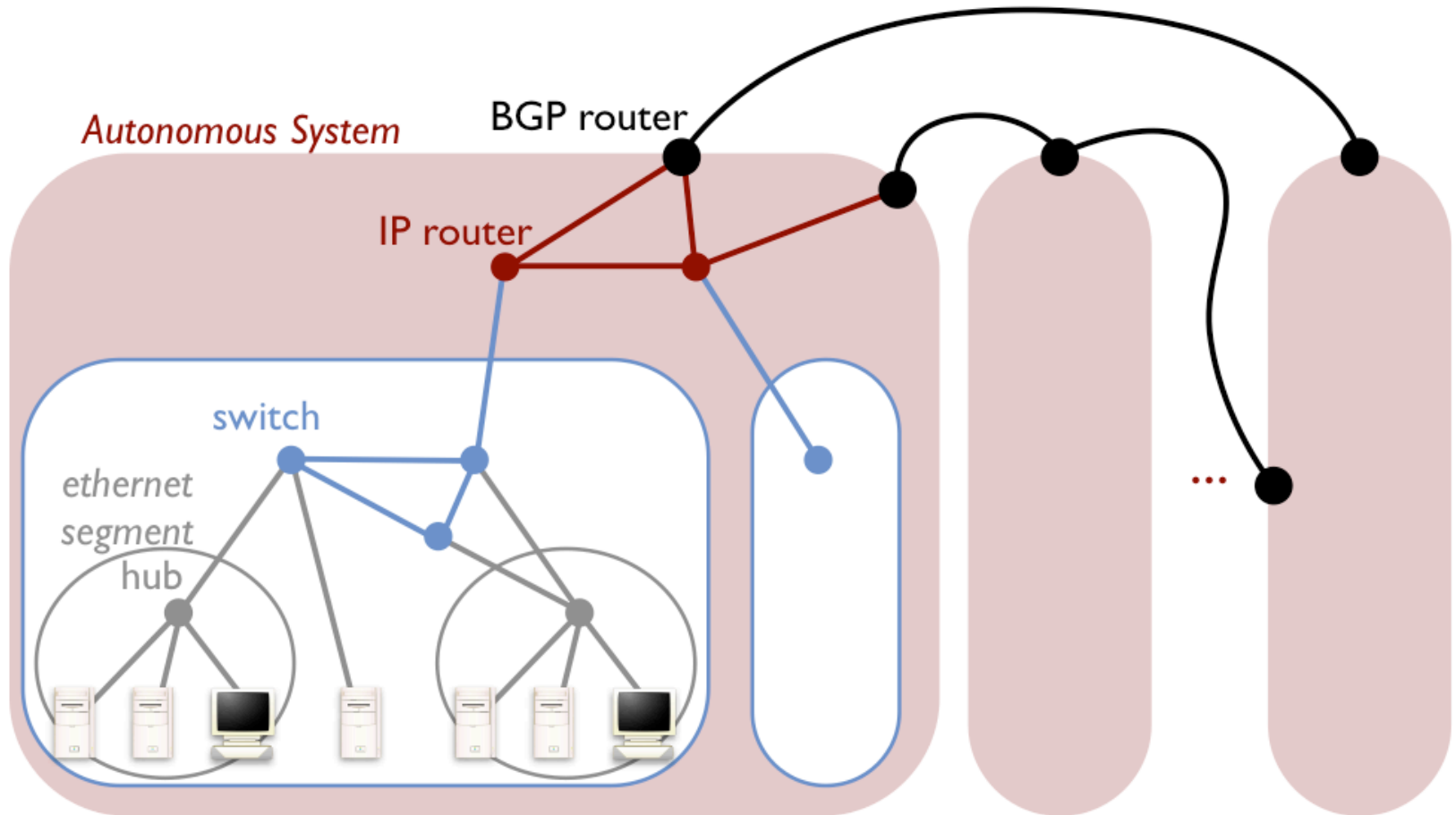




The Global Internet

[Big Picture of the Internet]



The Global Internet and Inter-domain Routing

- Why does Border Gateway Protocol (BGP) exist?
 - What is interdomain routing and why do we need it?
 - Why does BGP look the way it does?
- How does BGP work?
 - Path vector algorithm
 - Various boring details
- pay more attention to the “why” than the “how”



[Routing]

- Provides paths between networks
- We know several designs already
 - link-state
 - distance vector
- But previous lectures assumed single “domain”
 - all routers have same routing metric (shortest path)
 - no privacy issues, no policy issues



Internet is more complicated.....

- Internet not just unstructured collection of networks
- Internet is comprised of a set of “autonomous systems” (ASes)
 - Independently run networks, some are commercial ISPs
 - Autonomy of control: ex: company, university, etc
 - Currently around 35,000 Ases
- Enables hierarchical aggregation of routing information
- ASes are sometimes called “domains”
 - hence “interdomain routing”



[Autonomous Systems]

- Intradomain Routing (within an AS)
 - Performed using domain-specific algorithm
 - Selected by domain administrators
 - Allows heterogeneous interior gateway protocols (IGP)
- Interdomain Routing (between AS' s)
 - Performed using standard global algorithm
 - Homogeneous exterior gateway protocol (EGP)
 - Main goal: reachability



[Autonomous Systems]

- Common intradomain routing protocols
 - Routing Information Protocol (RIP)
 - From the early Internet
 - Part of Berkeley Software Distribution (BSD) Unix
 - Distance vector algorithm
 - Based on hop count (infinity set to 16 hops)
 - Open Shortest Path First (OSPF)
 - Internet Standard (RFC 2328)
 - Link state algorithm
 - Authenticates messages
 - Load balances across links

