

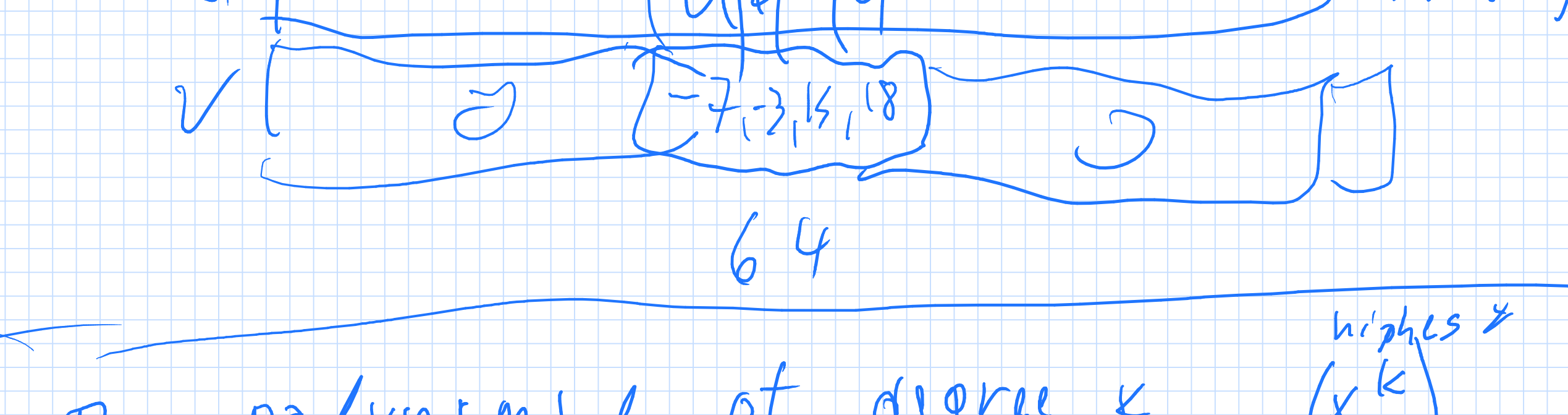
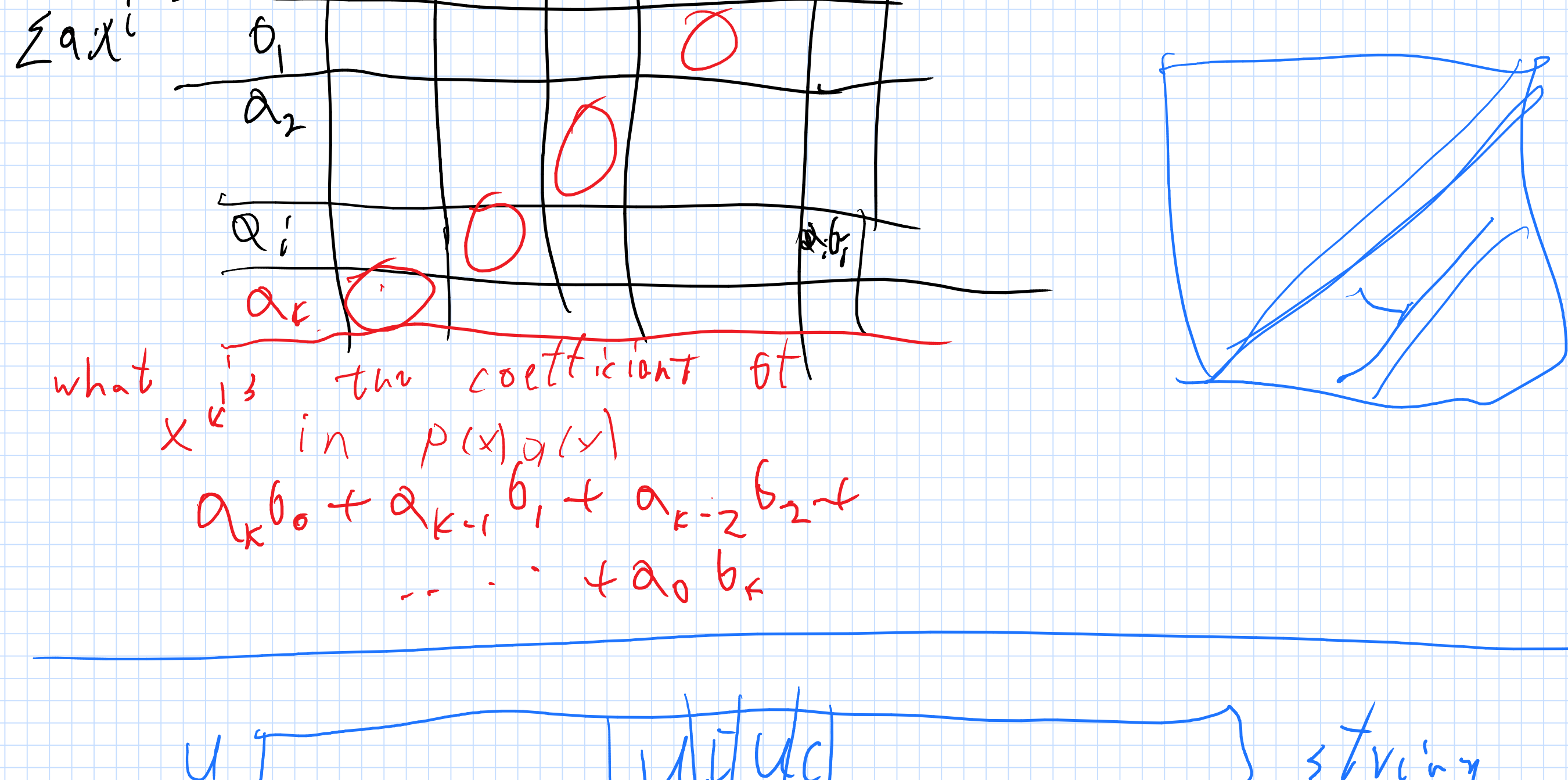
FFT Fast Fourier transform

