

Bakery Algorithm

1

Code for entry section:

```

Choosing[i] := true
Number[i] := max{Number[0], ..., Number[n-1]} + 1
Choosing[i] := false
for j := 0 to n-1 (except i) do
    wait until Choosing[j] = false
    wait until Number[j] = 0 or
        (Number[j],j) > (Number[i],i)
endfor
    
```

Code for exit section:

```

Number[i] := 0
    
```

2-Processor Mutex Algorithm

2

Code for entry section:

```

1 W[i] := 0
2 wait until W[1-i] = 0 or Priority = i
3 W[i] := 1
4 if (Priority = 1-i) then
5     if (W[1-i] = 1) then goto Line 1
6     else wait until (W[1-i] = 0)
    
```

Code for exit section:

```

7 Priority := 1-i
8 W[i] := 0
    
```

Deadlocks

□ Necessary conditions for deadlocks

- Non-shareable resources (locked objects)
- No preemption on locks
- Hold & Wait
- Circular Wait (Wait-for graph)

Validation of Transactions

Backward validation of transaction T_v

```

boolean valid = true;
for (int  $T_i = startT_n+1; T_i \leq finishT_n; T_i++$ ) {
    if (read set of  $T_v$  intersects write set of  $T_i$ ) valid = false;
}
    
```

Forward validation of transaction T_v

```

boolean valid = true;
for (int  $T_{id} = active1; T_{id} \leq activeN; T_{id}++$ ) {
    if (write set of  $T_v$  intersects read set of  $T_{id}$ ) valid = false;
}
    
```

Instructor's Guide for Coullouris, Dollimore, Kindberg and Blair, Distributed Systems: Concepts and Design, 5th ed., © Pearson Education 2012

Link Reversal Algorithm

5

Links are bi-directional
But algorithm imposes logical directions on them

Maintain a directed acyclic graph (DAG) for each destination, with the destination being the **only sink**

This DAG is for **destination node D**

Peer pointers (2): finger tables

Finger Table at N80

i	$f_t[i]$
0	96
1	96
2	96
3	96
4	96
5	112
6	16

Say $m=7$

i th entry at peer with id n is first peer with id $\geq n + 2^i \pmod{2^m}$