

Secure System Development Mechanisms

CS460 Cyber Security Lab
Spring 2010

Reading Material

- Web sites
 - Microsoft links from last lecture
 - Linux Capabilities - “man 7 capabilities” or <http://www.linuxjournal.com/article/5737>
- Papers
 - “The Security Architecture of qmail”, Hafiz, Johnson, and Afandi. PLoP, 2004. <http://hillside.net/plop/2004/papers/mhafiz1/PLoP2004>.
 - Setuid Demystified Hao Chen, David Wagner, and Drew Dean. 11th USENIX Security Symposium, 2002.

Outline

- Two security problems and solutions in Windows and Linux
 - Compromise of high privilege program
 - Running code as other users

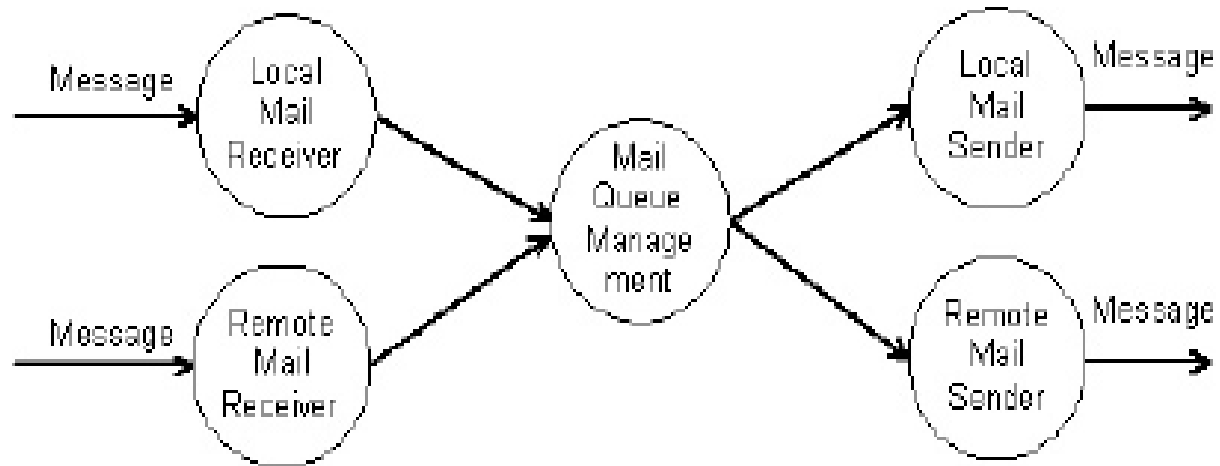
Problem: Exploit on High Privilege Program

- Attacker exploits bug in program or tricks user into running something unexpected
 - Exploits poor input processing on program
 - Surreptitiously causes exploit to be run when viewing mail
- Program is being run as high privilege user (e.g., root in Unix or Administrator in Windows)
 - Exploit is now also running at high privilege and can do most anything to the system

