

# Homework 0

Algorithms for Big Data  
CS498ABD Spring 2019

**Exercise 1: Balls and bins** Consider the standard balls and bins process. A collection of  $m$  identical balls are thrown into  $n$  bins. Each ball is thrown independently into a bin chosen uniformly at random.

- (a) What is the (precise) probability that a particular bin  $i$  contains exactly  $k$  balls at the end of the experiment.
- (b) Let  $X$  be the number of bins that contain exactly  $k$  balls. What is the expected value of  $X$ ?
- (c) What is the variance of  $X$ ?

**Exercise 2: Randomized max cut.** In the **max cut** problem, the input is a graph  $G = (V, E)$  with  $m = |E|$  edges and  $n = |V|$  vertices, and the goal is to partition  $V$  into two sets  $(A, B)$  (where  $B = V \setminus A$ ) maximizing the number of edges  $\{e = (u, v) \in E : u \in A, v \in B\}$  with endpoints in different sets. (Such an edge is said to be **cut** by the partition  $(A, B)$ ). This problem is known to be NP-Hard, but we will show that it is very easy to get a constant factor approximation.

- (a) Consider the following randomized algorithm.

```
random-partition( $G = (V, E)$ )
1.  $A, B \leftarrow \emptyset$ 
2. for each  $v \in V$ 
   A. with probability  $1/2$ 
     i.  $A \leftarrow A \cup \{v\}$ 
   B. else  $B \leftarrow B \cup \{v\}$ 
3. return  $(A, B)$ 
```

`random-partition` randomly partitions the vertices by assigning each vertex to  $A$  or  $B$  independently with equal probability. Show that this algorithm cuts  $m/2$  edges in expectation.

- (b) Let  $k \in \mathbb{N}$ . In the **max  $k$ -cut** problem, we want to partition  $V$  into  $k$  sets  $(A_1, \dots, A_k)$  maximizing the number of edges with endpoints in different parts. Consider the following randomized algorithm.

```
random- $k$ -partition( $G = (V, E)$ )
1.  $A_1, \dots, A_k \leftarrow \emptyset$ 
2. for each  $v \in V$ 
   //  $[k] = \{1, \dots, k\}$ 
   A. sample  $i \in [k]$  uniformly at random
   B.  $A_i \leftarrow A_i \cup \{v\}$ 
3. return  $(A_1, \dots, A_k)$ 
```

`random- $k$ -partition` randomly partitions the vertices into  $k$  sets analogously to `random-partition`. Show that this algorithm cuts  $(1 - 1/k)m$  edges in expectation.

**Exercise 3: Coupon Collectors.** In the coupon collectors problem, there are  $n$  coupons, and each round we are given one of the coupons uniformly at random. Coupons can repeat. We want to collect all  $n$  coupons, and in particular, we want to analyze the expected number of rounds before collecting all  $n$  coupons.

1. Suppose a coin flips heads with probability  $p$ . Show that the expected number of coin tosses until flipping heads is  $1/p$ .
2. For  $i \in [n]$ , show that the expected number of iterations between collecting the  $(i-1)$ th coupon and the  $i$ th coupon is  $\frac{n}{n+1-i}$ .
3. Show that the expected number of iterations until collecting all  $n$  coupons is  $nH_n$ , where  $H_n = 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n}$  is the  **$n$ th harmonic number** (and approximately  $\ln(n)$ ).