CS 425 / ECE 428
Distributed Systems
Fall 2018

Indranil Gupta (Indy)

Lecture 6: Failure Detection and Membership, Grids

A Challenge

- You’ve been put in charge of a datacenter, and your manager has told you, “Oh no! We don’t have any failures in our datacenter!”

- Do you believe him/her?

- What would be your first responsibility?
- Build a failure detector
- What are some things that could go wrong if you didn’t do this?

Failures are the Norm

… not the exception, in datacenters.

Say, the rate of failure of one machine (OS/disk/motherboard/network, etc.) is once every 10 years (120 months) on average.

When you have 120 servers in the DC, the mean time to failure (MTTF) of the next machine is 1 month.

When you have 12,000 servers in the DC, the MTTF is about once every 7.2 hours!

Soft crashes and failures are even more frequent!

To build a failure detector

- You have a few options

1. Hire 1000 people, each to monitor one machine in the datacenter and report to you when it fails.
2. Write a failure detector program (distributed) that automatically detects failures and reports to your workstation.

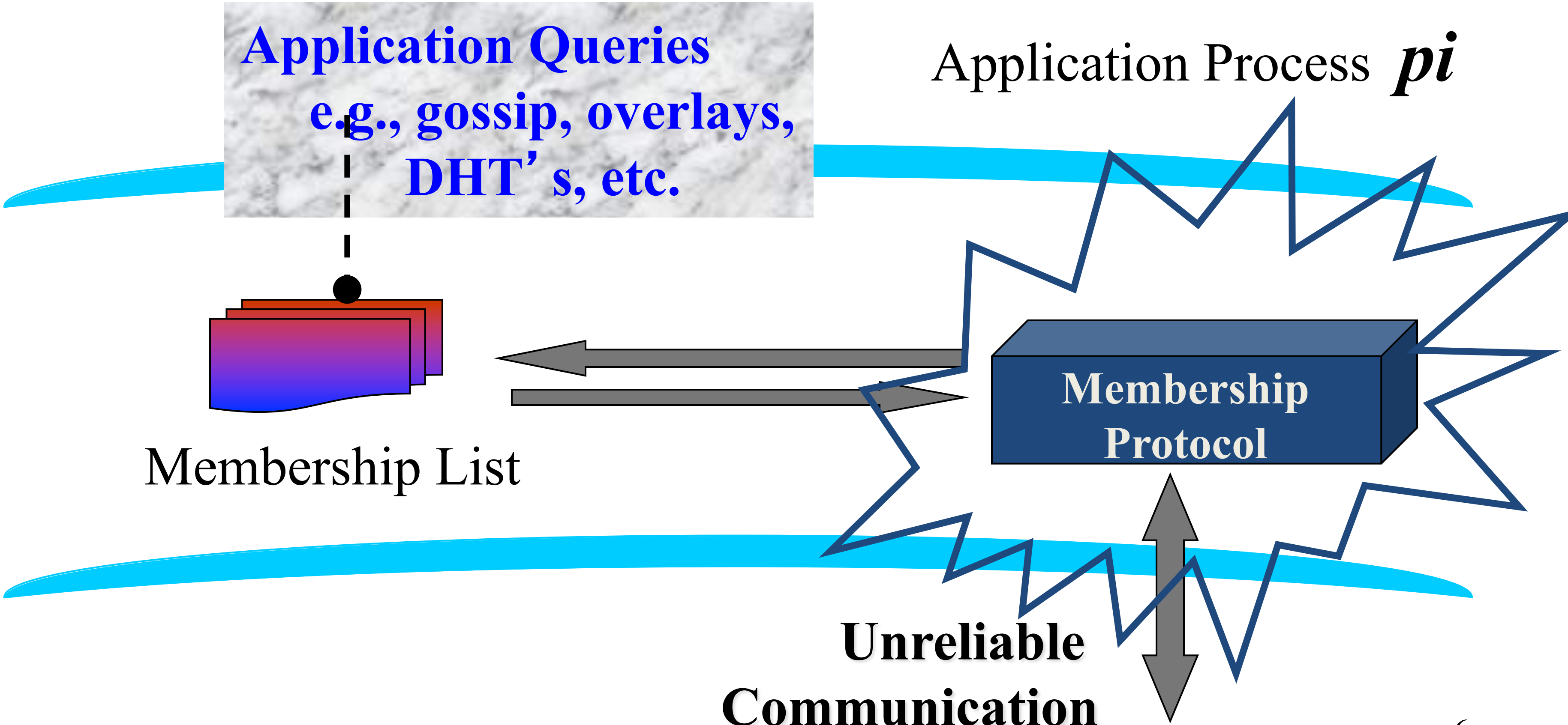
Which is more preferable, and why?

Target Settings

- Process ‘group’-based systems
  - Clouds/Datacenters
  - Replicated servers
  - Distributed databases

- Fail-stop (crash) process failures

Group Membership Service



Application Queries

e.g., gossip, overlays, DHT's, etc.