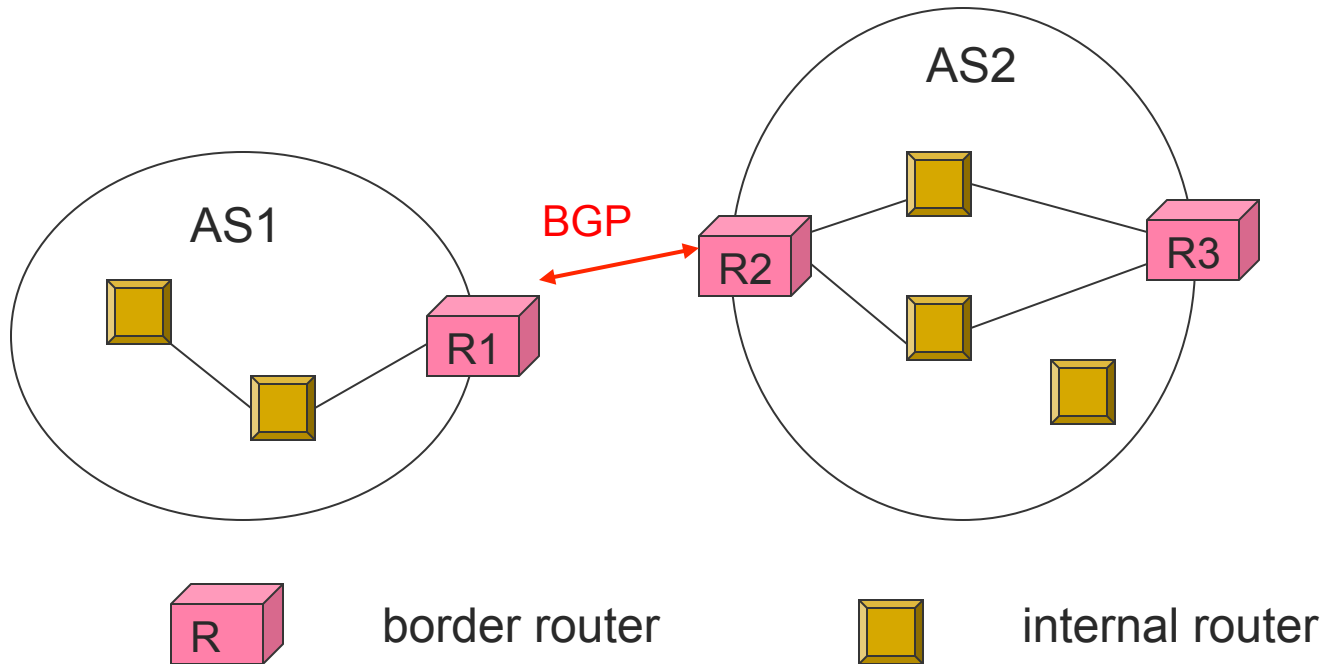


[Autonomous Systems]

- Standard interdomain routing protocols
 - General aspects
 - Very complex and difficult
 - Focuses on reachability rather than optimality
 - Must be loop free
 - Specify how reachability information should be exchanged
 - Exterior Gateway Protocol (EGP)
 - Defined on Internet with tree structure
 - Embodied (and enforced) tree structure
 - Had to be replaced eventually
 - Distance vector updates
 - Border Gateway Protocol (BGP)
 - Replaced EGP



[BGP]