$P(x) = a_0 + a_1x + a_2x^2 + \dots + a_{n-1}x^{n-1}$
 $q(x) = b_0 + b_1x + b_2x^2 + \dots + b_{n-1}x^{n-1}$

$p(x)q(x) = \sum_{i=0}^{n-1} \sum_{j=0}^{n-1} a_i b_j x^{i+j}$ n^2 terms
 $O(n^2)$
 pq degree $2n-2$

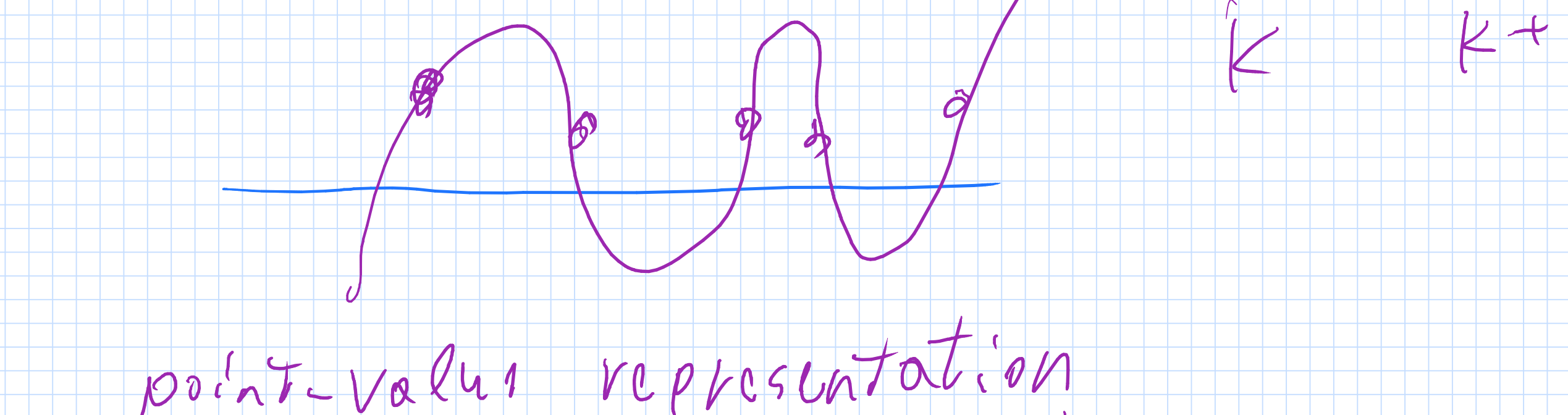
convolution algorithm

$u, v \in \mathbb{R}^n$
 $\langle u, v \rangle = \sum u_i v_i$
 $U \oplus i$ = shift U by i locations
 u $\begin{bmatrix} u_1 & u_2 & \dots & u_n \end{bmatrix}$
 $u \oplus i = \begin{bmatrix} 0 & \dots & 0 & u_1 & u_2 & \dots & u_n \end{bmatrix}$
 $\langle U \oplus i, V \rangle \quad i=1, \dots, n$



p polynomial of degree k (x^k)

If it is zero on $k+2$ different values $\Rightarrow p$ is the zero polynomial.



point-value representation $\{(x_1, y_1), (x_2, y_2), \dots, (x_{k+1}, y_{k+1})\}$ $k+1$ points

\Rightarrow define unique polynomial of degree k .

$g(x) = \sum_{i=1}^{k+1} y_i \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)}$ $g(x)$ polynomial of degree k