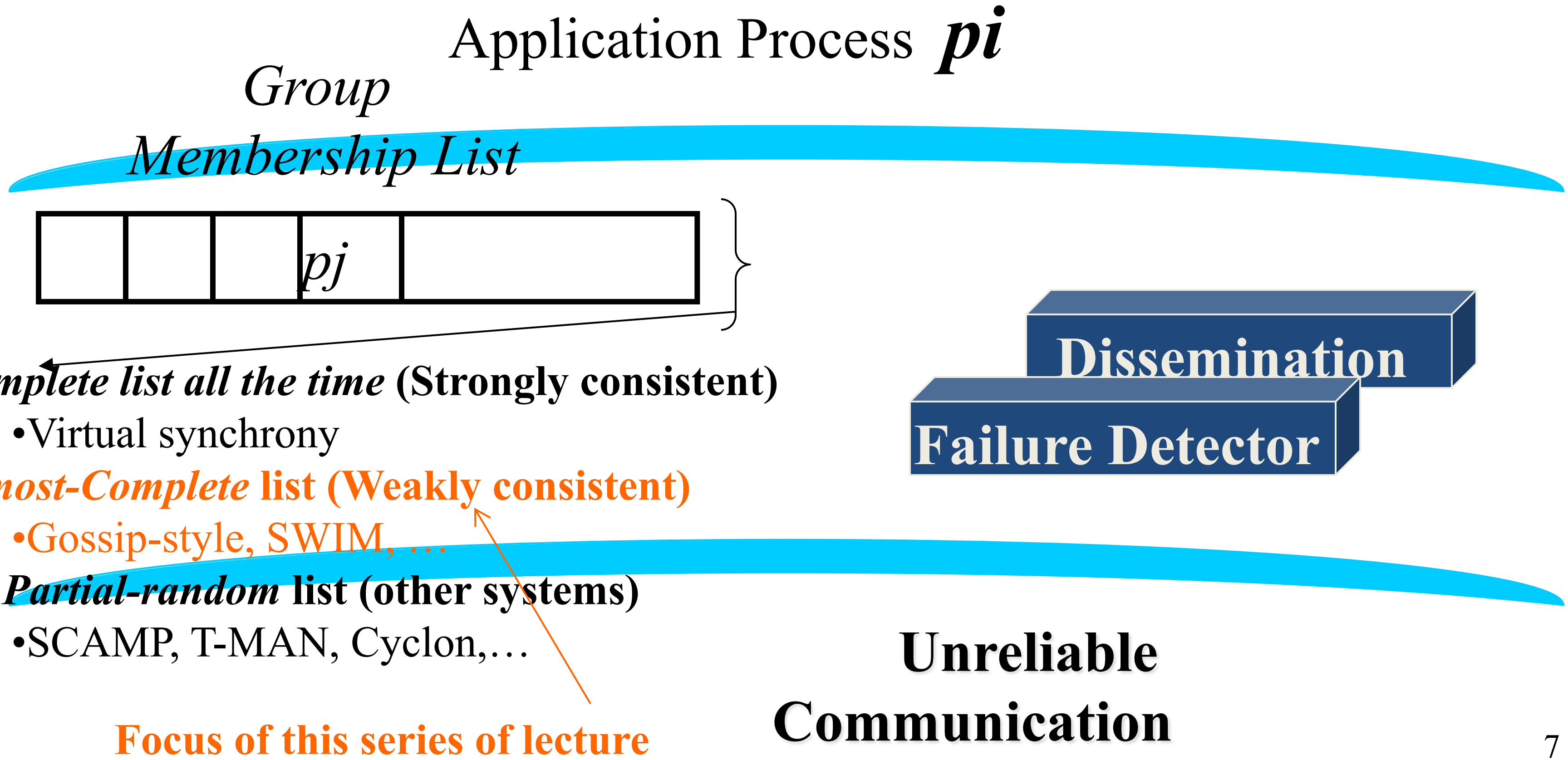
Application Process ***pi***

Membership List

Membership Protocol

Unreliable Communication

Two sub-protocols



Application Process ***pi***

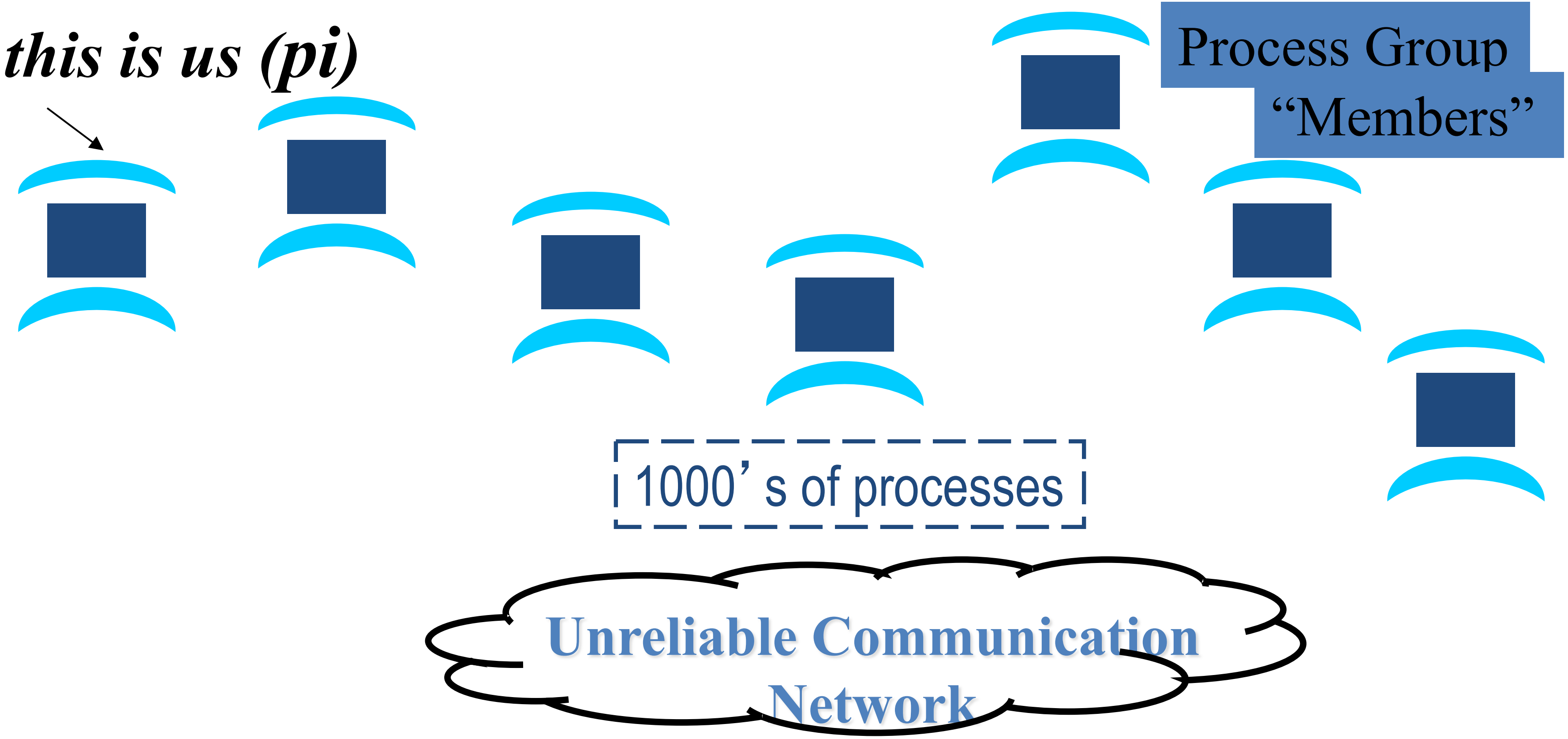
Group Membership List

- ***Complete list all the time*** **(Strongly consistent)**
  - Virtual synchrony
- ***Almost-Complete*** **list (Weakly consistent)**
  - Gossip-style, SWIM, …
- **Or *Partial-random* list (other systems)**
  - SCAMP, T-MAN, Cyclon,…

Focus of this series of lecture

Unreliable Communication

Large Group: Scalability A Goal



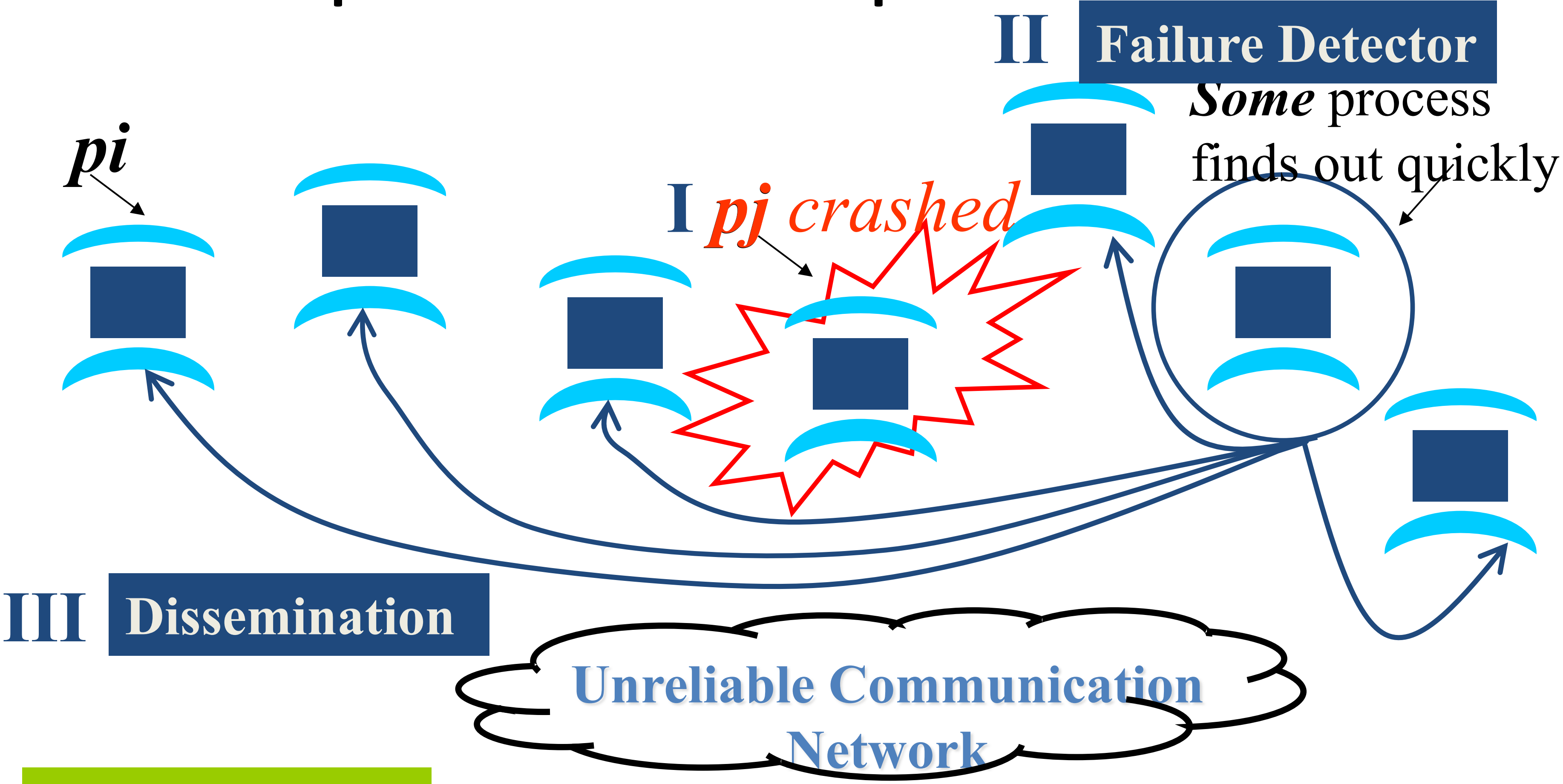
this is us (pi)

Process Group “Members”

1000’s of processes

Unreliable Communication Network

Group Membership Protocol



pi

I ***pj crashed***

II **Failure Detector**

Some process finds out quickly

III **Dissemination**

Unreliable Communication Network

Fail-stop Failures only

Next

- How do you design a group membership protocol?

I. *pj* crashes

- Nothing we can do about it!
- A frequent occurrence
- Common case rather than exception
- Frequency goes up linearly with size of datacenter

II. Distributed Failure Detectors: Desirable Properties

- Completeness = each failure is detected
- Accuracy = there is no mistaken detection
- Speed
  - Time to first detection of a failure
- Scale
  - Equal Load on each member
  - Network Message Load

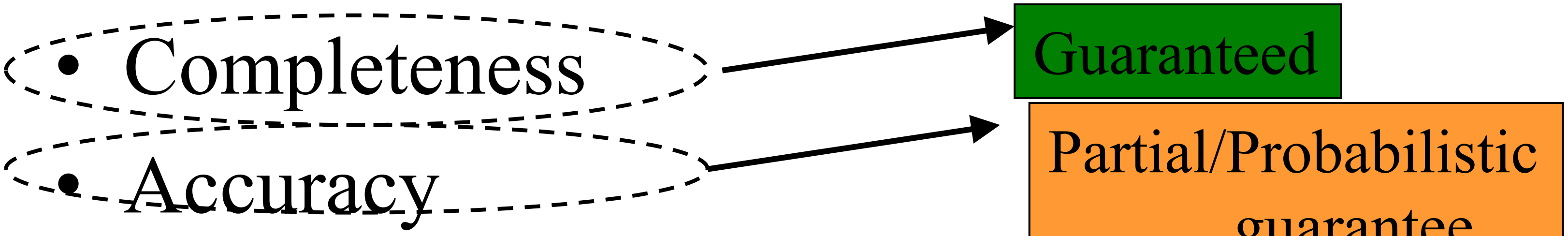
Distributed Failure Detectors: Properties

- Completeness
- Accuracy
- Speed
  - Time to first detection of a failur
- Scale
  - Equal Load on each member
  - Network Message Load

Impossible together in lossy networks [Chandra and Toueg]

If possible, then can solve consensus! (but consensus is known to be unsolvable in asynchronous systems)

What Real Failure Detectors Prefer



- Completeness
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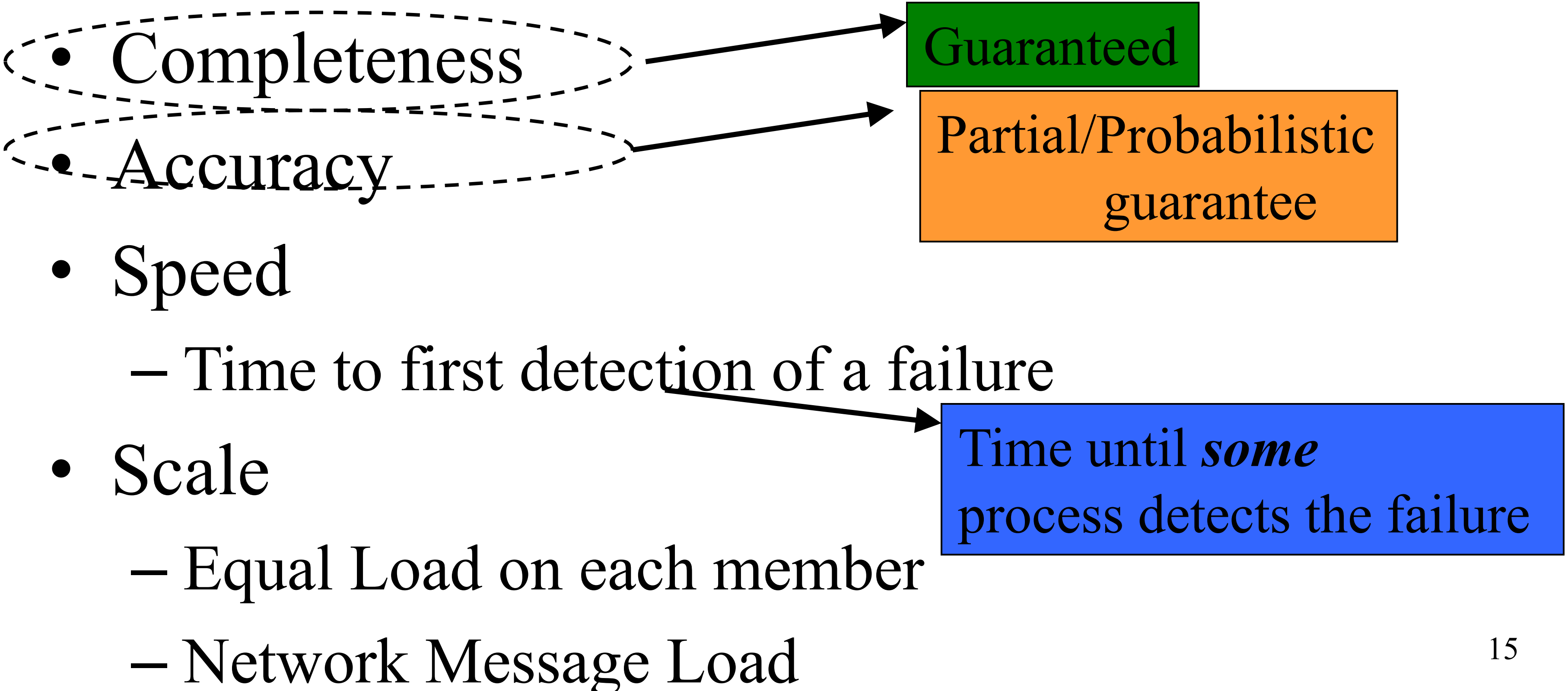
What Real Failure Detectors Prefer



