The time for this MP is rather short! So please start early! Start now!
You must work in groups of two for this MP. Please stick with the groups you formed for MP1. (see end of document for expectations from group members.)

Covfefe! Inc. (MP1) just got acquired by the fictitious FakeNews Inc. (this company’s business model is to detect, not generate, fake news – go figure!). This company liked your previous work while you were hired by Covfefe! Inc., so they’ve commissioned you to build a distributed group membership service for them.

You must work in groups of two for this MP.

This service maintains, at each machine in the system (at a daemon process), a list of the other machines that are connected and up. This membership list is a full membership list, and needs to be updated whenever:
1. A machine (or its daemon) joins the group;
2. A machine (or its daemon) voluntarily leaves the group; and
3. A machine (or its daemon) crashes from the group (you may assume that the machine does not recover for a long enough time).

There is only one group at any point of time. Since we’re implementing the crash/fail-stop model, when a machine rejoins, it must do so with an id that includes a timestamp - this distinguishes successive incarnations of the same machine (these are the ids held in the membership lists). Notice that the id also has to contain the IP address.

A machine failure must be reflected in at least one membership lists within 5 seconds (assuming synchronized clocks) – this is called time-bounded completeness, and it must be provided no matter what the network latencies are. A machine failure, join or leave must be reflected within 6 seconds at all membership lists, assuming small network latencies.

You’re told that at most three machines can fail simultaneously, and after three back-to-back failures, the next set of failure(s) don’t happen for at least 20 seconds (i.e., enough time for your system to converge back to a good topology). Your system must ensure completeness for all such failures (up to three
simultaneous failures). That is, whenever a process fails (and there are up to 3 simultaneous failures in the system), all membership lists that contained the process must be updated within 20 s.

*Your algorithm must be scalable to large numbers of machines.* For your experiments however, you may assume that you have \( N > 5 \) machines in the group at any given time. Note that this is not a limit on the set of machines eligible to be group members. Typical runs will involve about 7-10 VMs.

You **must use:** 1) a **ring as a backbone** (either uni- or bi-directional), and 2) the **ping-ack** style of failure detection discussed in class (SWIM-style, but without any indirect pings). DO NOT use heartbeating. Think of the “topology” (graph) you create among the processes for the failure detection, in addition to the ring backbone. Think about how the topology is reorganized when you have failures. You can fix the pinging topology, i.e., who pings whom. **Use the minimum number of monitors to achieve the MP specification.** (Higher numbers will result in loss of points). In other words, do not do random pinging and do not ping everyone (remember, you are told at most 3 machines fail simultaneously – use that to your advantage). In your report, justify your design choices and why your design: a) meets 5 second completeness, and b) meets completeness up to 3 failures, and c) would violate completeness with at least one combination of 4 simultaneous failures.

Your algorithm must use a small bandwidth (defined as Bytes per second, and NOT as messages per second) needed to meet the above requirements. So, for instance, don’t use all to all heartbeating or all-to-all or random pinging (it’s an overkill, and if you do it, FakeNews Inc. will slander you by generating fake news about you, and giving you fake points.).

For the failure detection or leaves, you cannot use a master or leader, since its failure must be detected as well. However, to enable machines to join the group, you can have a fixed contact machine that all potential members know about, which you already know is called the “introducer”. When the introducer is down, no new members can join the group until the contact has rejoined – but the rest of the group should proceed normally including failures should still being detected, and leaves being allowed.

Pay attention to the format of messages that are sent between machines. Ensure that any platform-dependent fields (e.g., ints) are marshaled (converted) into a platform-independent format. An example is Google’s Protocol Buffers (this is not a requirement, especially since it is not installed on CS VM Cluster). You can invent your own, but do specify it clearly in your report.
Make your implementation bandwidth-efficient. Your implementation must use UDP (cheap).

