a)

There are two shortest routes:
01010, 01110, 01111
01010, 01011, 01111

1b)
4 disjoint shortest paths. The number of node disjoint shortest paths is bounded above by the Hamming distance " d " between the two nodes.

At the final step, to get closer to the destination, the bit to be flipped must be a bit that is different. There are only $d$ nodes that differ from the target node in just one position, which the source node also differs from the target in. Then node-disjointness requires that each new path "uses up" one of those 2nd-to-last nodes.

One possible solution
00101, 10101, 10001, 10000, 10010
00101, 00001, 00000, 00010, 10010
00101, 00111, 10111, 10011, 10010
00101, 00100, 10100, 10110, 10010
2) Additions in the solution below are modulo 256.

The nodes are $9,10,20,24,90,105$ and 245 .
(a) $\mathrm{n}=9$

| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{n}+2^{\wedge} \mathrm{i}$ | 10 | 11 | 13 | 17 | 25 | 41 | 73 | 137 |
| $\mathrm{Ft}(\mathrm{i})$ | 10 | 20 | 20 | 20 | 90 | 90 | 90 | 245 |

(b)

Finger table of node 90 :

| I | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{n}+2^{\wedge} \mathrm{i}$ | 91 | 92 | 94 | 98 | 106 | 122 | 154 | 218 |
| $\mathrm{Ft}(\mathrm{i})$ | 105 | 105 | 105 | 105 | 245 | 245 | 245 | 245 |

Finger table of node 105:

| I | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{n}+2^{\wedge} \mathrm{i}$ | 106 | 107 | 109 | 113 | 121 | 137 | 169 | 233 |
| $\mathrm{Ft}(\mathrm{i})$ | 245 | 245 | 245 | 245 | 245 | 245 | 245 | 245 |

The path: 9, 90, 105, 245
3)
(a) Yes
(b) Yes.
(c) False.
(d) True: Write3(Y,5) -> Read2(Y,5) -> Read2(X,2).

Recommended exercise:

In the Chord p2p network above, determine which nodes will store the keys with the following hashed identifiers: 9, 10, 11, 29, 248 and 255.

Answer: 9, 10, 20, 90, 9, 9