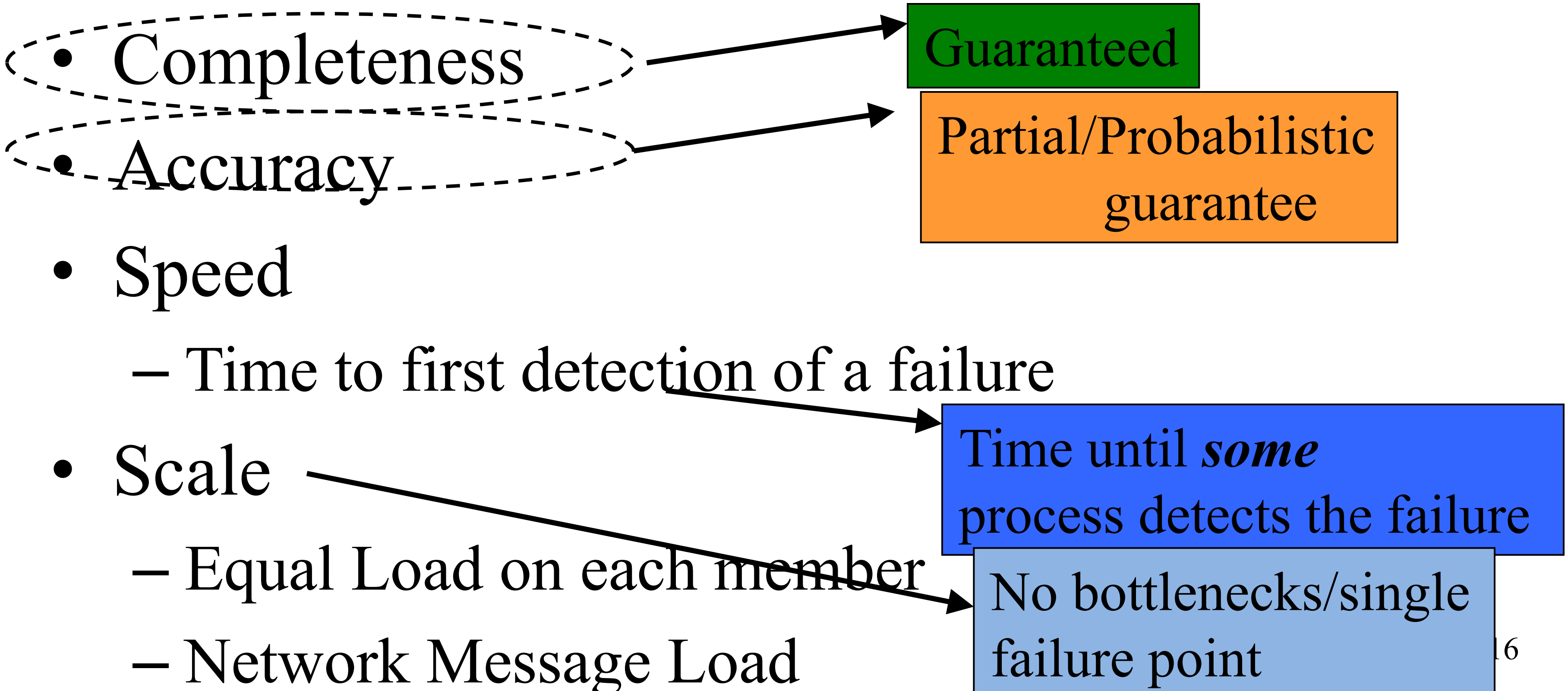
- Completeness
- Accuracy
- Speed
  - – Time to first detection of a failure
- Scale
  - – Equal Load on each member
  - – Network Message Load

Guaranteed

Partial/Probabilistic guarantee

Time until ***some*** process detects the failure

What Real Failure Detectors Prefer



- Completeness
- Accuracy
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  - Time to first detection of a failure
- Scale
  - Equal Load on each member
  - Network Message Load

Guaranteed

Partial/Probabilistic guarantee

Time until ***some*** process detects the failure

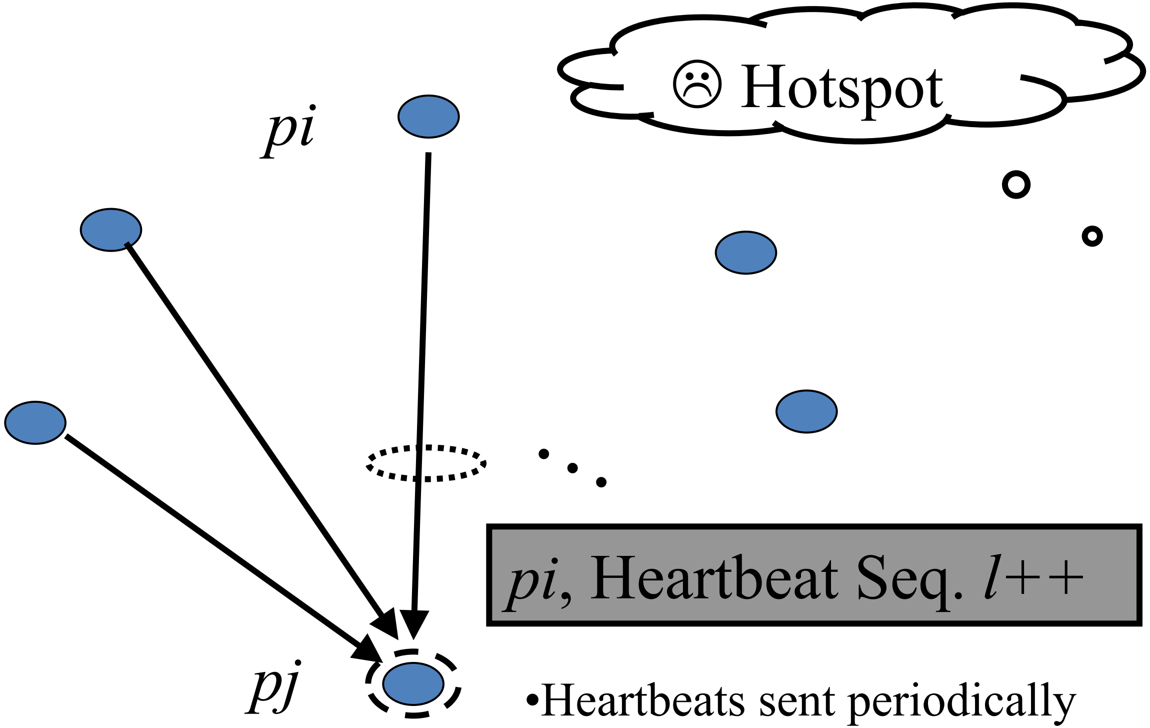
No bottlenecks/single failure point

Failure Detector Properties

- Completeness
- Accuracy
- Speed
  - Time to first detection of a failure
- Scale
  - Equal Load on each member
  - Network Message Load

