1. **Delay-Bandwidth Product for Links in Series**

Consider three nodes in series. Node A is connected to node B via a 250 Mbps fiber optic link, 3500 km in length. Node B is connected to node C via a 1 Mbps link, 2 km in length. The links are full duplex. The rate of transmission errors on the links, the time to switch a packet at node B, and the time to transmit an ACK are all negligible. A large file is to be sent from node A to node C, and there is no other traffic on the links. Packets are 1.5 KB, including headers.

*(Use Speed of Light = 3 \times 10^8 m/s)*

- a. Ignoring reliability and packet headers, what is the maximum throughput that can be achieved (in Mbps)? Explain.
- b. What is the round trip time from A to C?
- c. What is the roundtrip bandwidth delay product for the path from A to C? (Specify the units you use).
- d. Suppose an end-to-end sliding window protocol is used with SWS=RWS. What size of SWS is optimal?
- e. Why wouldn't you want SWS to be many times larger than the value you suggested in part d?

2. **Spanning Tree Algorithm for Intelligent Bridges**

Suppose the Perlman spanning tree algorithm and the bridge learning algorithm for forwarding are used for the network shown below.

```
B2
  B3
    B6
      B7
       B9
B4
  B1
    B5
      B8

A
  C

B
  D

B
  E

B
  F

B
  G

Venus
Jupiter
Mars
```

- a. Indicate which bridge is root, which ports are root ports (i.e. the preferred port for reaching the root bridge), which bridge is the designated bridge for each LAN, and which ports are designated ports (i.e. the ports that connect some LAN to its given
designated bridge). Hint: bridges that are not designated bridges for any LAN, and ports that are not either root ports or designated ports do not play a role in the routing of packets. The remaining bridges together with the LANs form a spanning tree.

Use a table with this format:

<table>
<thead>
<tr>
<th>Bridge</th>
<th>Root</th>
<th>Root Port</th>
<th>LANs</th>
</tr>
</thead>
<tbody>
<tr>
<td>B1</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>B2</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>B3</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>B4</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>B5</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>B6</td>
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<td>B7</td>
<td></td>
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<tr>
<td>B8</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>B9</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

b. Suppose after the configuration is complete, host Mars attaches to LAN B, host Venus attaches to LAN G, and host Jupiter attaches to LAN E. Suppose Mars sends a message to Venus, then Venus sends a message to Mars, then Jupiter sends a message to Venus. For each of the three messages, indicate which LANs the message is heard on.

3. Switch Fabrics

Banyan and Batcher networks are two types of self-routing fabrics often used to construct large switches from simpler components. A single stage of a $n \times n$ Banyan network consists of $n/2$ switches of dimension $2 \times 2$. An $n \times n$ Batcher network can be made from two Batcher networks of size $n/2 \times n/2$ plus a merge network with $n/2 \log_2 n$ switches.

a. For $n = 128$, how many stages are required to route packets from the inputs to the outputs of a Banyan Network?

b. How many $2 \times 2$ switches are required for the network in part a)?

c. Write down a recurrence relation $T(n)$ for the number of switches in a Batcher network of size $n \times n$.

d. Give the number of switches required for $n = 32$.

4. Server Bandwidth

Consider a server with direct memory access (DMA) in and out of main memory. Assume the server's I/O bus speed is 600Mbps and the memory bandwidth is 900Mbps.

a. How many switched 1.8Mbps links could be supported by the server?

b. Suppose the server switching time is such that it can forward packets at the rate of 1300 packets per second. Determine the throughput as a function of the packet size.

c. At what packet size does the I/O bus speed become the limiting factor?

5. Ethernet Timing

This problem is about the Ethernet/IEEE 802.11 access protocol. To be definite, suppose that if a host detects a transmission while it is transmitting a frame, then:

(i) if the host has already transmitted the 128 bit preamble, the host stops transmitting the frame and sends a 64 bit jamming sequence;
(ii) Else the host finishes transmitting the 128 bit preamble and then sends a 64 bit jamming sequence. For simplicity, assume a collision is detected as soon as an interfering signal first begins to reach a host.

Suppose the packets are 1024 bits long, which is the minimum length allowed. Hosts A and B are the only active hosts on a 10 Mbps Ethernet and the propagation time between them is 0.055ms, or 550 bit durations. Suppose A begins transmitting a frame at time $t = 0$, and just before the beginning of the frame reaches B, B begins sending a frame, and then almost immediately B detects a collision.

a. Does A finish transmitting the frame before it detects that there was a collision? Explain.

b. What time does A finish sending a jamming signal? What time does B finish sending a jamming signal?

c. What time does A first hear an idle channel again? What time does B first hear an idle channel again?

d. Suppose each host next decides to retransmit immediately after hearing the channel idle. After the resulting (second) collision: When does A next hear the channel idle? When does B next hear the channel idle?

e. Suppose after the second collision, A decides to wait 1024 bit durations to retransmit (if it hears silence after that long) and B decides to retransmit immediately after hearing a silent channel. Is the transmission of host B successful?

f. At the time A was planning to send its second retransmission, it senses a carrier present. Suppose at that particular time A decides to wait $2 \times 0.1024$ms more until its next retransmission. What time does host A finish sending its packet?

6. **Link-State Routing**

Show how the link-state algorithm builds the routing table for node A in the following network.

Use the same format as in P&D (Table 3.14, page 258).