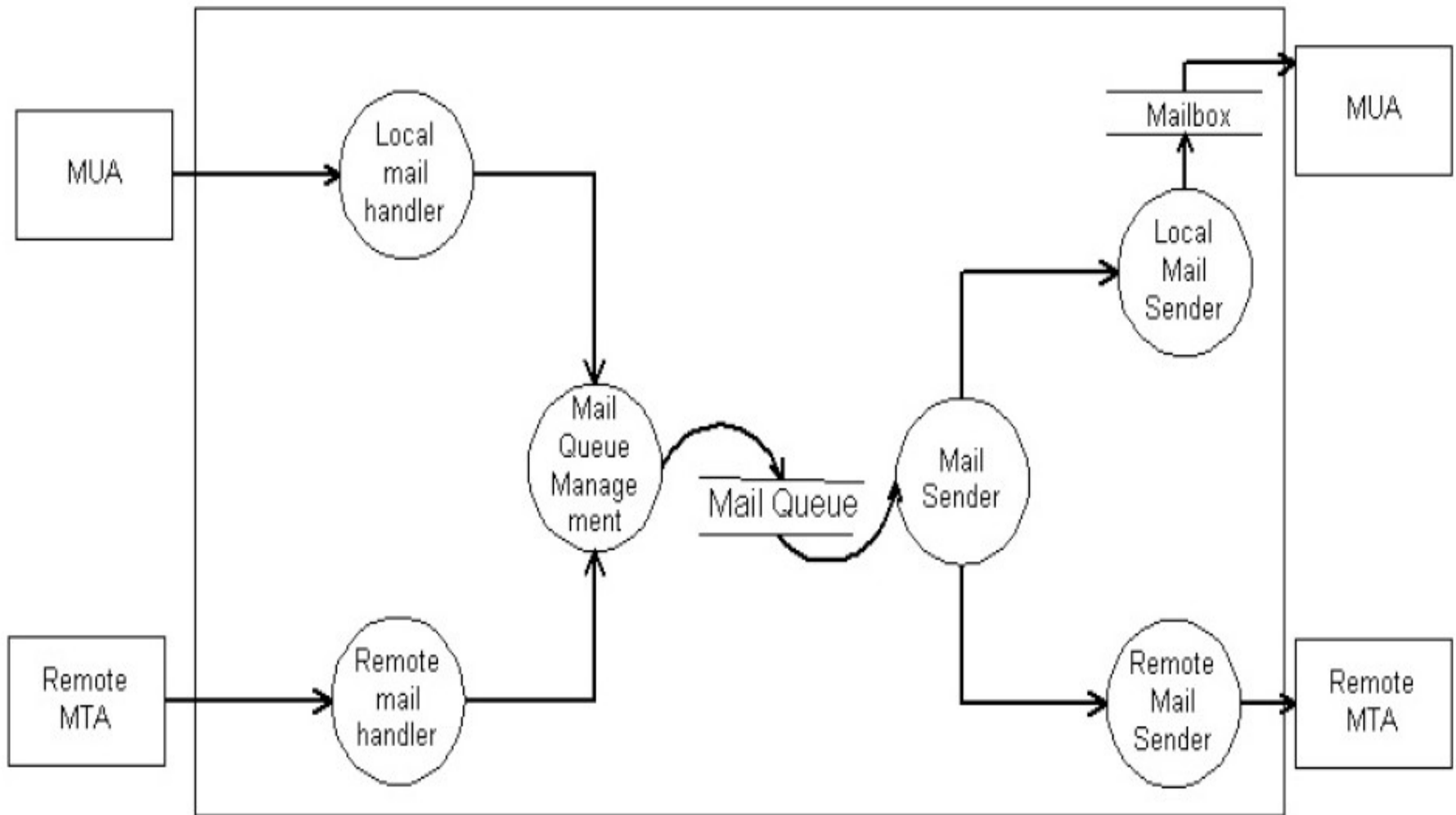
Solution: Modularity

- Divide program into smaller, communicating programs
 - Only subset of the processes need to run at high privilege
 - E.g., qmail as a redesigned MTA replacement for sendmail
- Get simplicity as a side effect
 - Easier to test and analyze for correctness

