

CS 473 ✦ Spring 2020

🌀 Homework 10 🌀

Due Wednesday, April 22, 2020 at 9pm

This is the last homework assignment.

- Give a linear-programming formulation of the **bipartite maximum matching** problem. The input is a bipartite graph $G = (U \cup V; E)$, where $E \subseteq U \times V$; the output is the largest matching in G . Your linear program should have one variable for each edge. (Don't worry about the optimal solution being integral; it will be.)
 - Now derive the dual of your linear program from part (a). What do the dual variables represent? What does the objective function represent? What problem is this!?
- Given points $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)$ in the plane, the **linear regression problem** asks for real numbers a and b such that the line $y = ax + b$ fits the points as closely as possible, according to some criterion. The most common fit criterion is minimizing the **L_2 error**, defined as follows:¹

$$\varepsilon_2(a, b) = \sum_{i=1}^n (y_i - ax_i - b)^2.$$

But there are many other ways of measuring a line's fit to a set of points, some of which can be optimized via linear programming.

- The **L_1 error** (or *total absolute deviation*) of the line $y = ax + b$ is defined as follows:

$$\varepsilon_1(a, b) = \sum_{i=1}^n |y_i - ax_i - b|.$$

Describe a linear program whose solution (a, b) describes the line with minimum L_1 error.

- The **L_∞ error** (or *maximum absolute deviation*) of the line $y = ax + b$ is defined as follows:

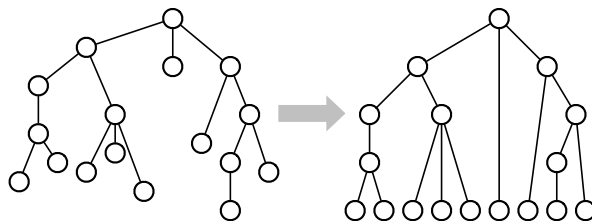
$$\varepsilon_\infty(a, b) = \max_{i=1}^n |y_i - ax_i - b|.$$

Describe a linear program whose solution (a, b) describes the line with minimum L_∞ error.

¹This measure is also known as *sum of squared residuals*, and the algorithm to compute the best fit is normally called (*ordinary/linear*) *least squares*.

3. Suppose you are given a rooted tree T , where every edge e has two associated values: a non-negative *length* $\ell(e)$ and a *cost* $\$(e)$ (which could be positive, negative, or zero). Your goal is to add a non-negative *stretch* $s(e) \geq 0$ to the length of every edge e in T , subject to the following conditions:

- Every root-to-leaf path π in T has the same total stretched length $\sum_{e \in \pi} (\ell(e) + s(e))$
- The total *weighted stretch* $\sum_e s(e) \cdot \$(e)$ is as small as possible.



- Describe an instance of this problem with no optimal solution.
- Give a concise linear programming formulation of this problem.
- Suppose that for the given tree T and the given lengths and costs, the optimal solution to this problem is unique. Prove that in the optimal solution, $s(e) = 0$ for every edge on some longest root-to-leaf path in T . In other words, prove that the optimally stretched tree has the same depth as the input tree. [Hint: What is a basis in your linear program? When is a basis feasible?]
- Describe and analyze an algorithm that solves this problem in $O(n)$ time. Your algorithm should either compute the minimum total weighted stretch, or report correctly that the total weighted stretch can be made arbitrarily negative.