In spite of arbitrary simultaneous process failures

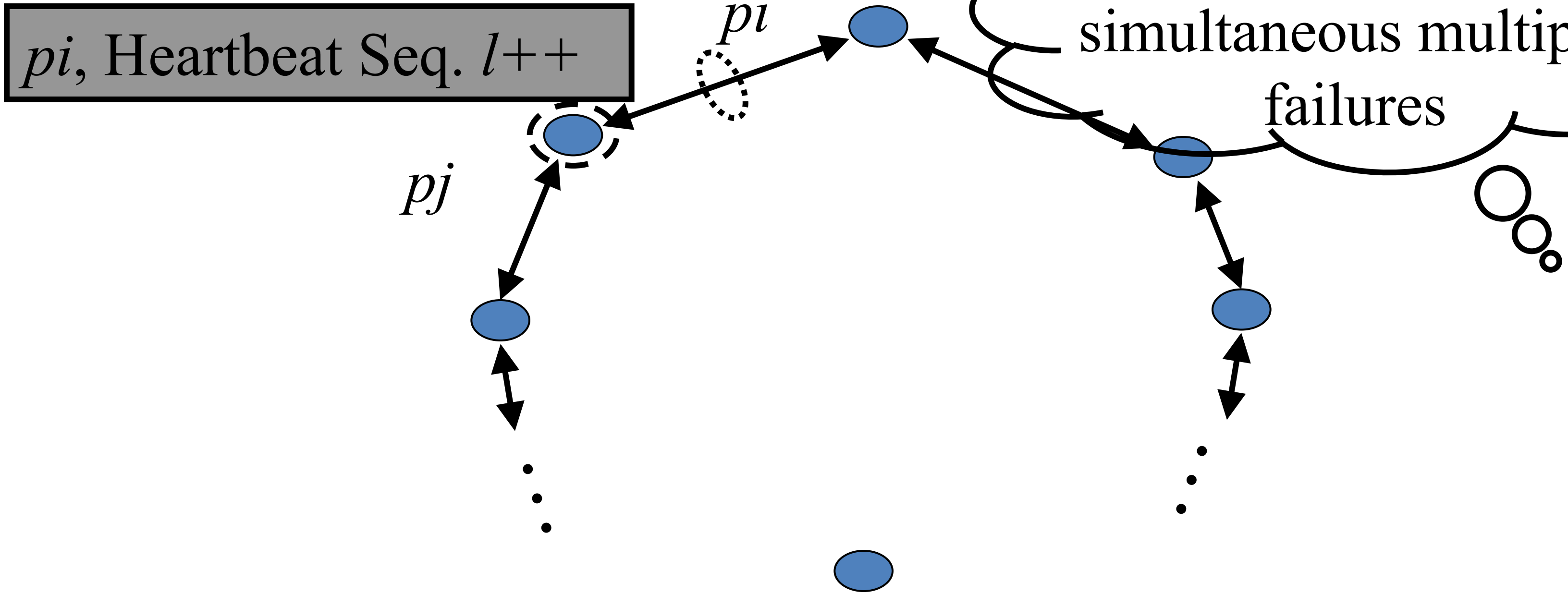
Centralized Heartbeating



pi, Heartbeat Seq. *l*++

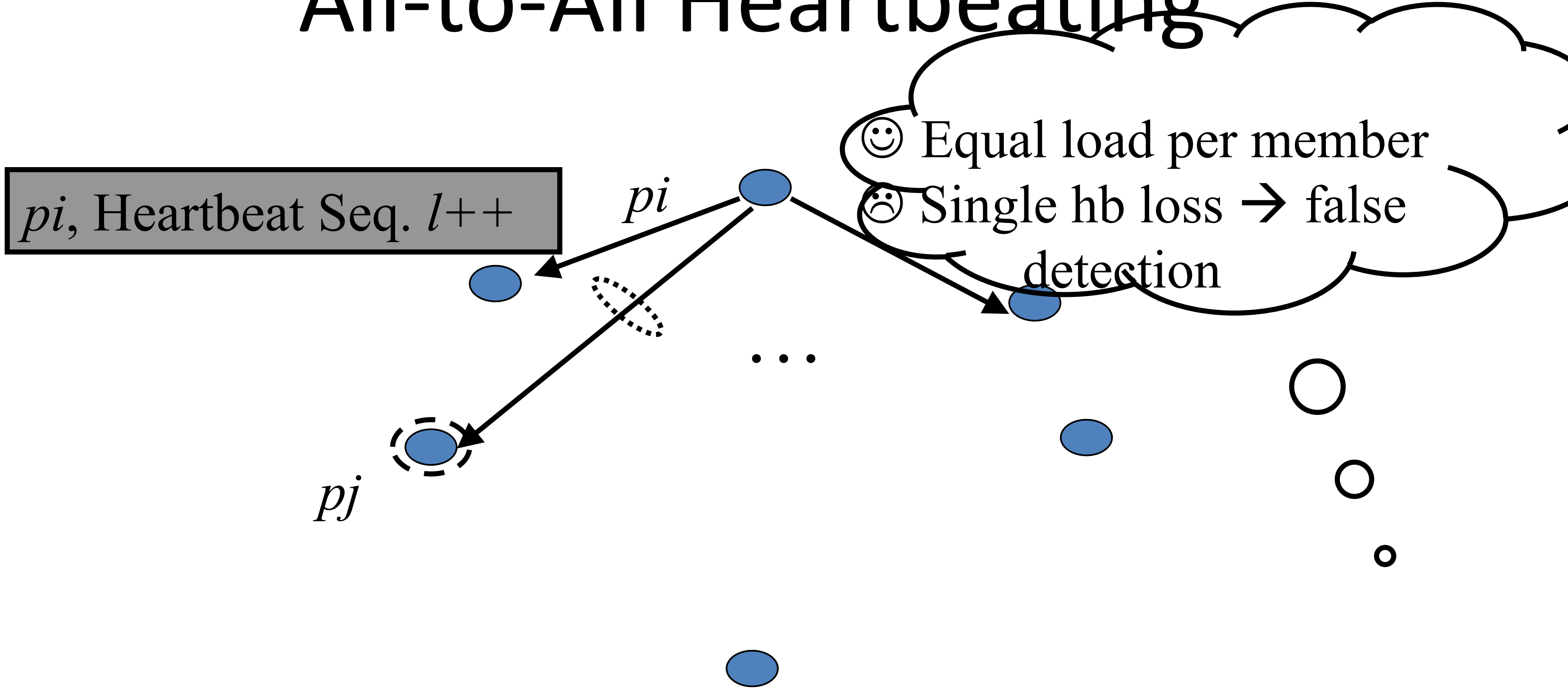
- Heartbeats sent periodically
- If heartbeat not received from *pi* within timeout, mark *pi* as failed

Ring Heartbeating



☹ Unpredictable on simultaneous multiple failures

All-to-All Heartbeating



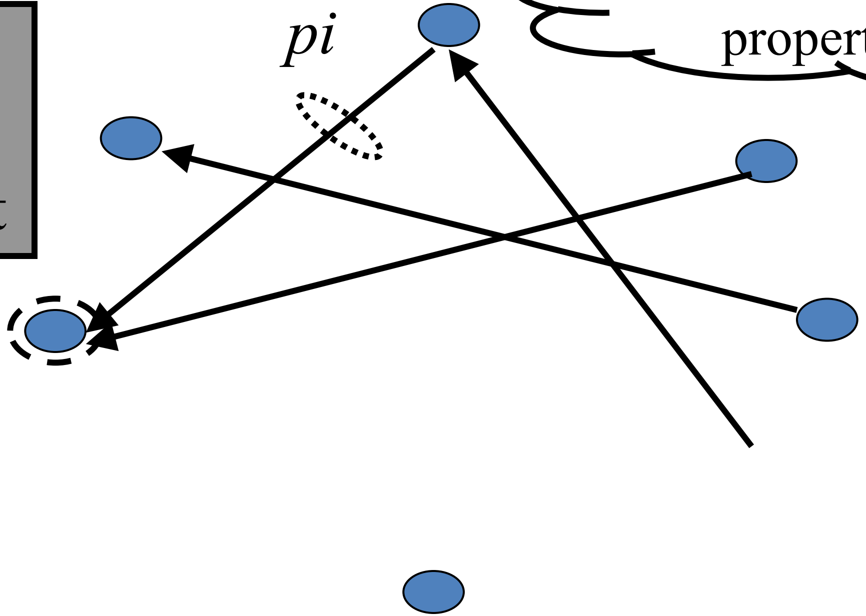
- ☺ Equal load per member
- ☹ Single hb loss → false detection

Next

- How do we increase the robustness of all-to-all heartbeating?

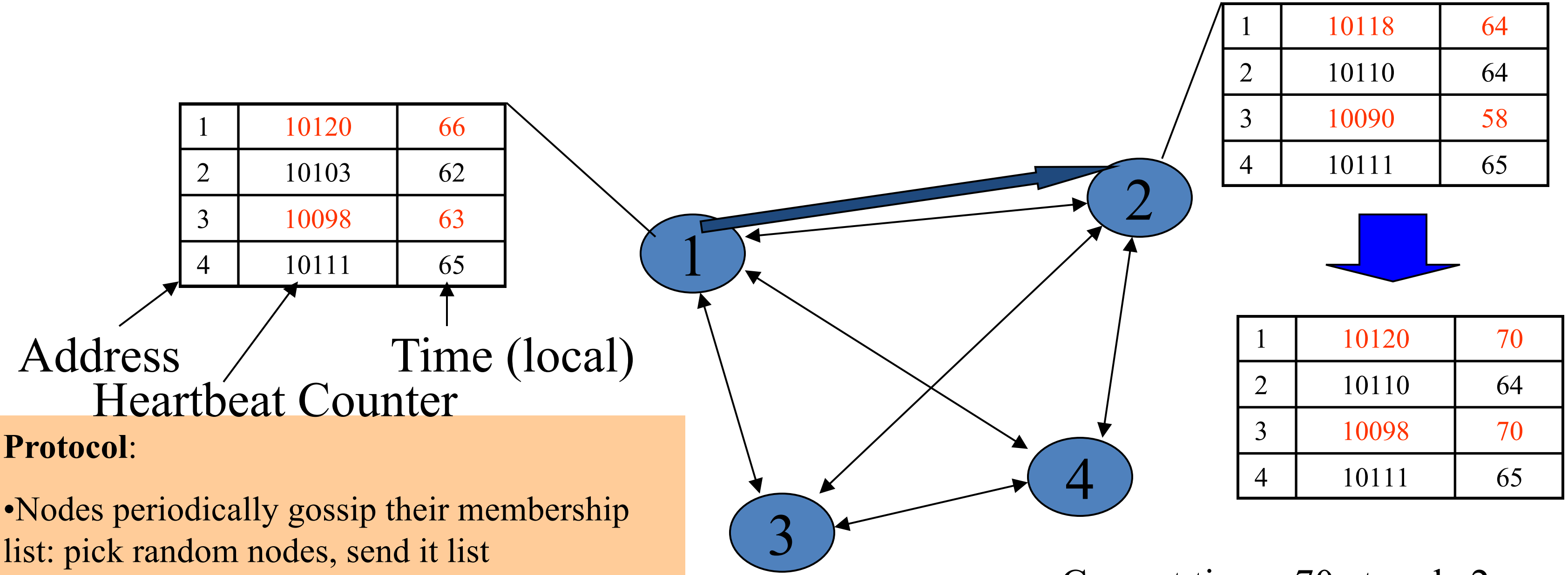
Gossip-style Heartbeating

Array of
Heartbeat Seq. *l*
for member subset



☺ Good accuracy properties

Gossip-Style Failure Detection



1	10120	66
2	10103	62
3	10098	63
4	10111	65

Address

Heartbeat Counter

Time (local)

1	10118	64
2	10110	64
3	10090	58
4	10111	65

1	10120	70
2	10110	64
3	10098	70
4	10111	65

Protocol:

- Nodes periodically gossip their membership list: pick random nodes, send it list
- On receipt, it is *merged* with local membership list
- When an entry times out, member is marked as failed

Current time : 70 at node 2

(asynchronous clocks)

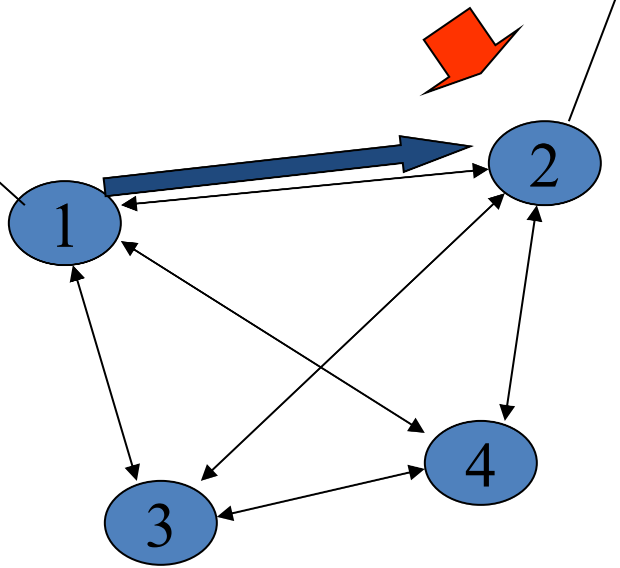
Gossip-Style Failure Detection

- If the heartbeat has not increased for more than T_{fail} seconds,
  the member is considered failed
- And after a further $T_{cleanup}$ seconds, it will delete the member from the list
- Why an additional timeout? Why not delete right away?