 $g(x_i) = y_i \quad \forall i$

$P(x) = \sum a_i x^i$
 $(x_i, P(x_i)) \quad i=1, \dots, n+1$

$(x_i, p(x_i))$
 $(x_i, q(x_i))$
 $p, q \equiv$ point-value representation $(x_i, p(x_i), q(x_i)) \quad \forall i$
 $O(n)$ time

$p(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + \dots + a_{n-1}x^{n-1}$
 $s(x) = \sum_{i=0}^{n/2} a_{2i} x^{2i}$
 $t(x) = \sum_{i=0}^{n/2} a_{2i+1} x^{2i}$

$p(x) = s(x) + x t(x)$
 $s(x) = \sum_{i=0}^{n/2} a_{2i} x^i$
 $t(x) = \sum_{i=0}^{n/2} a_{2i+1} x^i$

$s(x) = \tilde{s}(x^2)$
 $t(x) = \tilde{t}(x^2)$
 \tilde{s}, \tilde{t} at degree $n/2$

$p(x) = \tilde{s}(x^2) + x \tilde{t}(x^2)$ $X = 'n'$ values
 $O(n)$

$X \Rightarrow X^2 = \{x^2 \mid x \in X\}$
 $\tilde{s}(X^2) \quad \tilde{t}(X^2)$
 $|X^2| = n$

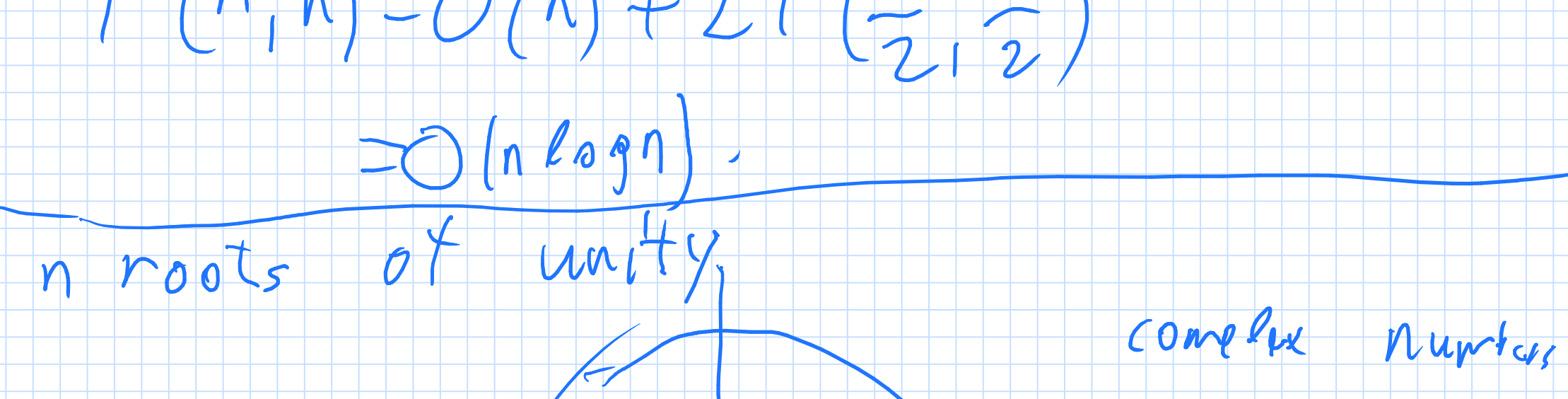
$T(n, n) = O(n) + 2T(\frac{n}{2}, \frac{n}{2})$
 $= O(n^2)$

