## CS 473: Fundamental Algorithms, Spring 2011

## Discussion 4

February 8, 2011

## 1. Recurrences

Solve the following recurrences.

- $T(n)=5 T(n / 4)+n$ and $T(n)=1$ for $1 \leq n<4$.
- $T(n)=2 T(n / 2)+n l o g n$
- $T(n)=2 T(n / 2)+3 T(n / 3)+n^{2}$

2. Tree Traversal.

Let $T$ be a rooted binary tree on $n$ nodes. The nodes have unique labels from 1 to $n$.

- Given the preorder and postorder node sequences for $T$, give a recursive algorithm to reconstruct a tree that satisfies the preorder and postorder sequences. Is this reconstruction unique?
- Given the preorder and inorder node sequences for $T$, give a recursive algorithm to reconstruct a tree that satisfies the preorder and inorder sequences. Is this reconstruction unique?


## 3. Divide and Conquer.

Let $p=(x, y)$ and $p^{\prime}=\left(x^{\prime}, y^{\prime}\right)$ be two points in the Euclidean plane given by their coordinates. We say that $p$ dominates $p^{\prime}$ if and only if $x>x^{\prime}$ and $y>y^{\prime}$. Given a set of $n$ points $P=\left\{p_{1}, \ldots, p_{n}\right\}$, a point $p_{i} \in P$ is undominated in $P$ if there is no


Figure 1: The undominated points are shown as unfilled circles.
other point $p_{j} \in P$ such that $p_{j}$ dominates $p_{i}$. Describe an algorithm that given $P$ outputs all the undominated points in $P$; see figure. Your algorithm should run in time asymptotically faster than $O\left(n^{2}\right)$
4. Convex hull.

You are given a set $P$ of $n$ points in the plane, and you would like to compute their convex-hull (i.e., that is the shortest perimeter polygon that contains all the points). To see how the convex-hull looks like, think about the plane as being a wood board, and place a nail at each point. Now, you shrink a rubber band around the points. The rubber shrinks into the convex-hull. Clearly, the vertices of the convex-hull are a subset of the input points. Show an $O(n \log n)$ time algorithm for computing the convex-hull. (Hint: Split the plane by a vertical line, compute the convex-hulls on both sides, and then figure out how to stitch the two convex-hulls together. To get a handle on this stitching problem, think about rotating a tangent around one of the convex-hulls.)

