CS 573: Algorithms, Fall 2012 Homework 0, due Monday, September 3, 23:59:59, 2012

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This homework tests your S 373—many of these per too help you identify gape your own. Chapters 1-sexts as well.	oroblems have appos in your knowle	peared on homeworks of dge. You are respo	or exams in those class on sible for filling to	sses—primarily hose gaps on
Before you do anything page (http://courses. This web page gives instagether in order, write everything, use good En	engr.illinois.estructions on how your name and	edu/cs573/fa2012/fa to write and submit netID on every page,	aq/), and then check homeworks—staple	the box below. your solutions
I have	read the CS 5	573 Homework In	structions and I	FAQ.
Remember to do the quiz on moodle!				

"Be that as it may, it is to night school that I owe what education I possess; I am the first to own that it doesn't amount to much, though there is something rather grandiose about the gaps in it." – The tin drum, Gunter Grass

Required Problems

- 1. Moving numbers. [50 Points]
 - (A) [25 Points] The input is a multiset X of n positive integer numbers in the range 1 to k. Consider the famous algorithm:

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\begin{aligned} \mathbf{Play}(X): \\ \mathbf{while} \ X \ \text{contains more than one element } \mathbf{do} \\ \mathbf{if} \ X \ \text{contains the number 0 } \mathbf{then} \\ \text{Remove the number 0 from } X \\ \mathbf{continue} \\ \text{Two distinct elements } x_1 \ \text{and } x_2 \ \text{are chosen arbitrarily from } X \\ y_1 &= \min(x_1, x_2) - 1 \\ y_2 &= \max(x_1, x_2) + 1 \\ X \leftarrow (X \setminus \{x_1, x_2\}) \cup \{y_1, y_2\} \end{aligned}
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Here is an example of the execution of $Play(\{1, 2, 3, 4\})$.

$$\{1,2,3,4\} \implies \{1,2,2,5\} \implies \{0,2,3,5\} \implies \{2,3,5\} \implies \{2,2,6\}$$

$$\implies \{3,1,6\} \implies \{1,2,7\} \implies \{1,8,1\} \implies \{0,2,8\}$$

$$\implies \{2,8\} \implies \{1,9\} \implies \{0,10\} \implies \{10\} .$$

Prove (maybe using induction, but you do not have to) that **Play** always terminates. (Hint: Come up with an argument why in each step some non-trivial progress is being made.)

- (B) [25 Points] (Harder.) Prove that the algorithm always terminates after $O((kn)^2)$ steps. **Extra fun questions**: Since this problem is so much "fun", here are a few more questions. Do not submit a solution for this part (we also will not provide solutions to this part).
 - (I) (Hard.) Prove that, in the worst case, $\Theta((kn)^{3/2})$ steps are needed (that is, improve the upper bound, and provide a matching lower bound) if $k \leq n$. For k > n, prove that the right bound is $\Theta(kn^2)$.
 - (II) Show that if instead of increasing the maximum number by one, we increase it by two, then the number of steps in the worst case can be exponential.
- (III) (Easy.) Consider the modified algorithm, where the algorithm repeatedly takes a subset of numbers (out of the current set of numbers), decrease the smallest number in the subset by 1, and change the value of all the other numbers in an arbitrary fashion (again, we remove a number when it becomes zero). Prove that this algorithm always terminates.

2. Rolling to victory. [30 Points]

Alice and Bob each have a fair n-sided die. Alice rolls her die once. Bob then repeatedly throws his die until he rolls a number at least as big as the number Alice rolled. Each time Bob rolls, he pays Alice \$1. (For example, if Alice rolls a 5, and Bob rolls a 4, then a 3, then a 1, then a 5, the game ends and Alice gets \$4. If Alice rolls a 1, then no matter what Bob rolls, the game will end immediately, and Alice will get \$1.)

Exactly how much money does Alice expect to win at this game? Prove that your answer is correct. If you have to appeal to 'intuition' or 'common sense', your answer is probably wrong!

3. RANDOM WALK. [20 Points]

A **random walk** is a walk on a graph G, generated by starting from a vertex $v_0 = v \in V(G)$, and in the *i*th stage, for i > 0, randomly selecting one of the neighbors of v_{i-1} and setting v_i to be this vertex. A walk v_0, v_1, \ldots, v_m is of length m.

- (A) For a vertex $u \in V(G)$, let $P_u(m, v)$ be the probability that a random walk of length m, starting from u, visits v (i.e., $v_i = v$ for some i). Prove that a graph G with n vertices is connected, if and only if, for any two vertices $u, v \in V(G)$, we have $P_u(n-1, v) > 0$.
- (B) Prove that a graph G with n vertices is connected if and only if for any pair of vertices $u, v \in V(G)$, we have $\lim_{m\to\infty} P_u(m, v) = 1$.