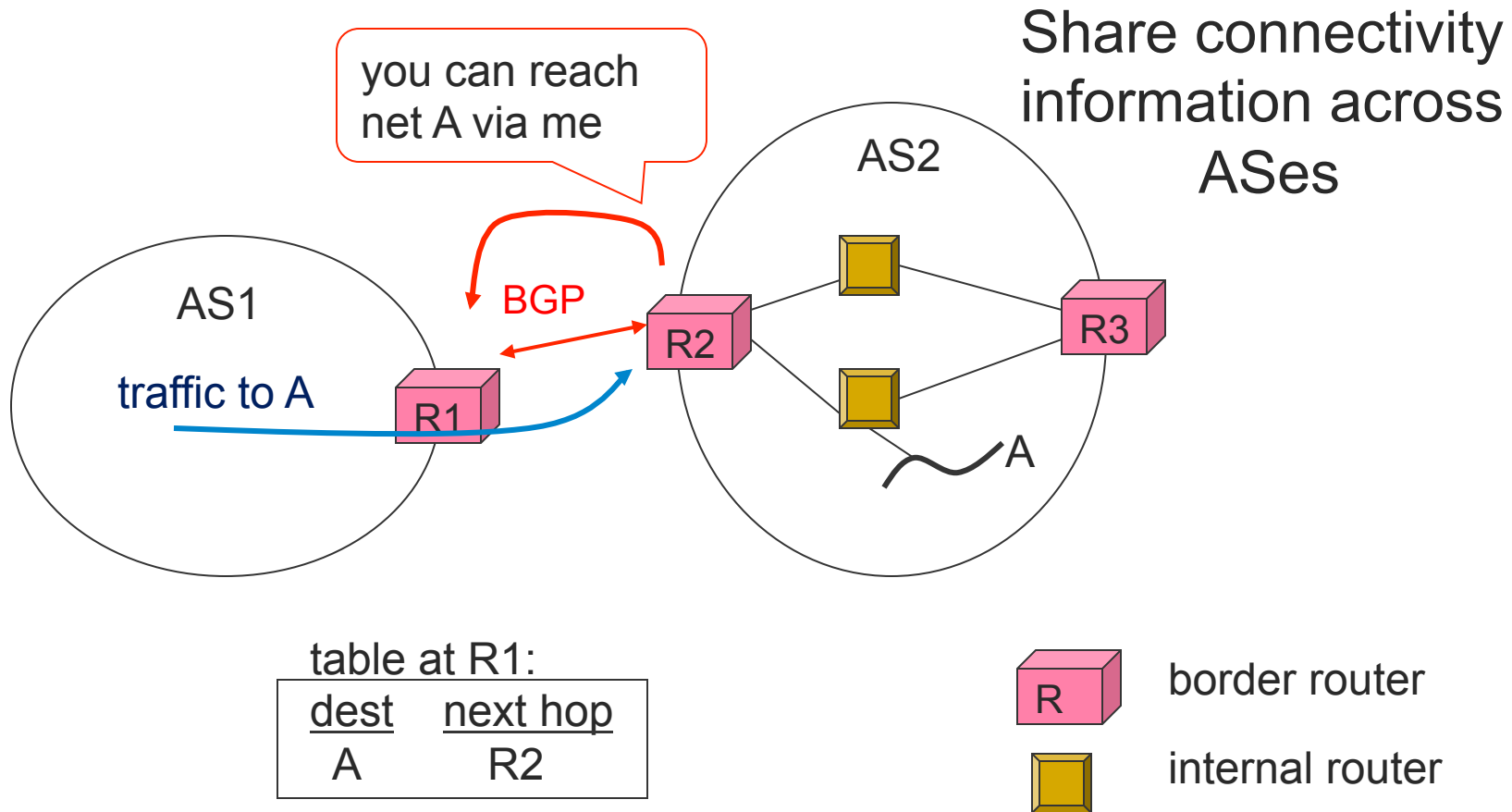


- Border routers

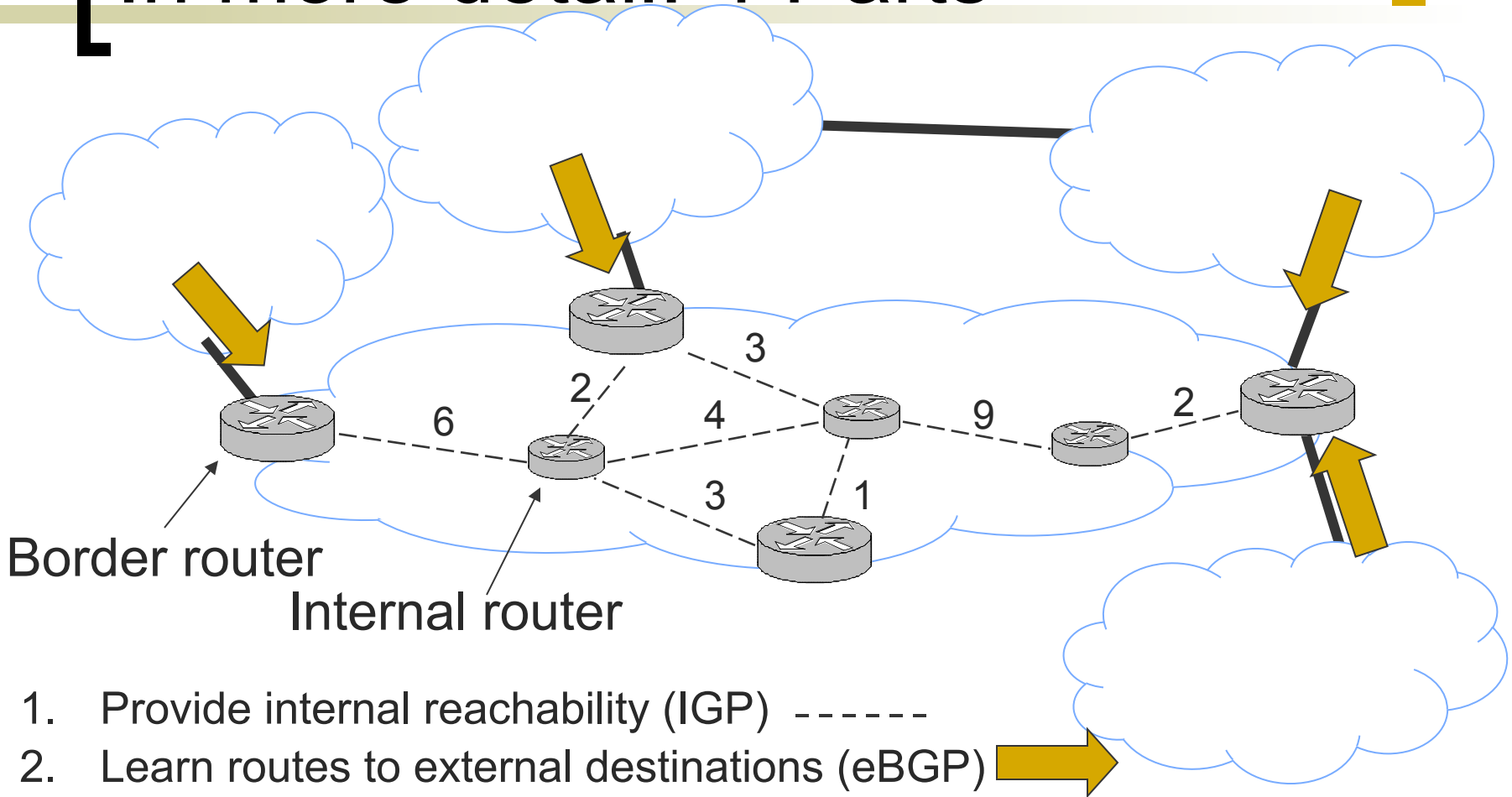
- Connects an AS to the Internet
- Used for default external route



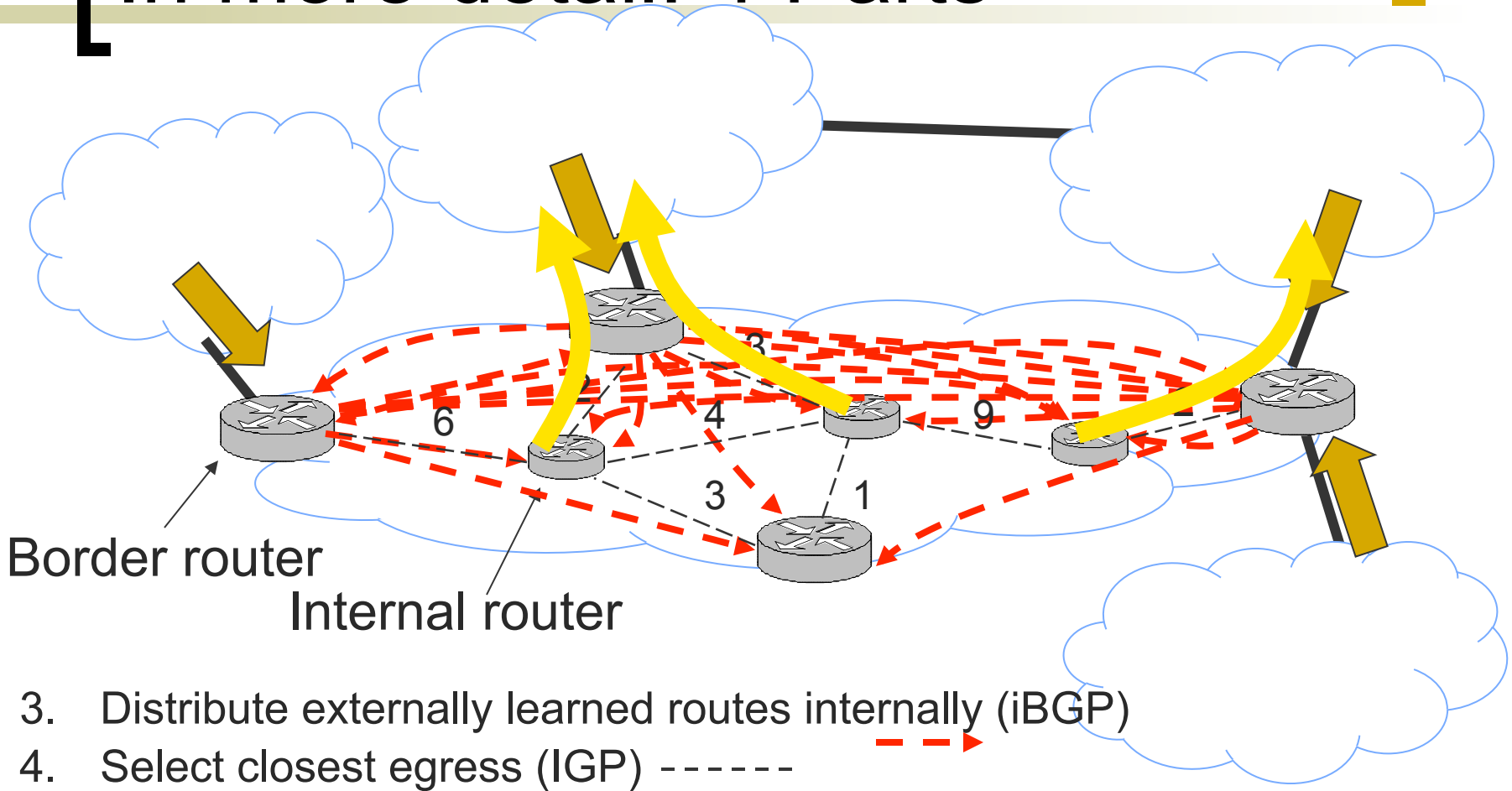
Autonomous Systems



[In more detail: 4 Parts]



[In more detail: 4 Parts]



The ‘A’ in AS really means Autonomous

- Want to choose their own internal routing protocol
 - Different algorithms and metrics
- Want freedom to route based on policy
 - “My traffic can’t be carried over my competitor’s network”
 - “I don’t want to carry transit traffic through my network”
 - Not expressible as Internet-wide “shortest path”!
- Want to keep their connections and policies private
 - Would reveal business relationships, network structure



