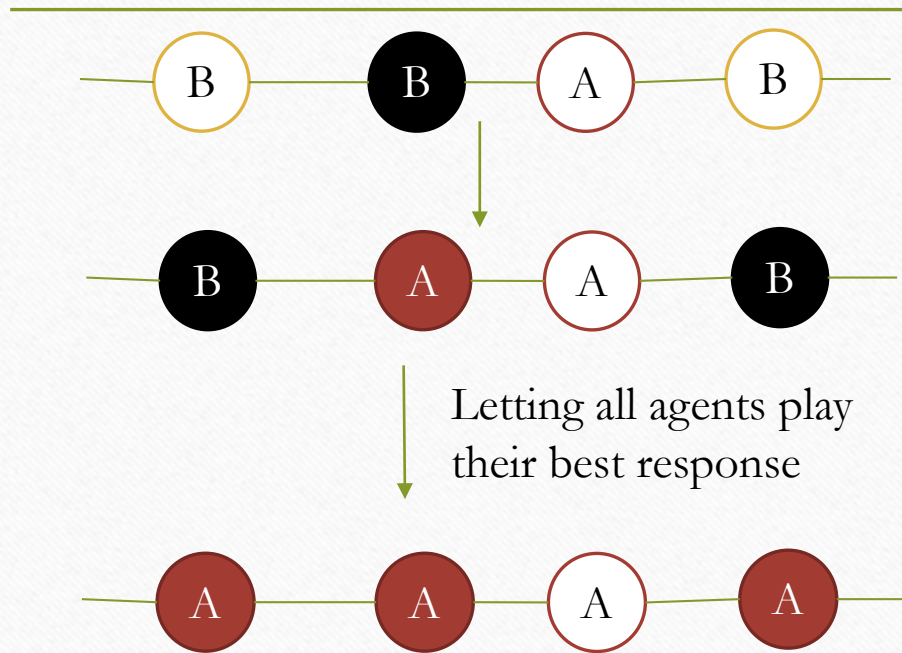


B switch to	Payoff before switch	Payoff after switch	Utility
A	B-B + B-A $q + 0$	B-A + A-A $0 + 1 - q$	$1 - q - q$
AB	B-B + B-A $q + 0$	B-AB + AB-A $q + 1 - q - 2r$	$1 - 2r - q$

Case 1: Direct A

- For B to switch to A,
 - Utility A should be positive:
 - $1 - 2q \geq 0$
 - $q \leq 1/2$
 - Utility A is greater than utility AB:
 - $1 - q - q \geq 1 - 2r - q$
 - $q \leq 2r$