Gossip-Style Failure Detection

- What if an entry pointing to a failed node is deleted right after T_{fail} (=24) seconds?

1	10120	66
2	10103	62
3	10098	55
4	10111	65



1	10120	66
2	10110	64
3 4	10098 10111	75 65
4	10111	65

Current time : 75 at node 2

Analysis/Discussion

- Well-known result: a gossip takes O(log(N)) time to propagate.
- So: Given sufficient bandwidth, a single heartbeat takes O(log(N)) time to propagate.
- So: N heartbeats take:
  - O(log(N)) time to propagate, if bandwidth allowed per node is allowed to be O(N)
  - O(N.log(N)) time to propagate, if bandwidth allowed per node is only O(1)
  - What about O(k) bandwidth?
- What happens if gossip period T_{gossip} is decreased?
- What happens to $P_{mistake}$ (false positive rate) as T_{fail} ,$T_{cleanup}$ is increased?
- Tradeoff: False positive rate vs. detection time vs. bandwidth

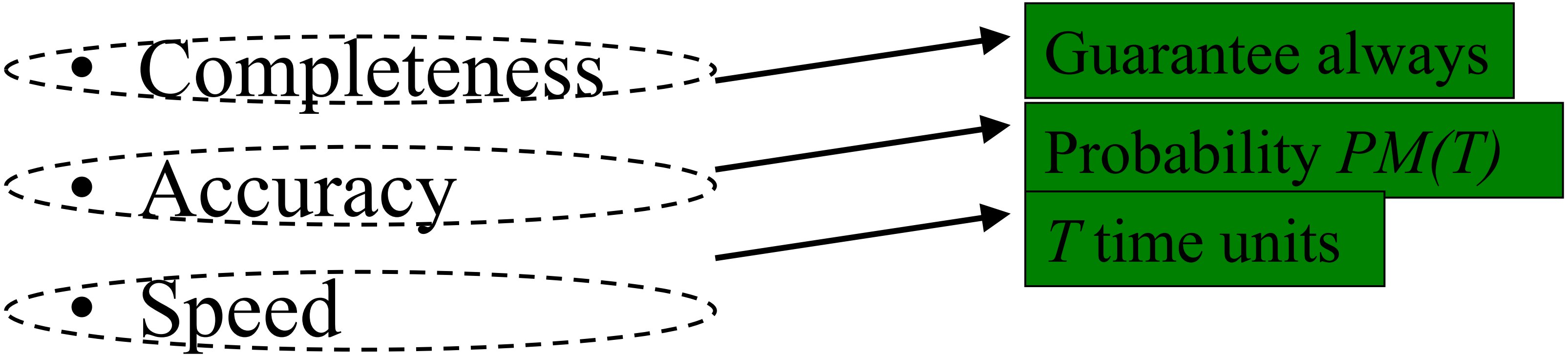
Next

- So, is this the best we can do? What is the best we can do?

Failure Detector Properties …

- Completeness
- Accuracy
- Speed
  - Time to first detection of a failure
- Scale
  - Equal Load on each member
  - Network Message Load

re application-defined Requirements



- Completeness
- Accuracy
- Speed
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- Scale
  - Equal Load on each member
  - Network Message Load

Guarantee always

Probability *PM(T)*

T time units

Are application-defined Requirements

- Completeness
- Accuracy
- Speed
  - Time to first detection of a failure
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  - Network Message Load

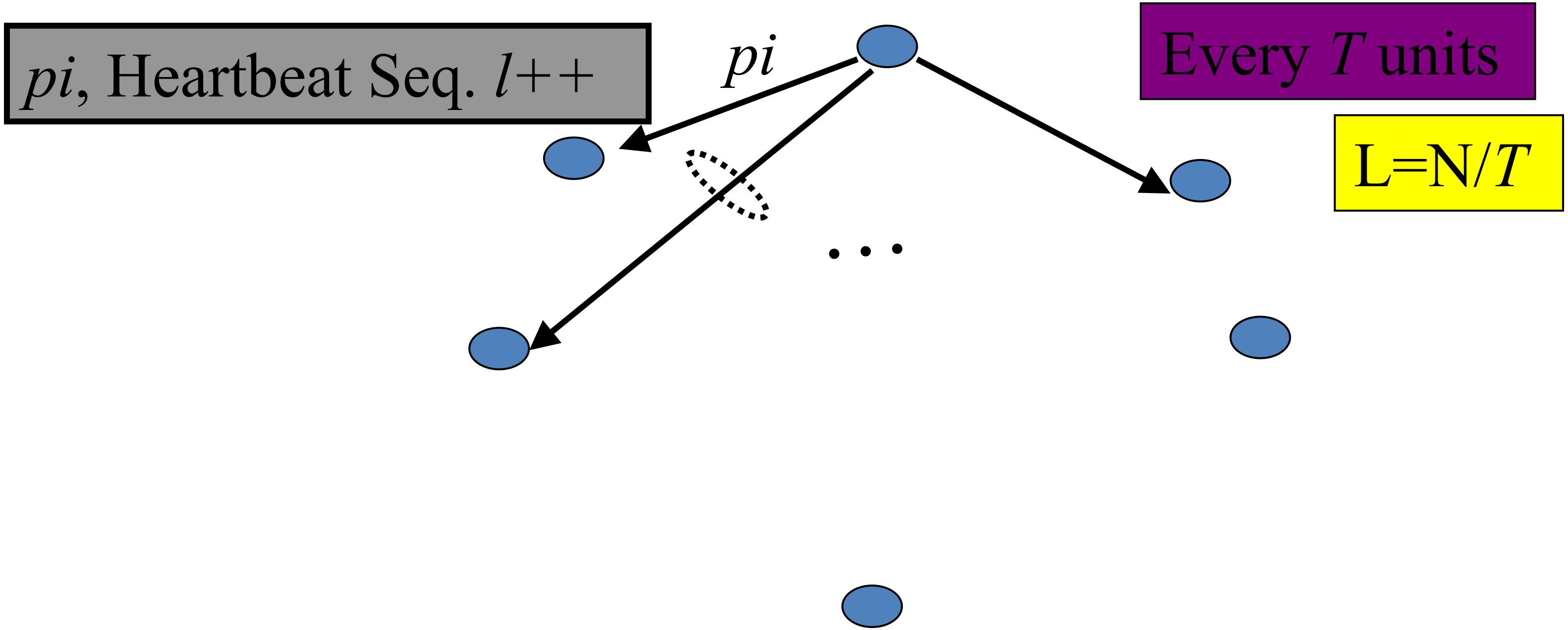
Guarantee always

Probability *PM(T)*

T time units

N*L: Compare this across protocols

All-to-All Heartbeating



pi, Heartbeat Seq. *l*++

pi

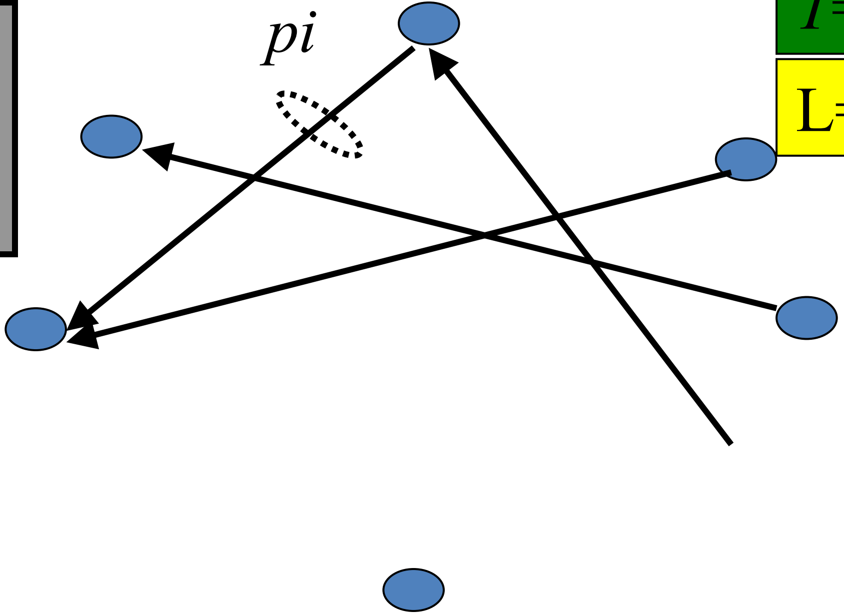
Every *T* units

L=N/*T*

Gossip-style Heartbeating

Array of Heartbeat Seq. *l* for member subset

Every tg units =gossip period, send O(N) gossip message



T=logN * tg

L=N/tg=N*logN/*T*

What's the Best/Optimal we can do?

- *Worst case* load L* per member in the group (messages per second)
  - as a function of *T*, *PM(T)*, N
  - Independent Message Loss probability p_{ml}
- $$\boxed{L^* = \frac{\log(PM(T))}{\log(p_{ml})} \cdot \frac{1}{T}}$$

Heartbeating

- Optimal L is independent of N (!)
- All-to-all and gossip-based: sub-optimal
    - L=O(N/T)
    - try to achieve simultaneous detection at ***all*** processes
    - fail to distinguish *Failure Detection* and *Dissemination* components

➲Can we reach this bound?