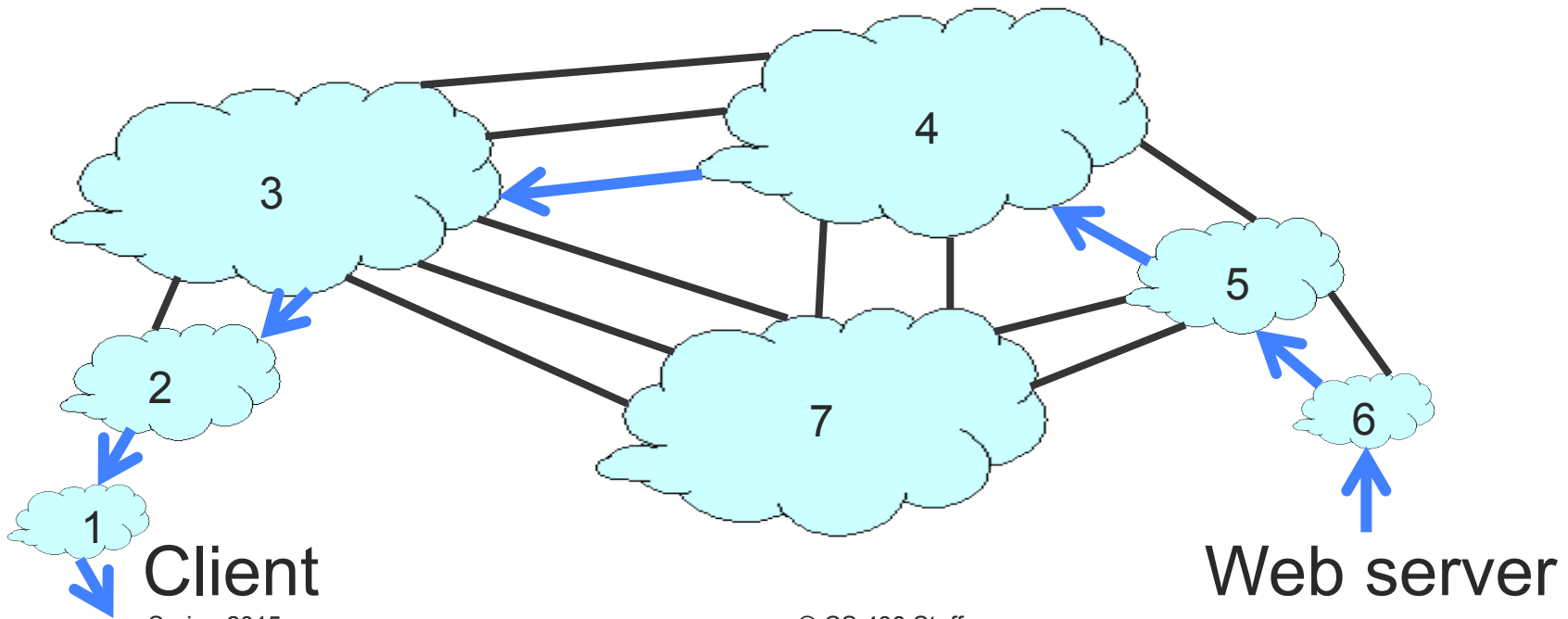
[AS' s are Buisnesses]

- Three kinds of common relationships between ASes
 - AS A can be AS B' s *customer*
 - AS A can be AS B' s *provider*
 - AS A can be AS B' s *peer*
- Business implications
 - Customer pays provider, peers don' t pay each other
- Policy implications
 - “When sending traffic, I prefer to route through customers over peers, and peers over providers”
 - “I don' t carry traffic from one provider to another provider”



[AS-Level Topology]

- Destinations are IP prefixes (e.g., 12.0.0.0/8)
- Nodes are Autonomous Systems (ASes)
- Links are connections & business relationships



[Autonomous Systems]

■ Challenges

○ Scale

- Border router must be able to forward any packet destined anywhere in the Internet

○ Autonomous routing in AS' s

- Impossible to calculate meaningful costs for paths that cross multiple AS' s

○ Trust

- One AS must trust the advertised routes of other AS' s

■ Goal

- Specify policies that lead to “good” paths (even if they are not optimal)



[Routing Choices]

- Key issues are *policy* and *privacy*
- Challenges
 - No universal metric
 - AS-specific Policy decisions
- Problems with link state
 - Metric used by routers not the same
 - Can't use shortest path - loops
 - LS database too large - entire Internet
 - Flooding may expose internal topology and policies to other AS' s



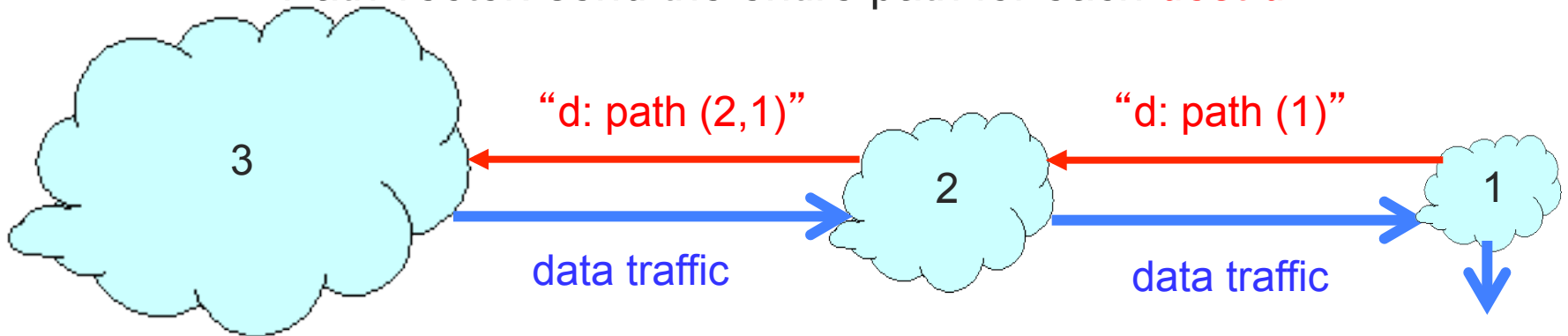
[Routing Choices]

- Key issues are *policy* and *privacy*
- Challenges
 - No universal metric
 - AS-specific Policy decisions
- Problems with distance-vector
 - Does not reveal any connectivity information
 - But still uses shortest path
 - Slow to converge



Solution: Path Vector Routing

- Extension of distance-vector routing
 - Support flexible routing policies
 - Faster loop detection (no count-to-infinity)
- Key idea: advertise the entire path
 - Distance vector: send distance metric per **dest d**
 - Path vector: send the entire path for each **dest d**



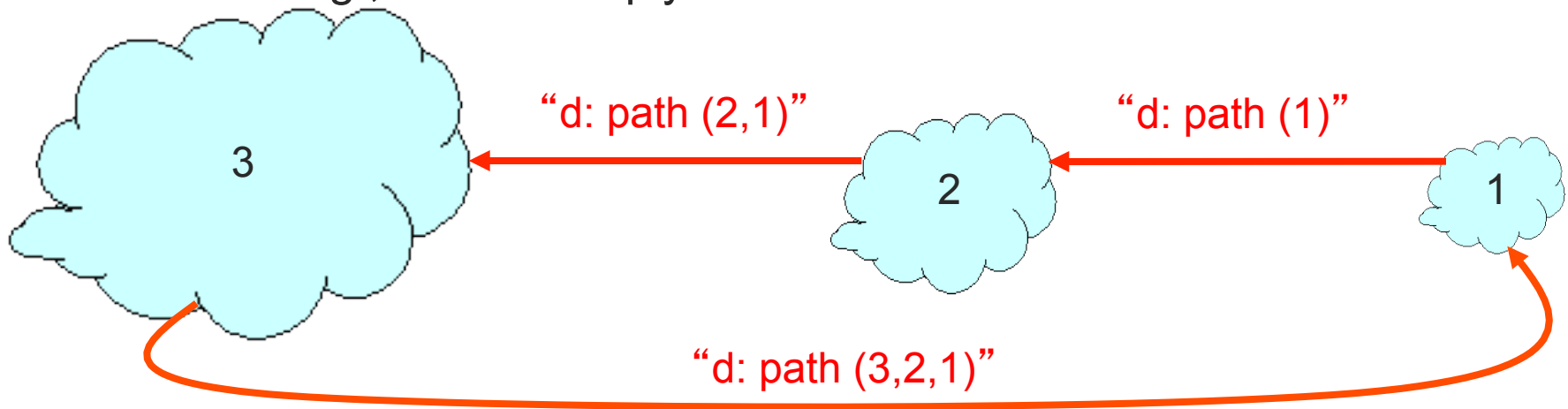
[Path Vectors]