Case 1: Direct A

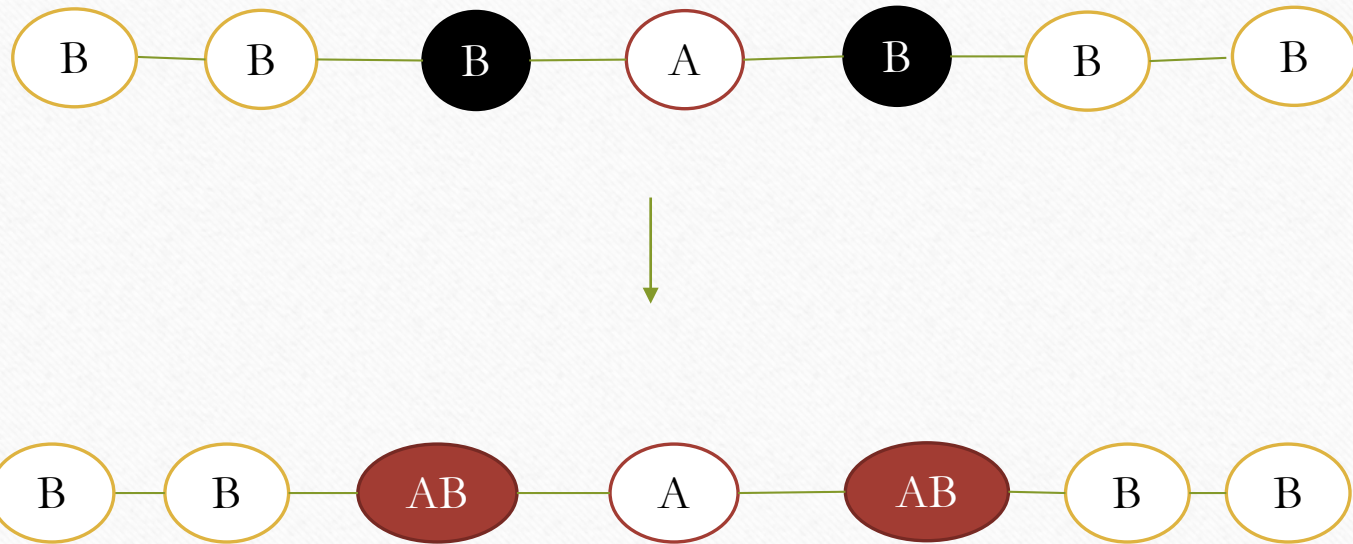


Conditions
 $q \leq 1/2$
 $q \leq 2r$

Case 2: Via AB

- For B switch to AB,
 - Utility AB should be positive:
 - $1 - 2r - q \geq 0$
 - $1 - q \geq 2r$
 - Utility AB is greater than utility A:
 - $1 - q - q \leq 1 - 4r - q$
 - $q \geq 2r$

Case 2: Via AB



Conditions
 $1 - q \geq 2r$
 $q \geq 2r$

Case 2: Via AB



B switch to	Payoff before switch	Payoff after switch	Utility
A	$AB-B + B-B$ $q + q$	$AB-A + A-B$ $1 - q$	$1 - q - 2q$
AB	$AB-B + B-B$ $q + q$	$AB-AB + AB-B$ $\max(q, 1-q) + q - 2r$	$\max(2q, 1) - 2r - 2q$

Case 2: Via AB

- For B to switch to AB.
 - Utility AB should be positive:
 - $\max(2q, 1) - 2r - 2q \geq 0$
 - $1 \geq 2r + 2q$
 - Utility AB should be greater than utility A:
 - $\max(2q, 1) - 2r - 2q \geq 1 - q - 2q$
 - $1 - 2r \geq 1 - q$
 - $q \geq 2r$

Conditions

$$1 - q \geq 2r$$

$$q \geq 2r$$

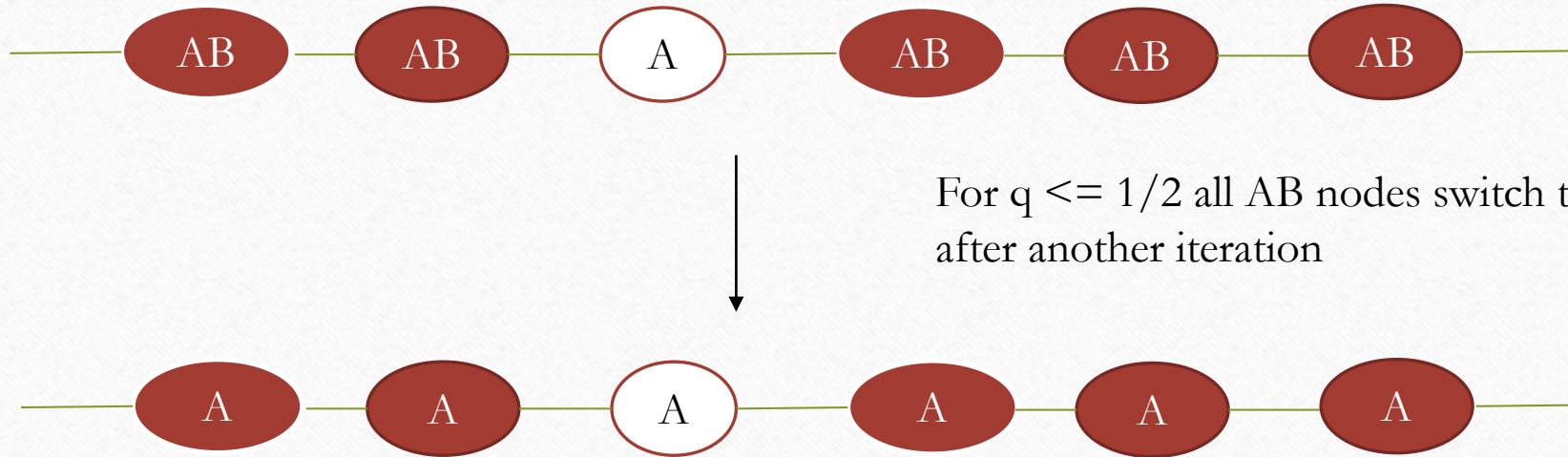
since $r > 0$ and for inequality to hold