Create logs at each machine, and if possible, use MP1 to debug. These logs are important as we will be asking you to grep them at demo time. You can make your logs as verbose as you want them, but at the least you must log: 1) each time a change is made to the local membership list (join, leave or failure) and 2) each time a failure is detected or communicated from one machine to another. We will request to see the log entries at demo time, via the MP1’s querier. Thus, make sure you integrate MP2 with MP1 to make this possible. You should also use your MP1 solution for debugging MP2 (and mention how useful this was in the report). (If this is not possible for MP2 because you didn’t finish MP1, then you can cook up a simple grep engine for MP2. But for future MP3 and onwards, make sure MP1 is integrated.)

We also recommend (but don’t require) writing unit tests for each of the join, leave, and failure functionalities. At the least, ensure that these actually work for a long series of join/leave/fail events.

**Machines:** We will be using the CS VM Cluster machines. You will be using 7-10 VMs for the demo. The VMs do not have persistent storage, so you are required to use git to manage your code. To access git from the VMs, use the same instructions as MP1.

**Demo:** Demos are usually scheduled on the Monday right after the MP is due. The demos will be on the CS VM Cluster machines. You must use all VMs for your demo (details will be posted on Piazza closer to the demo date). Please make sure your code runs on the CS VM Cluster machines, especially if you’ve used your own machines/laptops to do most of your coding. Please make sure that any third party code you use is installable on CS VM Cluster. Further demo details and a signup sheet will be made available closer to the date. We expect both partners to contribute equivalent amounts of effort during the entire MP execution (not just in the demo).

**Language:** Choose your favorite language! We recommend C/C++/Java/Go/Rust. We will release “Best MPs” from the class in these languages only (so you can use them in subsequent MPs).

**Report:** Write a report of less than 2 pages (12 pt font, typed only - no handwritten reports please!). Briefly describe your design (algorithm used, why it scales to large N, and marshaled message format, why it satisfies the three completeness requirements listed earlier), very briefly how useful MP1 was for debugging MP2, and measurements (not just mathematical calculations) of: (i) the background bandwidth usage (in Bps, not messages per second) for N=6
machines (assuming no membership changes), (ii) the average bandwidth usage whenever a node joins, leaves or fails (3 different numbers) for N=6 machines, and (iii) plot the false positive rate of your membership service when the message loss rate is 3%, 30% (you can emulate message losses at the sending end by dropping messages before sending them out to the network) – do this last part for a group with N=2, 6 machines (so, 4 total data points). On any plot, for each data point take at least as many readings as is necessary to get a non-zero false positive rate (at least 5 readings each), and plot averages and standard deviations and confidence intervals. Discuss your plots, don’t just put them on paper, i.e., discuss trends, and whether they are what you expect or not (why or why not). (Measurement numbers don’t lie, but we need to make sense of them!) Stay within page limit – for every line over the page limit you will lose 1 point!

Submission: There will be a demo of each group’s project code. On Gradescope submit your report by the deadline (11.59 PM Sunday). In the git, also by the same deadline, submit by the deadline your report as well as working code; please include a README explaining how to compile and run your code. Other submission instructions are similar to previous MPs.

When should I start? Start NOW. You already know all the necessary class material to do this MP. Each MP involves a significant amount of planning, design, and implementation/debugging/experimentation work. Do not leave all the work for the days before the deadline – there will be no extensions.

Evaluation Break-up: Demo [40%], Report (including design and plots) [40%], Code readability and comments [20%].

Academic Integrity: You cannot look at others’ solutions, whether from this year or past years. We will run Moss to check for copying within and outside this class – first offense results in a zero grade on the MP, and second offense results in an F in the course. There are past examples of students penalized in both those ways, so just don’t cheat. You can only discuss the MP spec and lecture concepts with the class students and forum, but not solutions, ideas, or code (if we see you posting code on the forum, that’s a zero on the MP). FakeNews Inc. is watching and will be very Sad!

We recommend you stick with the same group from one MP to the next (this helps keep the VM mapping sane on EngrIT’s end), except for exceptional circumstances. We expect all group members to contribute about equivalently to the overall effort. If you believe your group members are not, please have “the talk” with them first, give them a second chance. If that doesn’t work either, please approach Indy.
Happy Membership (from us and the fictitious FakeNews Inc.)!