- Each routing update carries the entire path
- Loops are detected as follows
 - When AS gets route check if AS already in path
 - If yes, reject route
 - If no, add self and (possibly) advertise route further
- Advantage
 - Metrics are local
 - AS chooses path, protocol ensures no loops



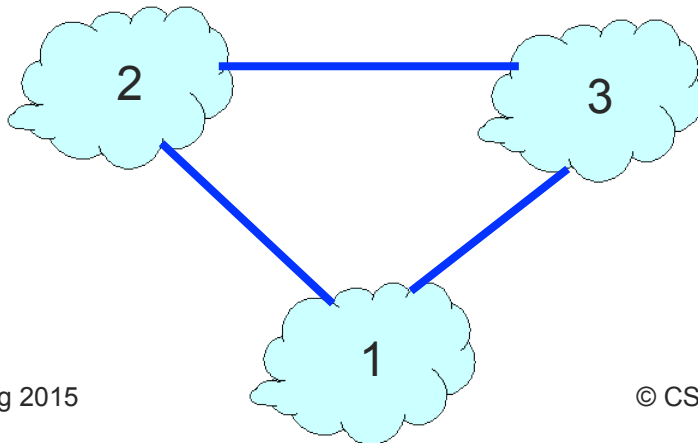
[Loop Detection]

- Node can easily detect a loop
 - Look for its own node identifier in the path
 - e.g., node 1 sees itself in the path “3, 2, 1”
- Node can simply discard paths with loops
 - e.g., node 1 simply discards the advertisement



Flexible Policies

- Each node can apply local policies
 - Path selection: Which path to use?
 - Path export: Which paths to advertise?
- Examples
 - Node 2 may prefer the path “2, 3, 1” over “2, 1”
 - Node 1 may not let node 3 hear the path “1, 2”

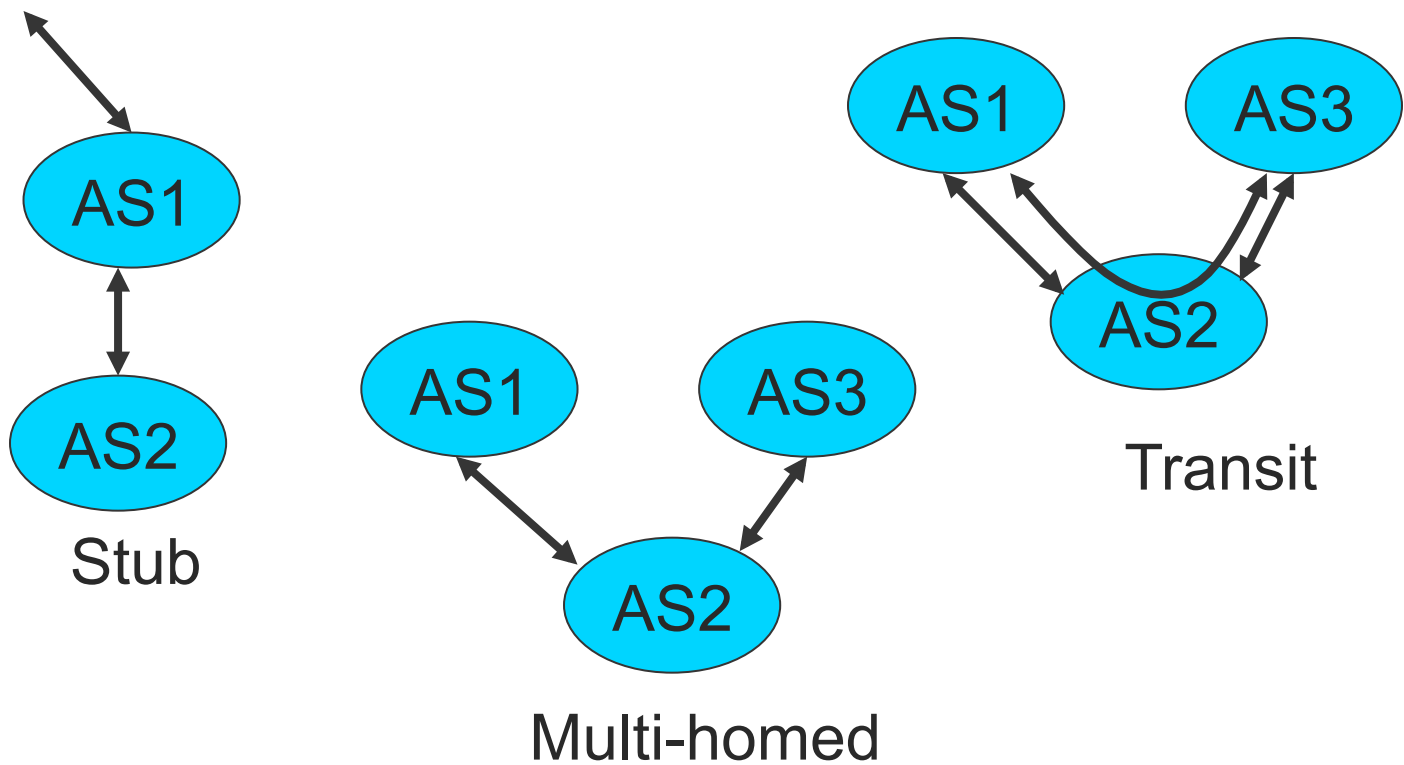


[AS Categories]

- Stub
 - An AS that has only a single connection to one other AS - carries only local traffic
- Multi-homed
 - An AS that has connections to more than one AS, but does not carry transit traffic
- Transit
 - An AS that has connections to more than one AS, and carries both transit and local traffic (under certain policy restrictions)



[AS Categories]