MTA structure



More MTA Structure



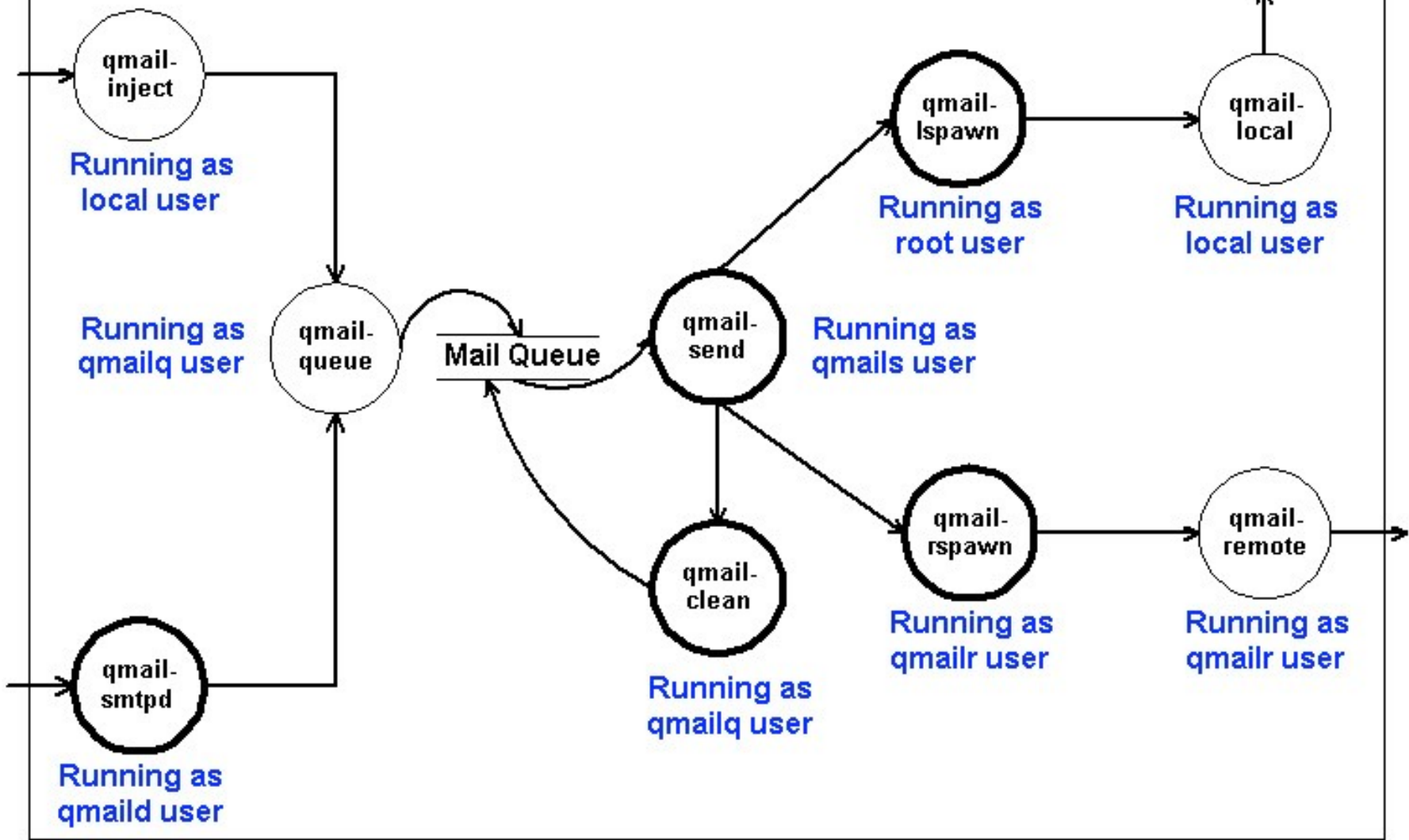
Security Patterns

- **Compartmentalization**
 - Failure in one part of system allows another part to be exploited
 - Put each part in separate security domain. If one part is compromised, the other parts remain secure
- **Distributed Responsibility**
 - A failure in a component can change any data in that component.
 - Partition data across components.