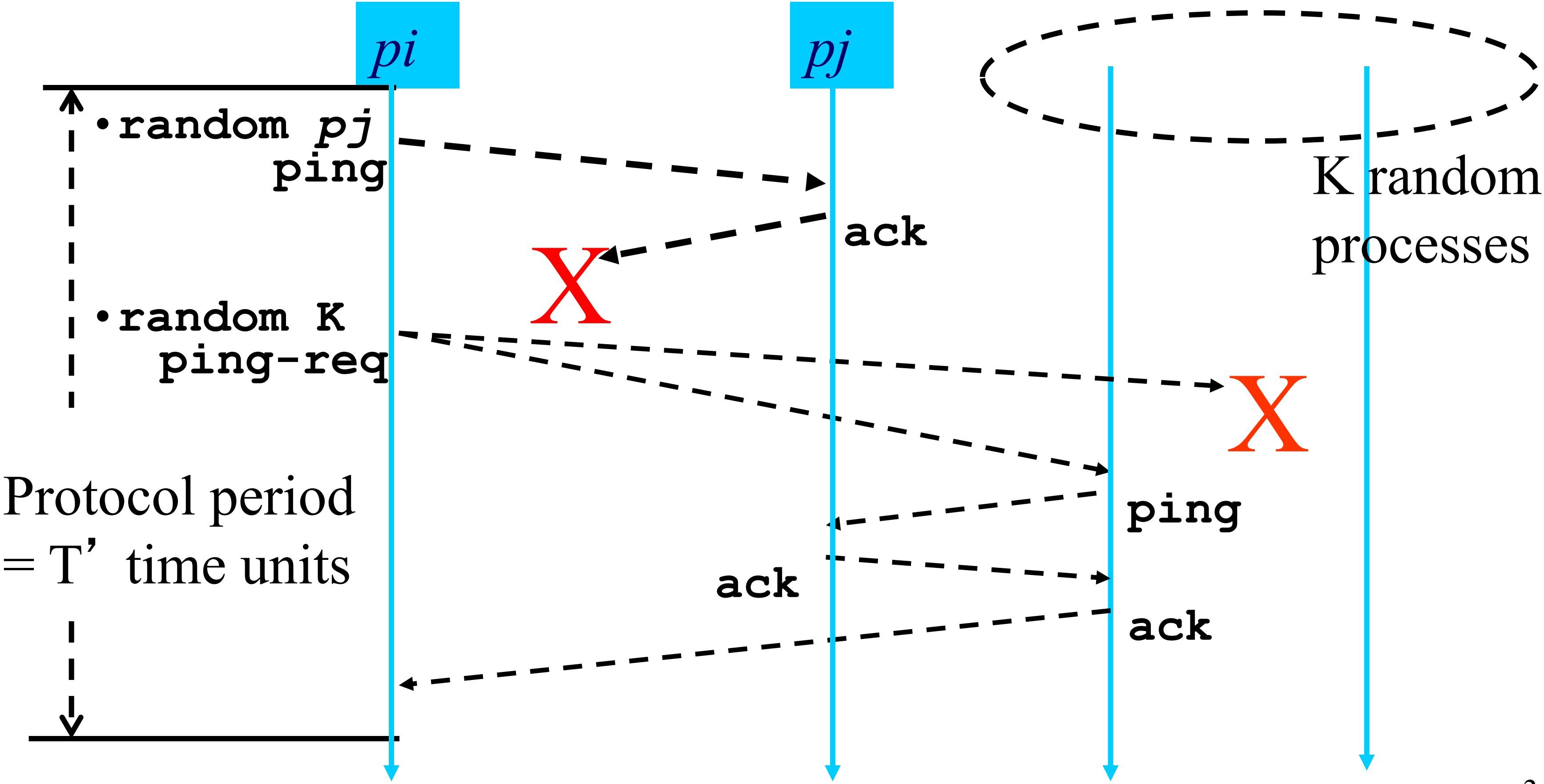
➲Key:

- Separate the two components
- Use a non heartbeat-based Failure Detection Component

Next

- Is there a better failure detector?

SWIM Failure Detector Protocol



Detection Time

- Prob. of being pinged in T'= $1-(1-\frac{1}{N})^{N-1}=1-e^{-1}$
- $\mathrm{E}[T] = \mathrm{T}'.\frac{e}{e-1}$
- Completeness: *Any* alive member detects failure
  - Eventually
  - By using a trick: within worst case *O(N)* protocol periods

Accuracy, Load

- *PM(T)* is exponential in -*K*. Also depends on *pml* (and *pf*)
  - See paper

- $\boxed{\frac{L}{L*} < 28}$ $\boxed{\frac{E[L]}{L*} < 8}$ for up to 15 % loss rates

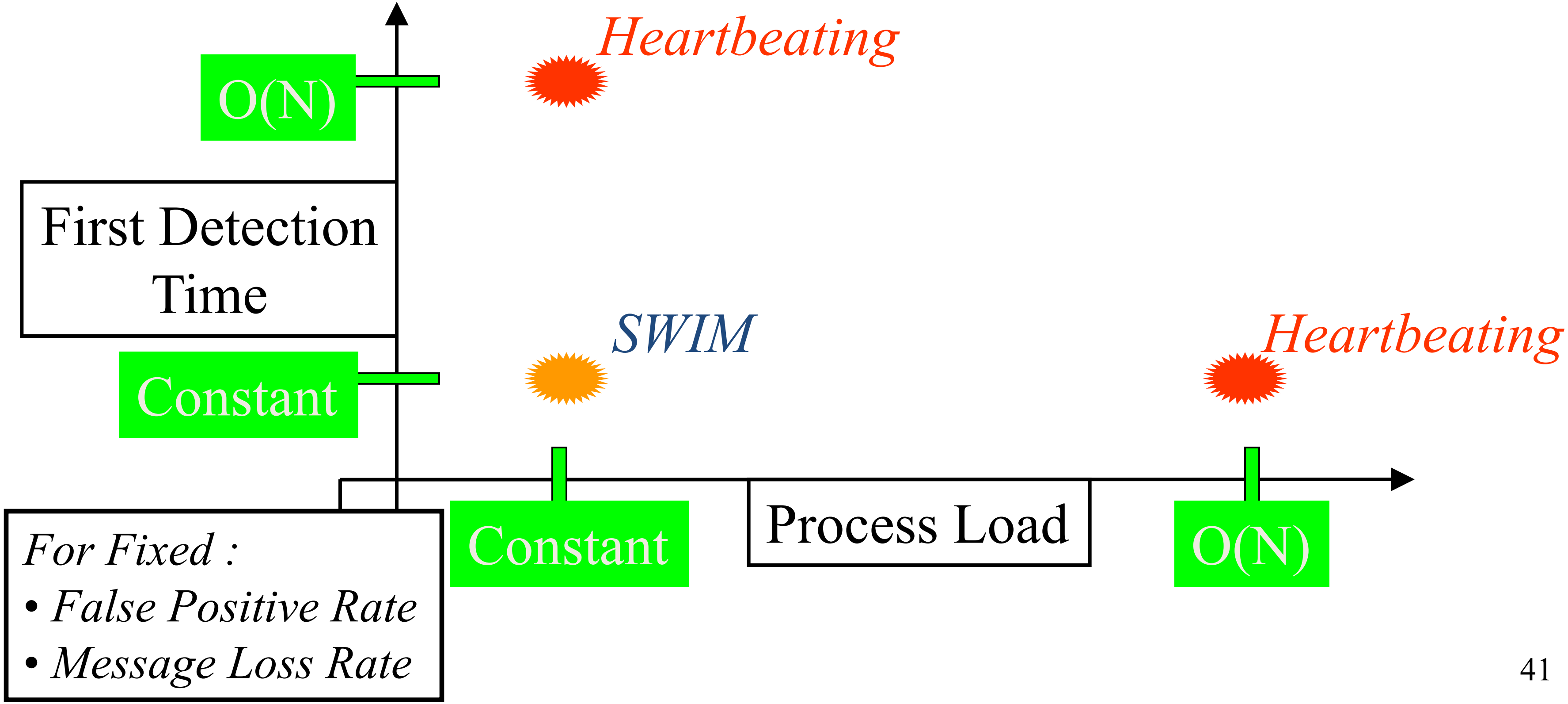
SWIM Failure Detector

Parameter	**SWIM**
First Detection Time	• Expected $\left[\frac{e}{e-1}\right]$ periods • Constant (independent of group size)
Process Load	• Constant per period • < 8 L* for 15% loss
False Positive Rate	• Tunable (via K) • Falls exponentially as load is scaled
Completeness	• Deterministic time-bounded • Within O(log(N)) periods w.h.p.

Time-bounded Completeness

- Key: select each membership element once as a ping target in a traversal
  - Round-robin pinging
  - Random permutation of list after each traversal
- Each failure is detected in worst case 2N-1 (local) protocol periods
- Preserves FD properties