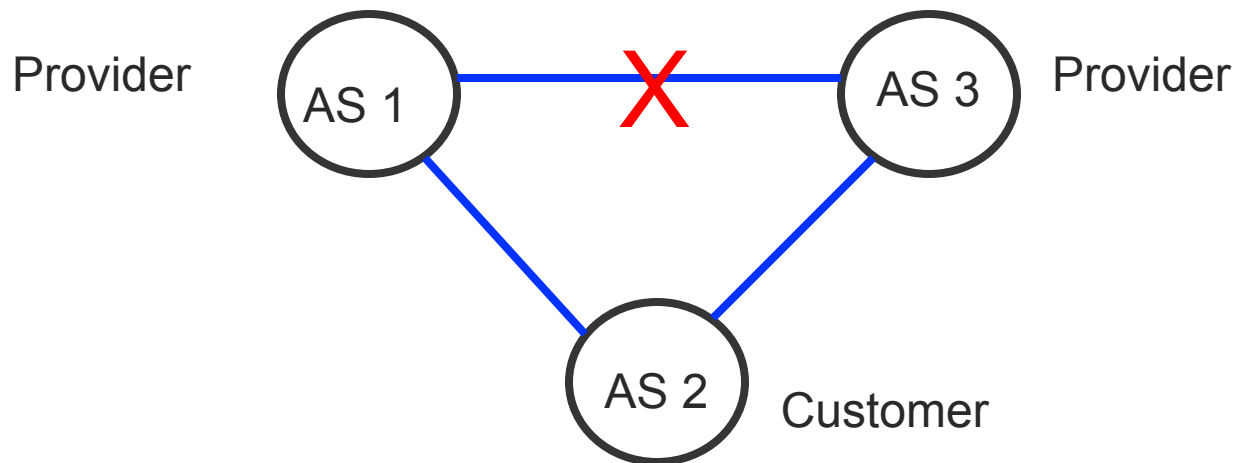
Issues with Path-Vector Policy Routing

- Reachability
- Security
- Performance
- Lack of isolation
- Policy oscillations



[Reachability]

- Normal routing
 - If graph is connected, reachability is assured
- Policy routing
 - Does not always hold



[Security]

- An AS can claim to serve a prefix that they actually don't have a route to (blackholing traffic)
 - Problem not specific to policy or path vector
 - Important because of AS autonomy
- Even worse: snoop on all traffic to almost any destination
 - Without anyone realizing that anything is wrong
- Fixable: make ASes “prove” they have a path
 - But not used in today's Internet



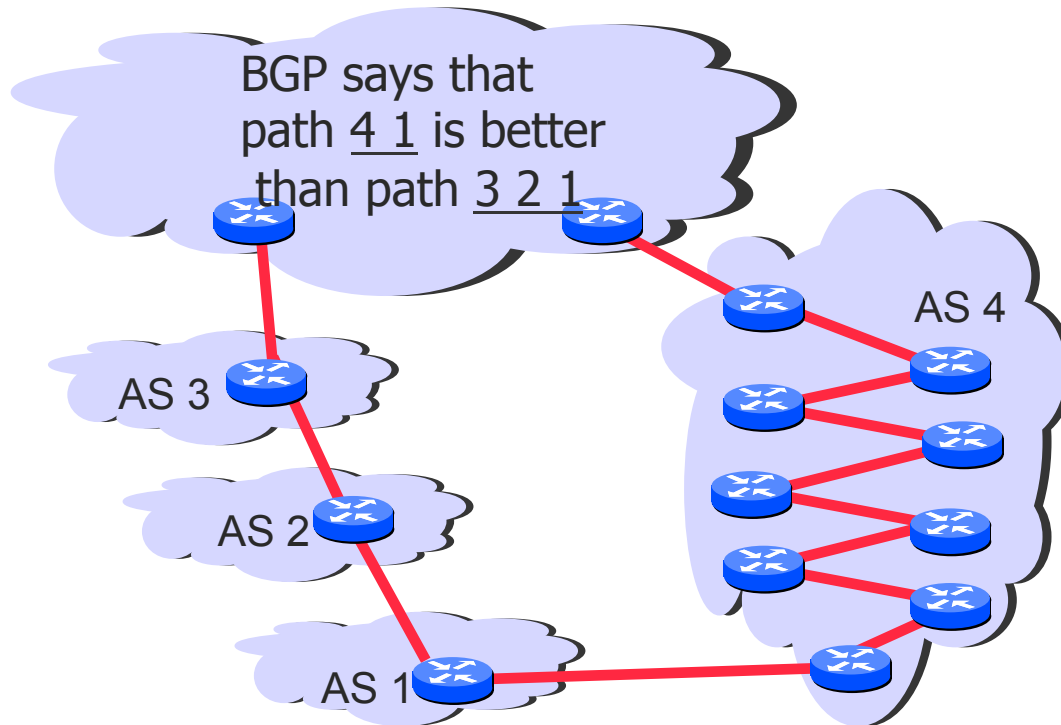
[Performance]

- BGP designed for policy not performance
- “Hot Potato” routing common but suboptimal
 - AS wants to hand off the packet as soon as possible
- Even BGP “shortest paths” are not shortest
 - Fewest AS' s != Fewest number of routers
- 20% of paths inflated by at least 5 router hops
- Not clear this is a significant problem