Legends

Spawned processes ○

Daemon processes ○



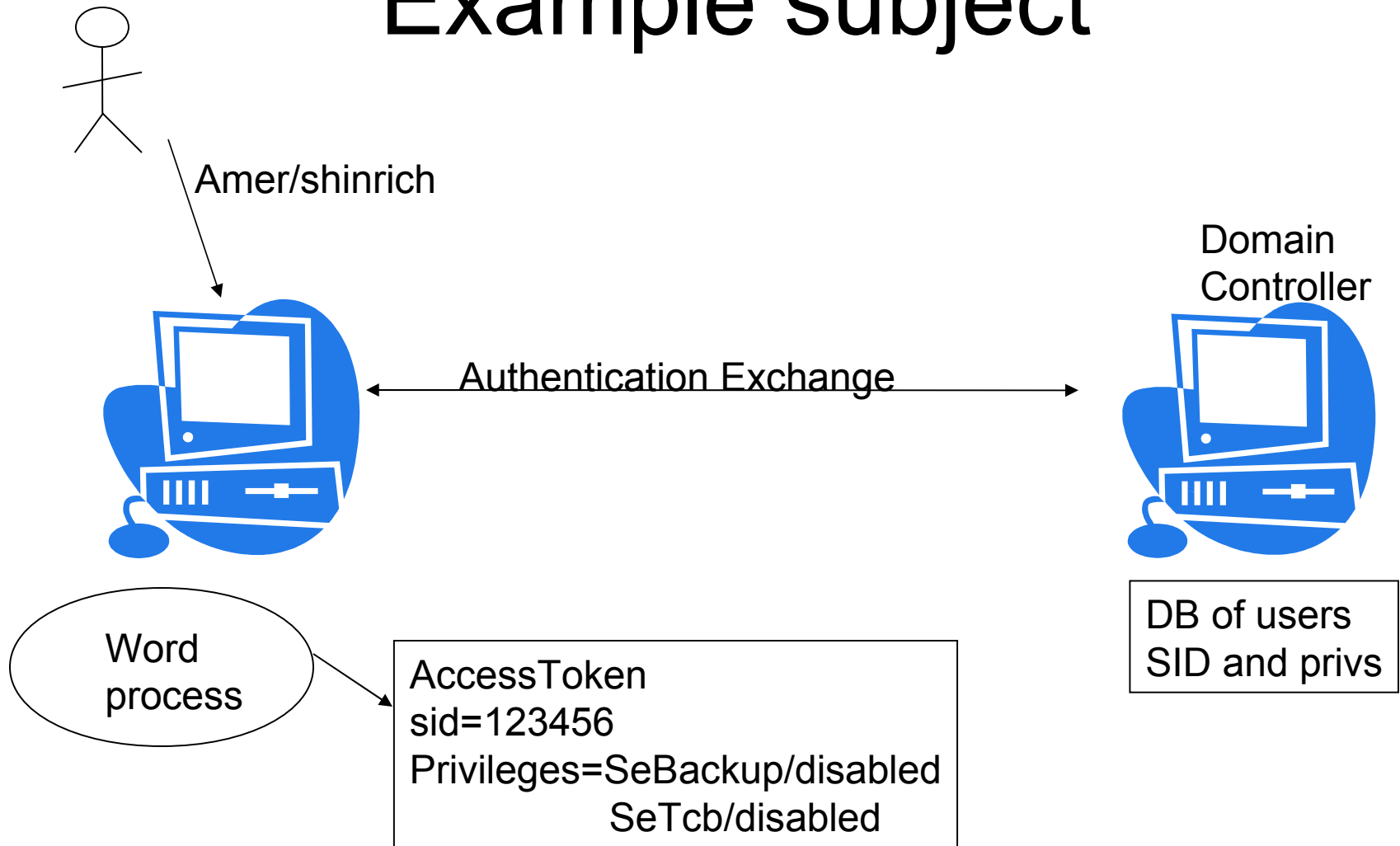
Solution: Least Privilege

- Even high privilege programs only need the extra powers for small parts of its execution
 - Turn off privilege when not needed
 - Permanently drop privileges that are never needed

Windows Security Elements

- Subject – Process or thread running on behalf of the system or an authenticated user
 - Security ID (SID) – A globally unique ID that refers to the subject (user or group)
 - Access token – the runtime credentials of the subject
 - Privilege – ability held by the subject to perform “system” operations. Usually breaks the standard security model
 - Associated with the access token
 - Generally disabled by default.
 - Can be enabled and disabled to run at least privilege
 - Example powerful privileges
 - **SeAssignPrimaryTokenPrivilege** – Replace process token
 - **SeBackupPrivilege** – Ignore file system restrictions to backup and restore
 - **SeIncreaseQuotaPrivilege** - Add to the memory quota for a process
 - **SeTcbPrivilege** – Run as part of the OS
 - Other privileges
- <http://msdn.microsoft.com/library/default.asp?url=/library/en-u>

Example subject



Running at reduced privilege

- Two system calls disable or remove privileges from the current access token
 - **AdjustTokenPrivileges** – enables/disables privileges
 - **CreateRestrictedToken** – permanently restrict or remove privileges

Example to Find Token Info

- ```
// find the buffer size
DWORD dwSize = 0;
PTOKEN_PRIVILEGES pPrivileges = NULL;
GetTokenInformation(hToken,
TokenPrivileges, NULL, dwSize, &dwSize);

// allocate the buffer
pPrivileges = (PTOKEN_PRIVILEGES)
GlobalAlloc(GPTR, dwSize);

// now that we have a buffer, try again
GetTokenInformation(hToken,
TokenPrivileges, pPrivileges, dwSize,
&dwSize);
```
- **MSDN pointer**  
[http://msdn.microsoft.com/en-us/library/aa446671\(VS.85\).aspx](http://msdn.microsoft.com/en-us/library/aa446671(VS.85).aspx)

