

CS 473: Algorithms, Fall 2009

HBS 7

1. Recall the problem of finding a set of undominated points. Given an input set of points $P = \{p_1, p_2, \dots, p_n\}$ in the plane (with distinct coordinates), a point p is undominated if there are no other points above and to the right of p . Describe a method to maintain the set of undominated points as new points are added to the set. Specifically, describe and analyze a data structure that stores the undominated points of P , and an algorithm $\text{INSERT}(x, y)$ that adds the point $p = (x, y)$ to P and returns `TRUE` or `FALSE` to indicate whether the set of undominated points has changed. Your data structure should use $O(n)$ space, and your INSERT algorithm should run in $O(\log n)$ amortized time.
2. Given a string of letters $Y = y_1y_2 \dots y_n$, a segmentation of Y is a partition of its letters into contiguous blocks of letters (also called words). Each word has a quality that can be computed by a given oracle (*e.g.* you can call $\text{QUALITY}(\text{"meet"})$ to get the quality of the word "meet"). The quality of a segmentation is equal to the sum over the qualities of its words. Each call to the oracle takes linear time in terms of the argument; that is $\text{QUALITY}(S)$ takes $O(|S|)$ time.

Using the given oracle, give an algorithm that takes a string Y and computes a segmentation of maximum total quality.

3. Consider two distinct horizontal lines l_1 and l_2 in the plane. There are n points on l_1 with distinct x -coordinates $A = a_1, a_2, \dots, a_n$ and there are n points on l_2 with distinct x -coordinates $B = b_1, b_2, \dots, b_n$. Design an algorithm to compute, given A and B , a largest set S of non-intersecting line segments subject to the following restrictions:
 - (a) Any segment in S connects a_i to b_i for some i ($1 \leq i \leq n$).
 - (b) Any two segments in S do not intersect.