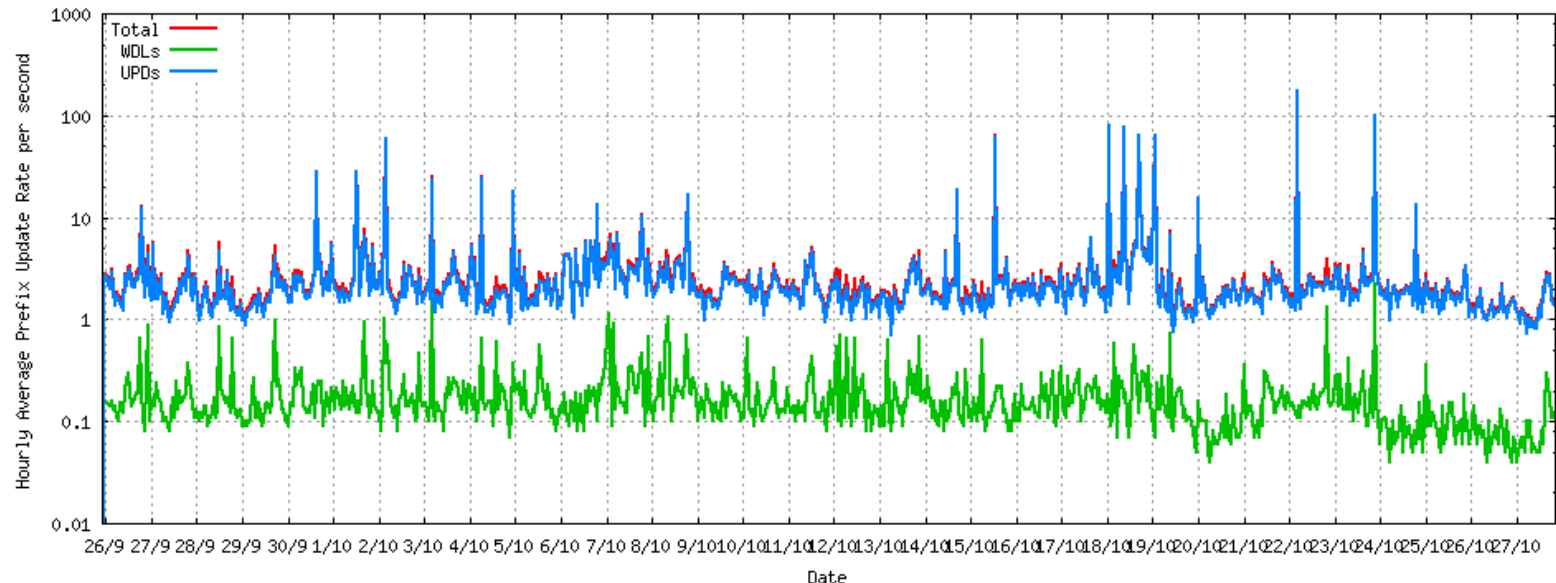
Performance

- AS path length can be misleading
 - An AS may have many router-level hops



Lack of Isolation: Dynamics

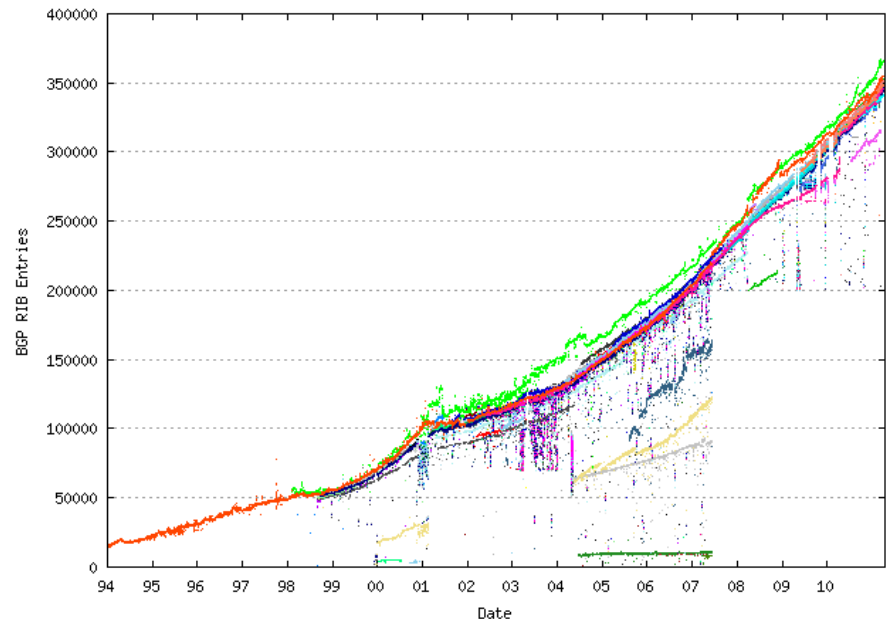
- Change in the path
 - Path must be re-advertised to every node using the path
 - “Route Flap Damping” supposed to help (but ends up causing more problems)



Lack of isolation: Routing Table Size

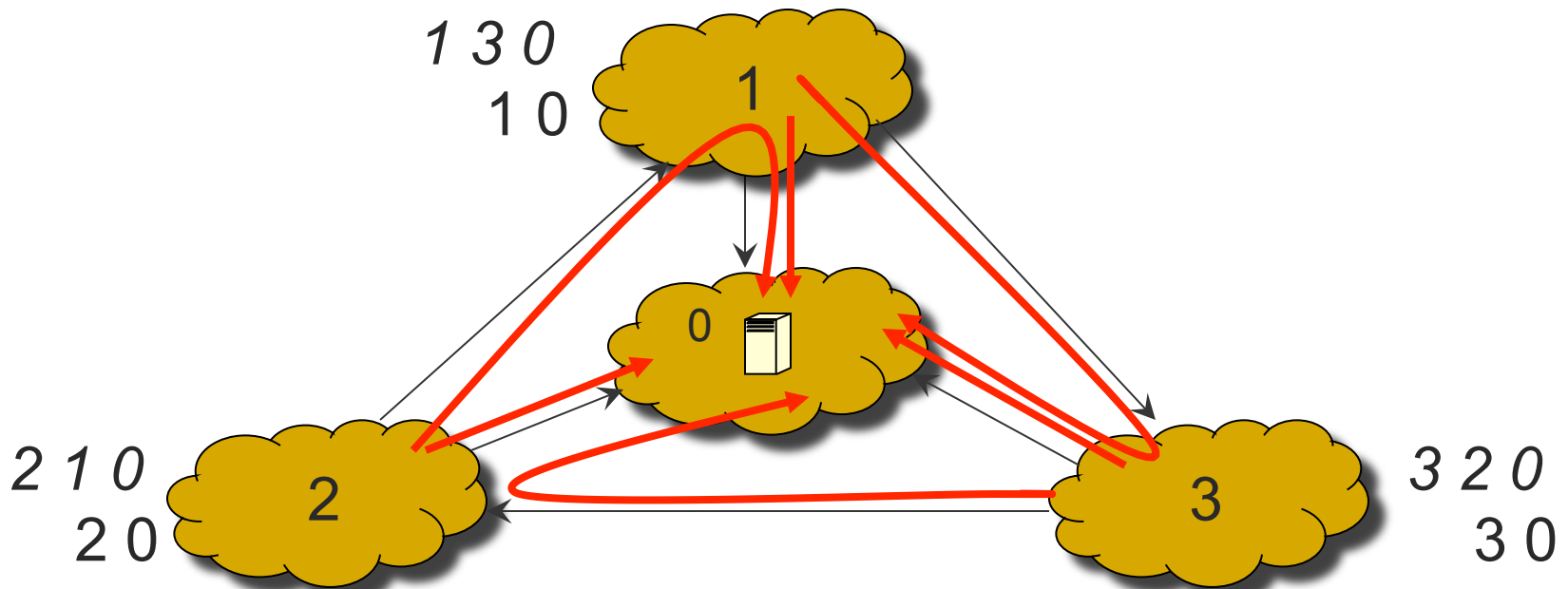
- Each BGP router must know path to every other IP prefix
 - But router memory is expensive and thus constrained
- Number of prefixes growing more than linearly
- Subject of current research

Number of prefixes in BGP table



Persistent Oscillations due to Policies

Depends on the interactions of policies



We are back to where we started!



[Policy Oscillations]

- Policy autonomy vs network stability
 - Focus of much recent research
- Not an easy problem
 - Difficult to decide whether given policies will eventually converge!
- However, if policies follow normal business practices, stability is guaranteed



Border Gateway Protocol (BGP)

- Interdomain routing protocol for the Internet
 - Prefix-based path-vector protocol
 - Policy-based routing based on AS Paths
 - Evolved during the past 15 years
- 1989 : BGP-1 [RFC 1105]
 - Replacement for EGP (1984, RFC 904)
- 1990 : BGP-2 [RFC 1163]
- 1991 : BGP-3 [RFC 1267]
- 1995 : BGP-4 [RFC 1771]
 - Support for Classless Interdomain Routing (CIDR)



BGP's job: maintain routing table

```
ner-routes>show ip bgp
```

```
BGP table version is 6128791, local router ID is 4.2.34.165
```

```
Status codes: s suppressed, d damped, h history, * valid, > best, i - internal
```

```
Origin codes: i - IGP, e - EGP, ? - incomplete
```

Network	Next Hop	Metric	LocPrf	Weight	Path
* i3.0.0.0	4.0.6.142	1000	50	0	701 80 i
* i4.0.0.0	4.24.1.35	0	100	0	i
* i12.3.21.0/23	192.205.32.153	0	50	0	7018 4264 6468 ?
* e128.32.0.0/16	192.205.32.153	0	50	0	7018 4264 6468 25 e



[BGP Operations]