# Linux/POSIX Privilege Model

- Privileges called capabilities
  - <http://www.linuxjournal.com/article/5737>
  - Each process has three capability sets
    - Effective – Set of currently activated privileges
    - Permitted – Set of privileges that process can use
    - Inheritable – Passed onto child processes created by exec
- Can remove capabilities globally
  - Global 32 bit mask that bounds capabilities that can be enabled on the system
  - `/proc/sys/kernel/cap-bound` can be accessed by `lcap` utility
  - `/usr/include/sys/capability.h`

# Example

- `lcap CAP_SYS_CHOWN`
  - Once done, it becomes impossible to change a file's owner:
- **`chown nobody test.txt`**
- *chown: changing ownership of `test.txt':*
- *Operation not permitted*



# Set of capabilities

- lcap
- Current capabilities: 0xFFFFDFCFF
- 0) \*CAP\_CHOWN
- 1) \*CAP\_DAC\_OVERRIDE
- 2) \*CAP\_DAC\_READ\_SEARCH
- 3) \*CAP\_FOWNER
- 4) \*CAP\_FSETID
- 5) \*CAP\_KILL
- 6) \*CAP\_SETGID
- 7) \*CAP\_SETUID
- 8) \*CAP\_SETPCAP
- 9) \*CAP\_LINUX\_IMMUTABLE
- 10) \*CAP\_NET\_BIND\_SERVICE
- 11) \*CAP\_NET\_BROADCAST
- 12) \*CAP\_NET\_ADMIN
- 13) \*CAP\_NET\_RAW
- 14) \*CAP\_IPC\_LOCK
- 15) \*CAP\_IPC\_OWNER
- 16) \*CAP\_SYS\_MODULE
- 17) CAP\_SYS\_RAWIO
- 18) \*CAP\_SYS\_CHROOT
- 19) \*CAP\_SYS\_PTRACE
- 20) \*CAP\_SYS\_PACCT
- 21) \*CAP\_SYS\_ADMIN
- 22) \*CAP\_SYS\_BOOT
- 23) \*CAP\_SYS\_NICE
- 24) \*CAP\_SYS\_RESOURCE
- 25) \*CAP\_SYS\_TIME
- 26) \*CAP\_SYS\_TTY\_CONFIG
- 27) \*CAP\_MKNOD
- 28) \*CAP\_LEASE
- 29) \*CAP\_AUDIT\_WRITE
- 30) \*CAP\_AUDIT\_CONTROL

\* ≡ Capabilities currently allowed

# Linux Privileges/Capabilities

- Can disable or remove capabilities per process
  - Libcap or setcap/getcap system calls
  - Can specify the affected process, the process group, or all processes
  - Can specify the capability mask for all three sets of capabilities
- Limited by lack of file system support

# Problem: Run privileged program portions as regular user

- File server program must have portions run at high privilege, but ultimately only returns information that the invoking user has access to
- More frequently allow low privilege user to run high privilege program