SWIM versus Heartbeating

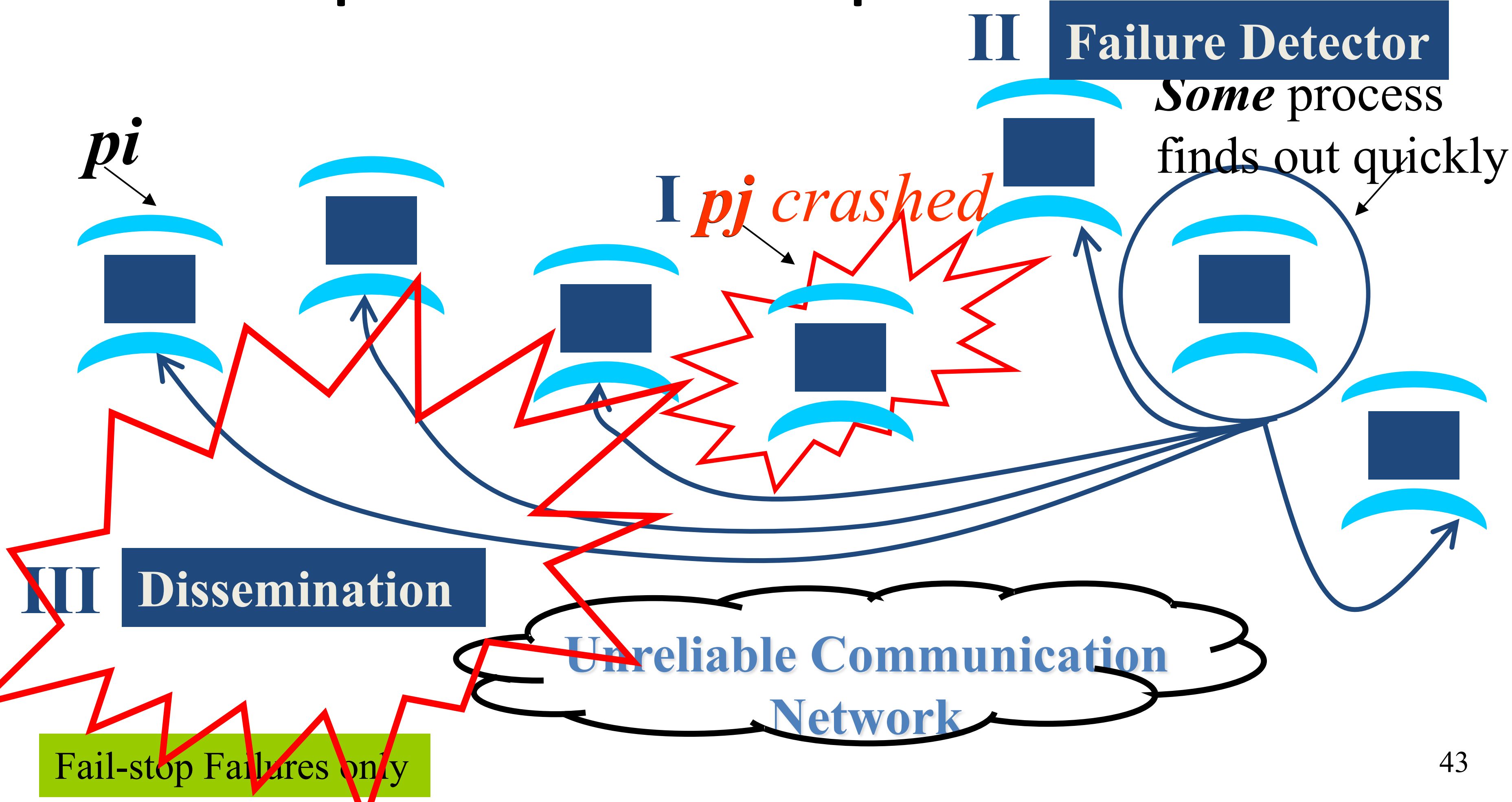


For Fixed :
- *False Positive Rate*
- *Message Loss Rate*

Next

- How do failure detectors fit into the big picture of a group membership protocol?
- What are the missing blocks?

Group Membership Protocol



II **Failure Detector**

Some process finds out quickly

pi

I ***pj crashed***

III **Dissemination**

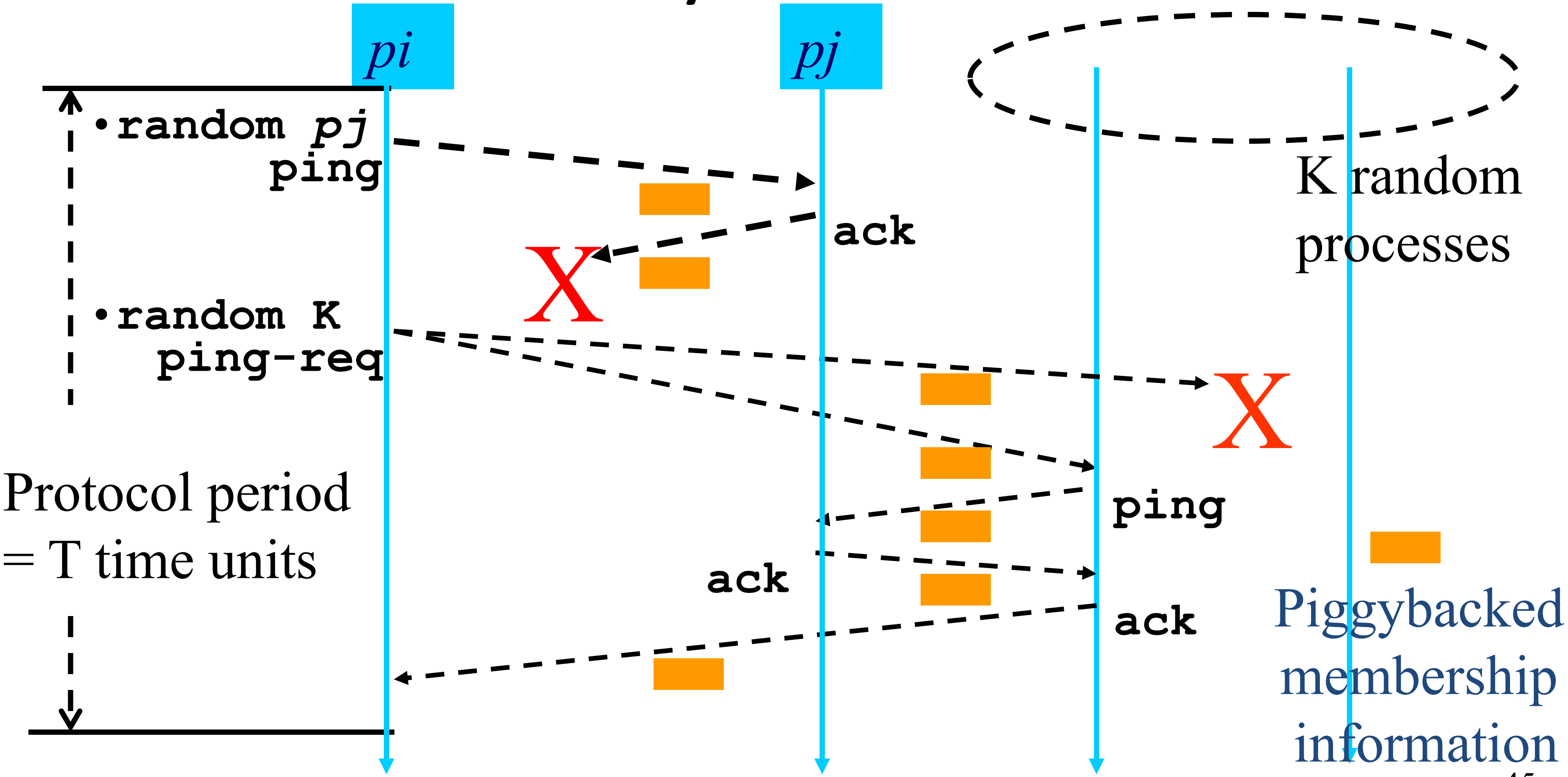
Unreliable Communication Network

Fail-stop Failures only

Dissemination Options

- Multicast (Hardware / IP)
  - unreliable
  - multiple simultaneous multicasts
- Point-to-point (TCP / UDP)
  - expensive
- Zero extra messages: Piggyback on Failure Detector messages
  - Infection-style Dissemination

Infection-style Dissemination



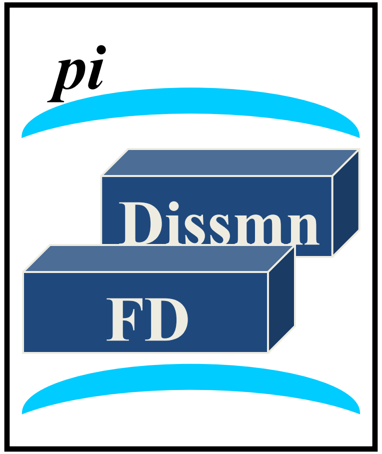
Infection-style Dissemination

- Epidemic/Gossip style dissemination
  - After $\lambda.\log(N)$ protocol periods, $N^{-(2\lambda-2)}$ processes would not have heard about an update
- Maintain a buffer of recently joined/evicted processes
  - Piggyback from this buffer
  - Prefer recent updates
- Buffer elements are garbage collected after a while
  - After $\lambda.\log(N)$ protocol periods, i.e., once they've propagated through the system; this defines weak consistency

Suspicion Mechanism

- False detections, due to
  - Perturbed processes
  - Packet losses, e.g., from congestion
- Indirect pinging may not solve the problem
- Key: *suspect* a process before *declaring* it as failed in the group

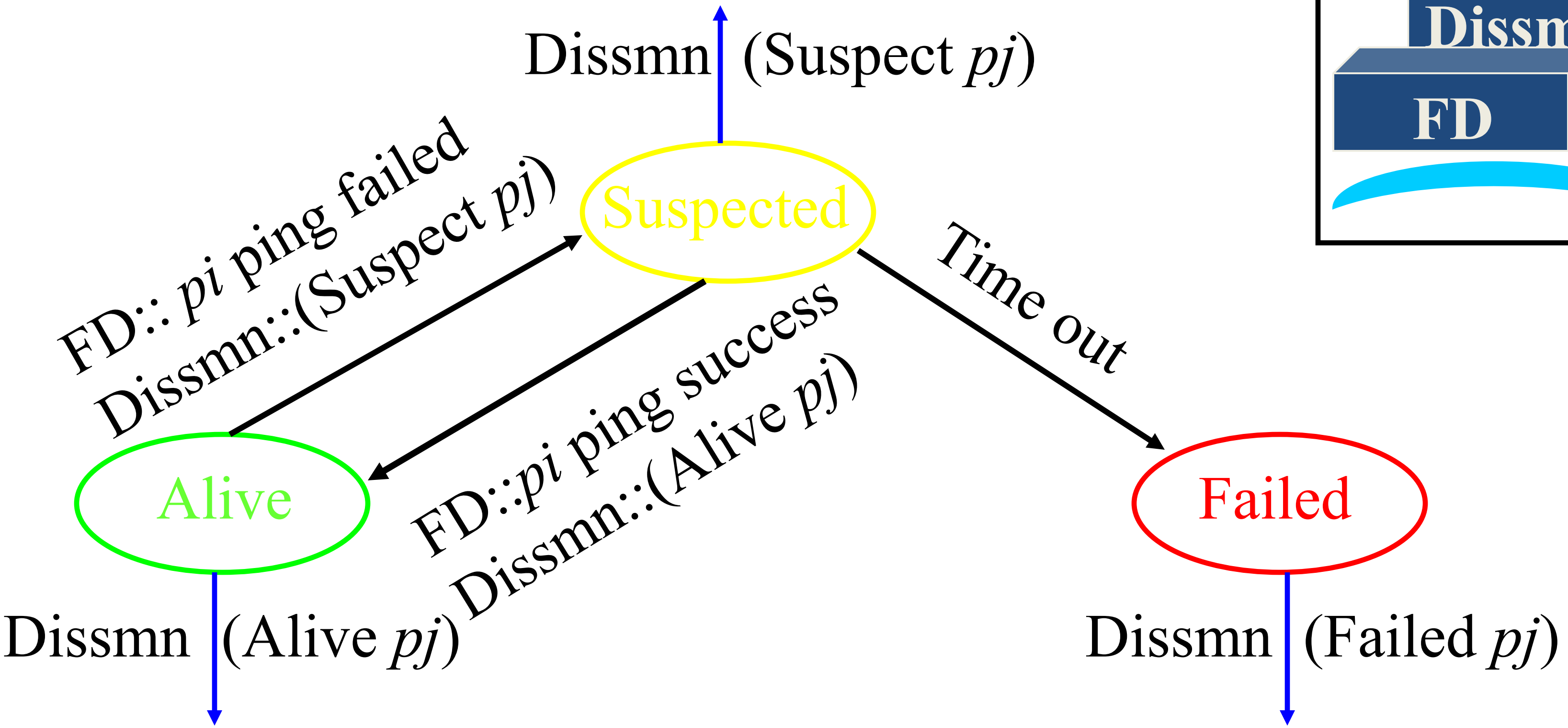
Suspicion Mechanism



pi

Dissmn

FD



Dissmn (Suspect *pj*)