$$\max(2q, 1) - 2q \geq 2r \geq 0$$

$$\max(2q, 1) \geq 2q$$

$$\max(2q, 1) = 1 \text{ and } q < 1/2$$

Case 2: Via AB



For $q \leq 1/2$ all AB nodes switch to A after another iteration

Epidemic region for infinite line

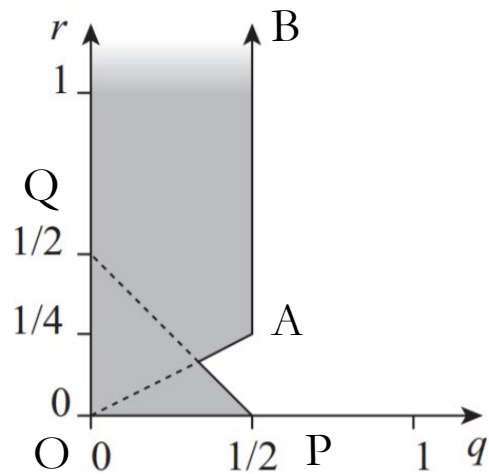


Figure 1: The region of the (q, r) plane for which technology A can become epidemic on the infinite line.

- Direct A (Region OAB)
 - $q \leq 1/2$
 - $q \leq 2r$
- Via AB (Region OPQ)
 - $1 \geq 2r + 2q$
 - $q \geq 2r$
 - $q \leq 1/2$

Epidemic Regions for other graphs

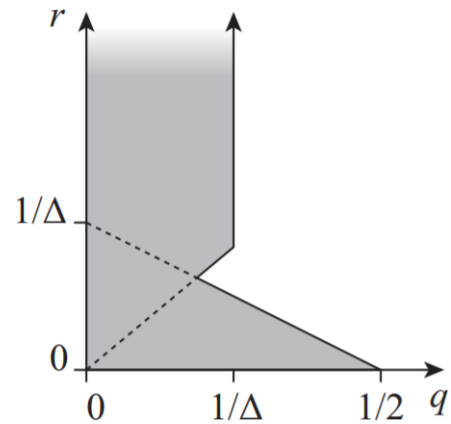


Figure 4: Epidemic regions for the infinite Δ -regular tree

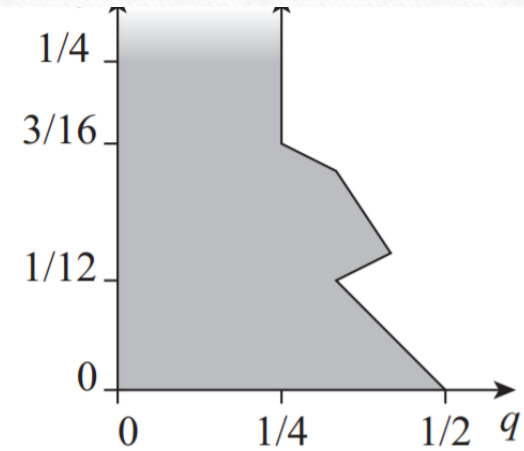


Figure 3: Epidemic regions for the infinite grid

Interpretation for epidemic region

- r is very small
 - Cheap to adopt AB, and on another update nodes switch to A (better)
- r is very large
 - Adoption cost of AB is very high, so nodes having A, B neighbors will switch to A
- r is intermediate
 - Allows a boundary of AB between adopters of A and B.