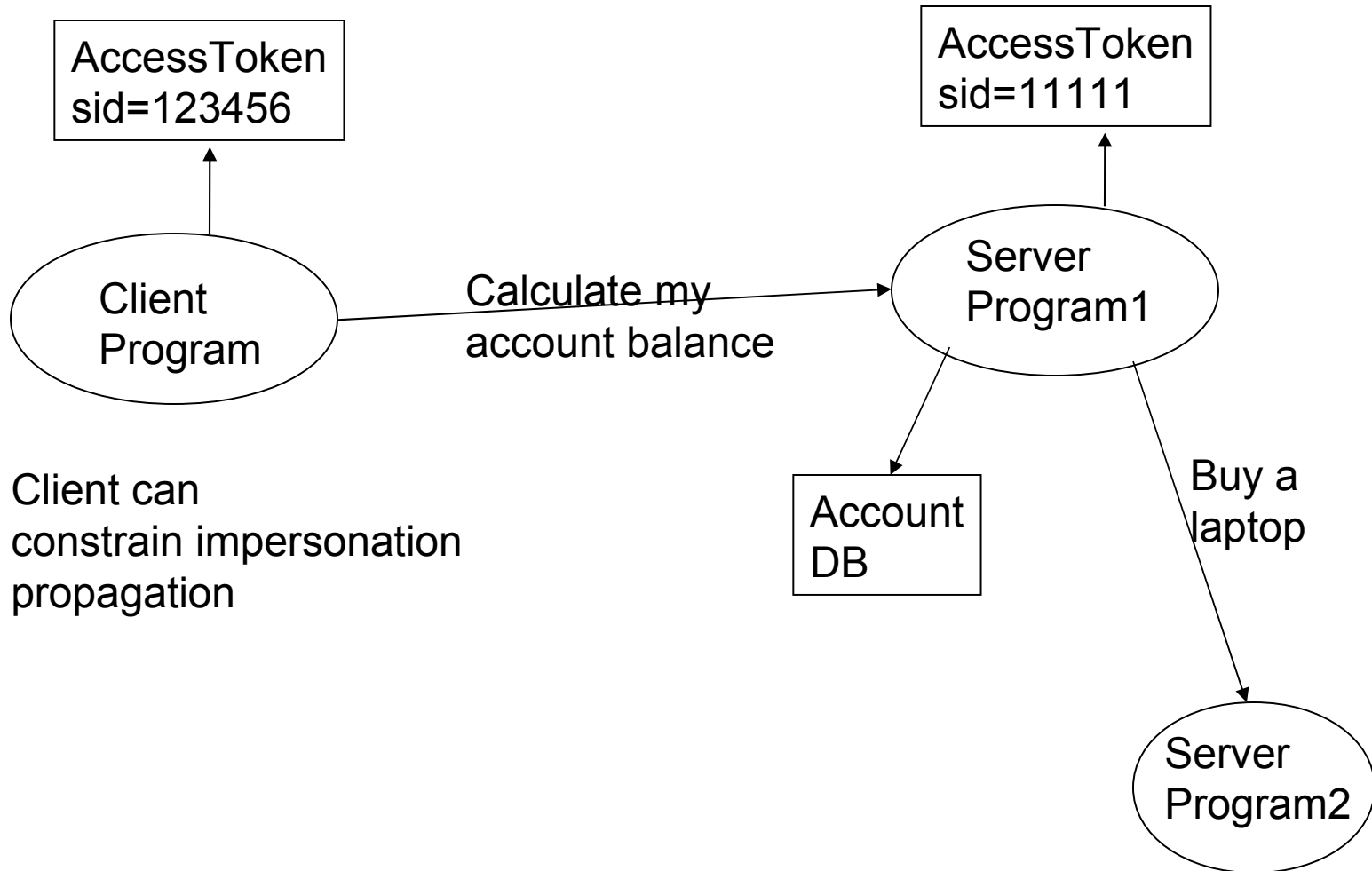
# Solution: Impersonation

- Client program runs as end user
- Client program communicates with privileged daemon or service
- Privileged service picks up client's identity
- “Impersonates” client while acting on behalf of the client

# Windows Impersonation

- Each process has three access tokens associated
  - Real access token
  - Effective access token
  - Saved access token
- Server program can run with client access token
  - **ImpersonateLoggedOnUser** - runs under the access token of the logged on user
    - Several variations of this system call which pull the impersonation token from various sources
  - **RevertToSelf** to return to the original user
  - **SelfImpersonatePrivilege** has been introduced
- Presumably client has lower privilege than server
- Multiple impersonation levels to restrict token propagation

# Example impersonation



# Impersonation problems

- Knowledgeable exploit can use `RevertToSelf`
- Base user is most likely a privileged user

# Solution: Set User ID

- Mark executable so it runs as a different user than the invoking user
  - Mark file system program to run as privileged user
- Rely on system calls to reset user ID to less privileged user



# Unix Set UID

- Each Unix process has three user ID's associated
  - Effective – Used in access checks
  - Real
  - Saved
- `setresuid` system call enables application to set all three
  - Assuming caller meets requirements, e.g., regular user cannot set UID to 0

# SetUID File Bit

- Normally, new process will run under UID of invoking process
- If SetUID bit is set
  - New process will run under executable File's UID for effective UID
  - Real UID will still be that of invoking user.
  - Setting SetUID bit is restricted for normal user

# Setting SetUID bit

- Consider executable Foo
  - Owned by Bob
    - What does this mean when run by Fred?
  - Owned by root
    - What does this mean when run by Fred?

# SetUID system calls

- This concept has been in Unix since the beginning
- The concept has evolved over time
  - Slightly different calls and semantics in different flavors of Unix
- In general for all flavors
  - Effective user ID of 0, can set effective UID to any value
  - Otherwise can only set effective UID to real or saved UID

# Unix Set UID

- Example
  - `setuid(getuid())`
  - Run as non-root user to permanently clear the root privilege
  - Simple API hides details and may reveal exploitable vulnerability