$|X| = n \quad |X^2| = \frac{n}{2} \quad \left(\frac{-x^2}{-x} \right) = \frac{1}{2}$
 $X = \{-\frac{n}{2}, \frac{n}{2}\}$
 $(-a)^2 = a^2$

$|X| = n = 2^r$
 $X^2 = \{x^2 \mid x \in X\}$
 $|X^2| = \frac{n}{2}$ collapsible set

$T(n, n) = O(n) + 2T(\frac{n}{2}, \frac{n}{2})$
 $= O(n \log n)$

n roots of unity

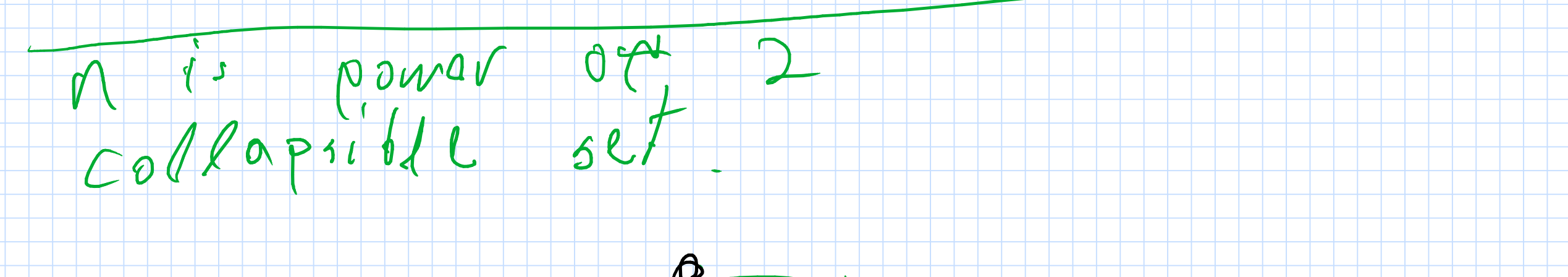


n th root of unity $\gamma_i = \cos \frac{2\pi i}{n} + i \sin \frac{2\pi i}{n}$

$\gamma_0, \dots, \gamma_{n-1} = (\gamma_1)^j \quad j=0, \dots, n-1$

$\gamma_j = \cos \frac{2\pi j}{n} + i \sin \frac{2\pi j}{n}$

n is power of 2 collapsible set.



$\forall i$ (γ_i, y_i) recover standard representation of polynomial

$y_i = p(\gamma_i) = \sum_{t=0}^{n-1} a_t \gamma_i^t \quad \forall i$

$\begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix} = \begin{pmatrix} 1 & \gamma_1 & \gamma_1^2 & \dots & \gamma_1^{n-1} \\ 1 & \gamma_2 & \gamma_2^2 & \dots & \gamma_2^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \gamma_n & \gamma_n^2 & \dots & \gamma_n^{n-1} \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{pmatrix}$

V Vandermonde matrix

$\vec{y} = V \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{pmatrix} \equiv \text{FFT}$

$\vec{y} = V \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{pmatrix}$

$V^{-1} \vec{y} = V^{-1} V \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{pmatrix} = \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{pmatrix}$

$V^{-1} = \frac{1}{n} \begin{pmatrix} 1 & \beta_1 & \beta_1^2 & \dots & \beta_1^{n-1} \\ 1 & \beta_2 & \beta_2^2 & \dots & \beta_2^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \beta_n & \beta_n^2 & \dots & \beta_n^{n-1} \end{pmatrix} \quad \beta_i = \frac{1}{\gamma_i} = \gamma_{-i}$