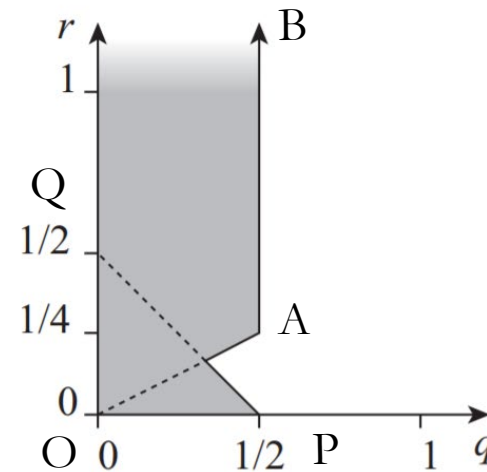


Figure 1: The region of the (q, r) plane for which technology A can become epidemic on the infinite line.

Interpretation for epidemic region

- The interesting insight is that worse technology can still survive if the cost of bilingual is in the middle interval.

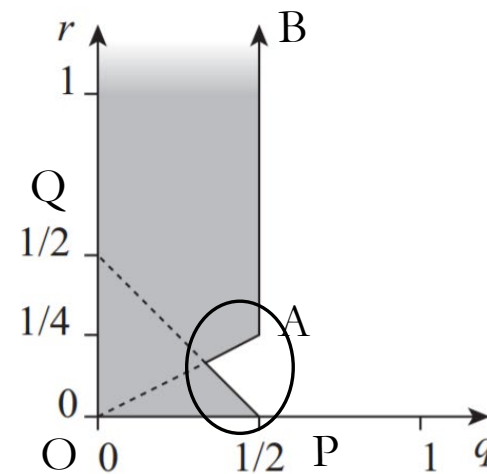


Figure 1: The region of the (q, r) plane for which technology A can become epidemic on the infinite line.

Characterization:

- Contagion games have well defined and stable equilibria.
 - If she adopts technology A she never discards it and once discards technology B, she never readapts it. Thus, after infinite best responses, each converge to a single strategy.
- Independent of the order of best responses.
 - Outcome is same for all schedules.

Blocking Structures

- A cannot become epidemic if (G, q, r) possess a certain blocking structure.
- The inequalities are linear in (q, r) so the epidemic region is the union of bounded or unbounded polygons.

DEFINITION 4.5. Consider a contagion game (G, q, r) . A pair (S_{AB}, S_B) of disjoint subsets of $V(G)$ is called a blocking structure for this game if for every vertex $v \in S_{AB}$,

$$\deg_{S_B}(v) > \frac{r}{q}\Delta,$$

and for every vertex $v \in S_B$,

$$(1 - q) \deg_{S_B}(v) + \min(q, 1 - q) \deg_{S_{AB}}(v) > (1 - q - r)\Delta,$$

and

$$\deg_{S_B}(v) + q \deg_{S_{AB}}(v) > (1 - q)\Delta,$$

where $\deg_S(v)$ denotes the number of neighbors of v in the set S .

Blocking Structures

- At $r \rightarrow \infty$, inequality 1 will fail
- At $r \rightarrow 0$, inequality 2 will fail
- At $q \rightarrow 1$, inequality 2 will fail
- At $q \rightarrow 0$, inequality 3 will fail
- So, for few extreme cases A will become epidemic

DEFINITION 4.5. Consider a contagion game (G, q, r) . A pair (S_{AB}, S_B) of disjoint subsets of $V(G)$ is called a blocking structure for this game if for every vertex $v \in S_{AB}$,

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and

$$\deg_{S_B}(v) + q \deg_{S_{AB}}(v) > (1 - q)\Delta,$$

where $\deg_S(v)$ denotes the number of neighbors of v in the set S .

Modeling Compatibility and Interoperability

- Interoperability
 - positive benefit (x) in A-B interactions ($x < q < 1 - q$)
- Extension to three technologies

Interoperability

- The q, r terms can be rescaled in terms of x and on rescaling by addition or division the behavior of the game remains unaffected.
- x can only strict the blocking structure; if A is epidemic for G ; then A is epidemic for G' .

Three technologies-Compatibility

- Suppose B and C are incumbent technologies currently in equilibrium
- If a better technology A appears, can they save themselves?
 - If B and C increases their compatibility by a calculated amount, they can resist an epidemic of A
- There are cases when compatibility is harmful to both parties.

Are the following assumptions realistic?

- All nodes have same degree but real world networks follow power law distribution.
- Every neighbor has the same influence in this model but in real-world, relationship strength might be a factor.