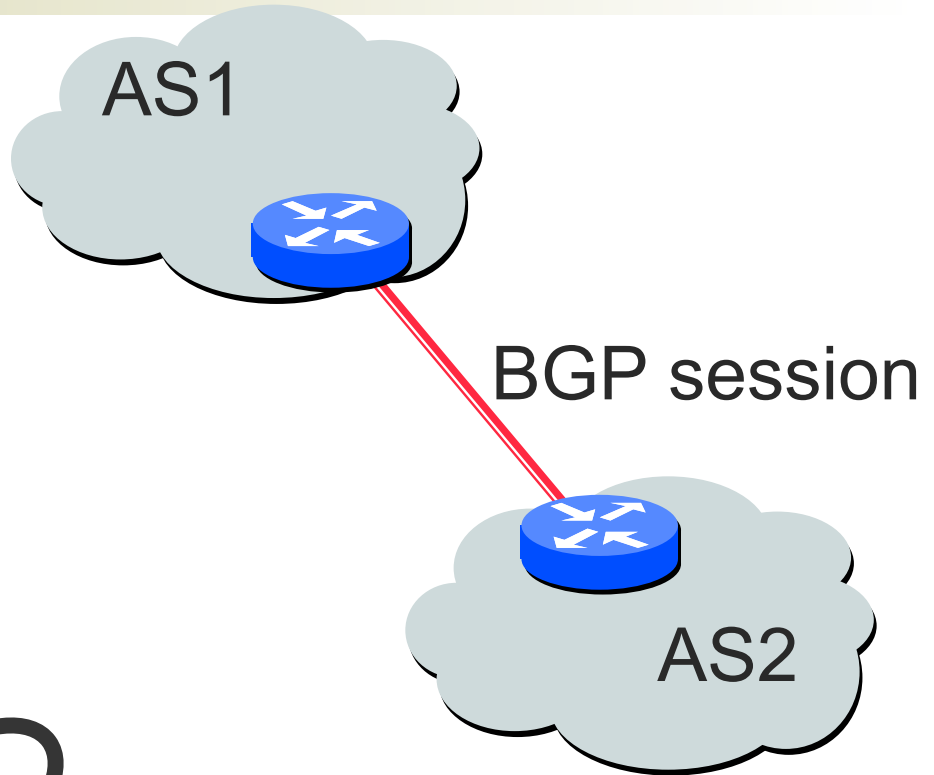
Establish session on TCP port 179



Exchange all active routes



Exchange incremental updates



While connection is ALIVE exchange route UPDATE messages



[Incremental Protocol]

- A node learns multiple paths to destination
 - Stores all of the routes in a routing table
 - Applies policy to select a single active route
 - ... and may advertise the route to its neighbors
- Incremental updates
 - Announcement
 - Upon selecting a new active route, add node id to path
 - ... and (optionally) advertise to neighbors
 - Withdrawal
 - If the active route is no longer available
 - ... send a withdrawal message to the neighbors



[BGP Route Processing]

Open ended programming.
Constrained only by vendor configuration language

Receive
BGP
Updates

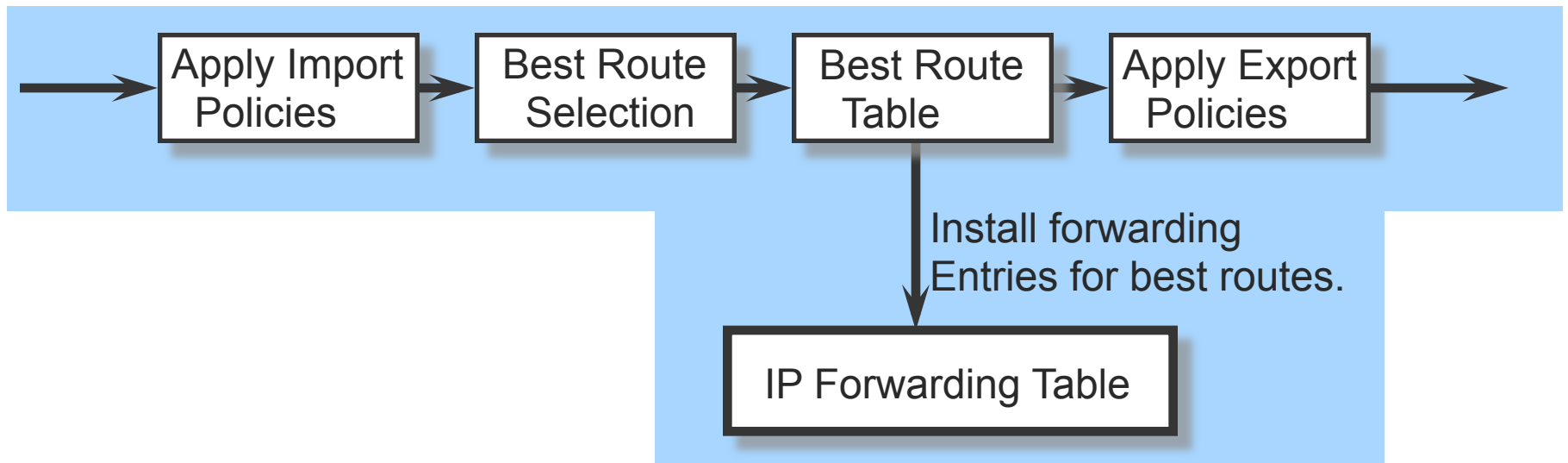
Apply Policy =
filter routes &
tweak attributes

Based on
Attribute
Values

Best
Routes

Apply Policy =
filter routes &
tweak attributes

Transmit
BGP
Updates



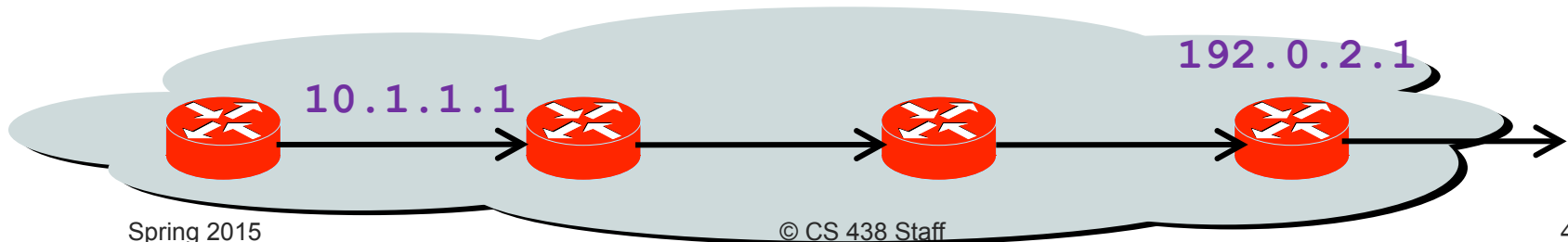
Selecting the best route

- Route Attributes
 - Set/modified according to operator instructions
- Route Choices
 - Compared based on attributes using (mostly) standardized rules



Joining BGP and IGP Information

- Border Gateway Protocol (BGP)
 - Announces reachability to external destinations
 - Maps a destination prefix to an egress point
 - `128.112.0.0/16` reached via `192.0.2.1`
- Interior Gateway Protocol (IGP)
 - Used to compute paths within the AS
 - Maps an egress point to an outgoing link
 - `192.0.2.1` reached via `10.1.1.1`



[Summary]

- BGP is essential to the Internet
 - ties different organizations together
- Poses fundamental challenges....
 - leads to use of path vector approach
- ...and myriad details