FD:: *pi* ping failed
Dissmn::(Suspect *pj*)

Suspected

Time out

FD::*pi* ping success
Dissmn::(Alive *pj*)

Alive

Failed

Dissmn (Alive *pj*)

Dissmn (Failed *pj*)

Suspicion Mechanism

- Distinguish multiple suspicions of a process
  - Per-process *incarnation number*
  - *Inc* # for *pi* can be incremented only by *pi*
    - e.g., when it receives a (Suspect, *pi*) message
  - Somewhat similar to DSDV (routing protocol in ad-hoc nets)
- Higher inc# notifications over-ride lower inc#'s
- Within an inc#: (Suspect inc #) > (Alive, inc #)
- (Failed, inc #) overrides everything else

SWIM In Industry

- First used in Oasis/CoralCDN
- Implemented open-source by Hashicorp Inc.
  - Called “Serf”
  - Later “Consul”
- Today: Uber implemented it, uses it for failure detection in their infrastructure
  - See “ringpop” system

Wrap Up

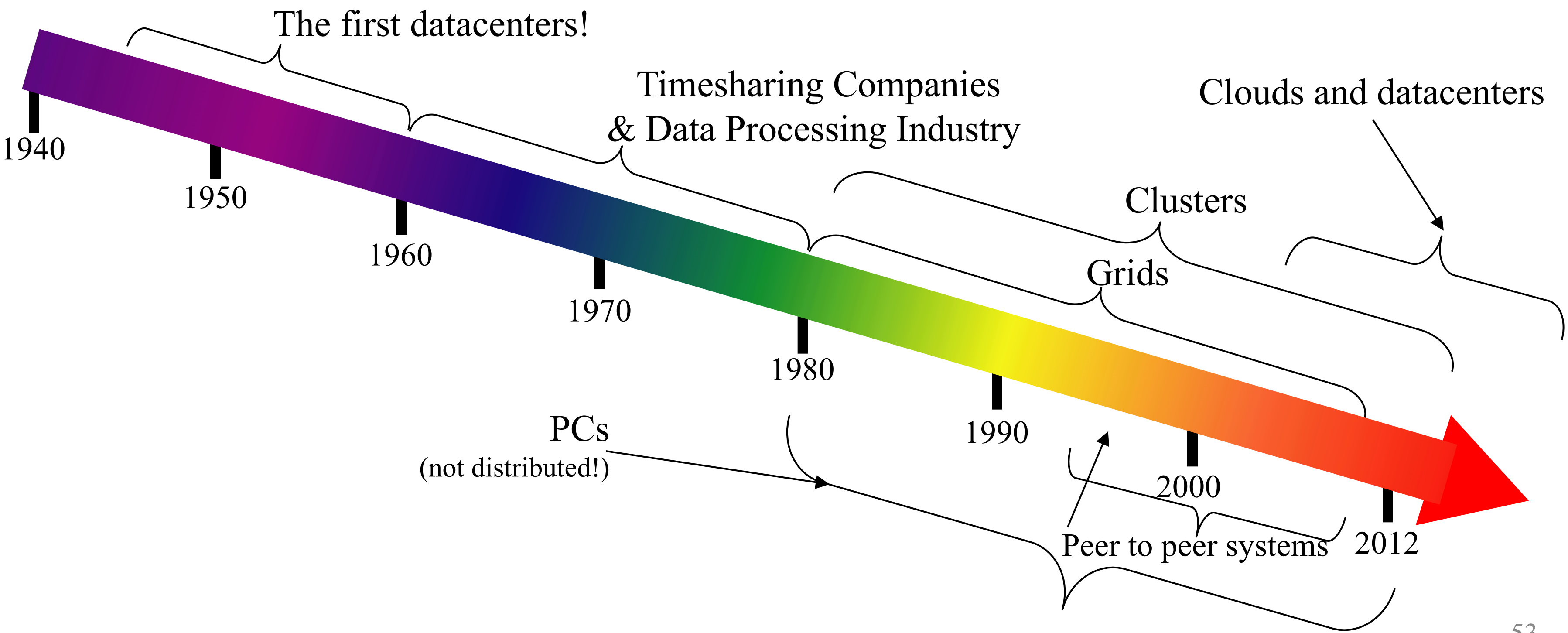
- Failures the norm, not the exception in datacenters
- Every distributed system uses a failure detector
- Many distributed systems use a membership service

- Ring failure detection underlies
  - IBM SP2 and many other similar clusters/machines

- Gossip-style failure detection underlies
  - Amazon EC2/S3 (rumored!)

Grid Computing

“A Cloudy History of Time”



"A Cloudy History of Time"

First large datacenters: ENIAC, ORDVAC, ILLIAC
Many used vacuum tubes and mechanical relays

1940

1950

1960

1970

1980

1990

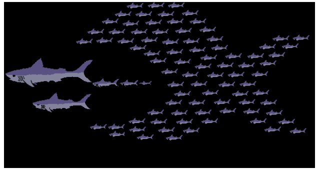
2000

2012

Clouds

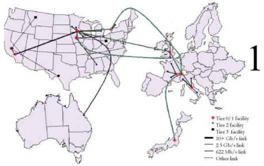
Data Processing Industry
- 1968: $70 M. 1978: $3.15 Billion

Timesharing Industry (1975):
- Market Share: Honeywell 34%, IBM 15%,
- Xerox 10%, CDC 10%, DEC 10%, UNIVAC 10%
- Honeywell 6000 & 635, IBM 370/168, Xerox 940 & Sigma 9, DEC PDP-10, UNIVAC 1108



Berkeley NOW Project
Supercomputers
Server Farms (e.g., Oceano)

P2P Systems (90s-00s)
- Many Millions of users
- Many GB per day

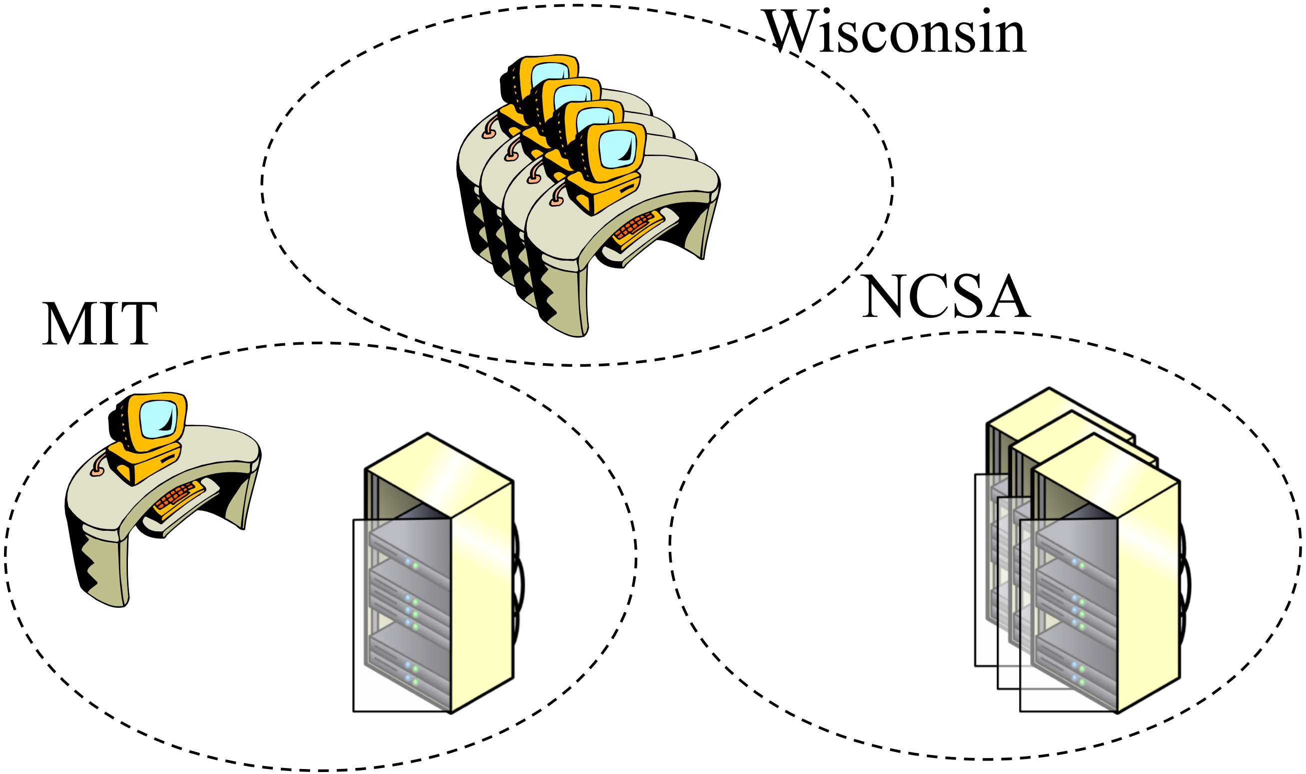


Grids (1980s-2000s):
- GriPhyN (1970s-80s)
- Open Science Grid and Lambda Rail (2000s)
- Globus & other standards (1990s-2000s)

Example: Rapid Atmospheric Modeling System, ColoState U

- Hurricane Georges, 17 days in Sept 1998
  - "RAMS modeled the mesoscale convective complex that dropped so much rain, in good agreement with recorded data"
  - Used 5 km spacing instead of the usual 10 km
  - Ran on 256+ processors
- Computation-intenstive computing (or HPC = high performance computing)
- *Can one run such a program without access to a supercomputer?*

Distributed Computing Resources

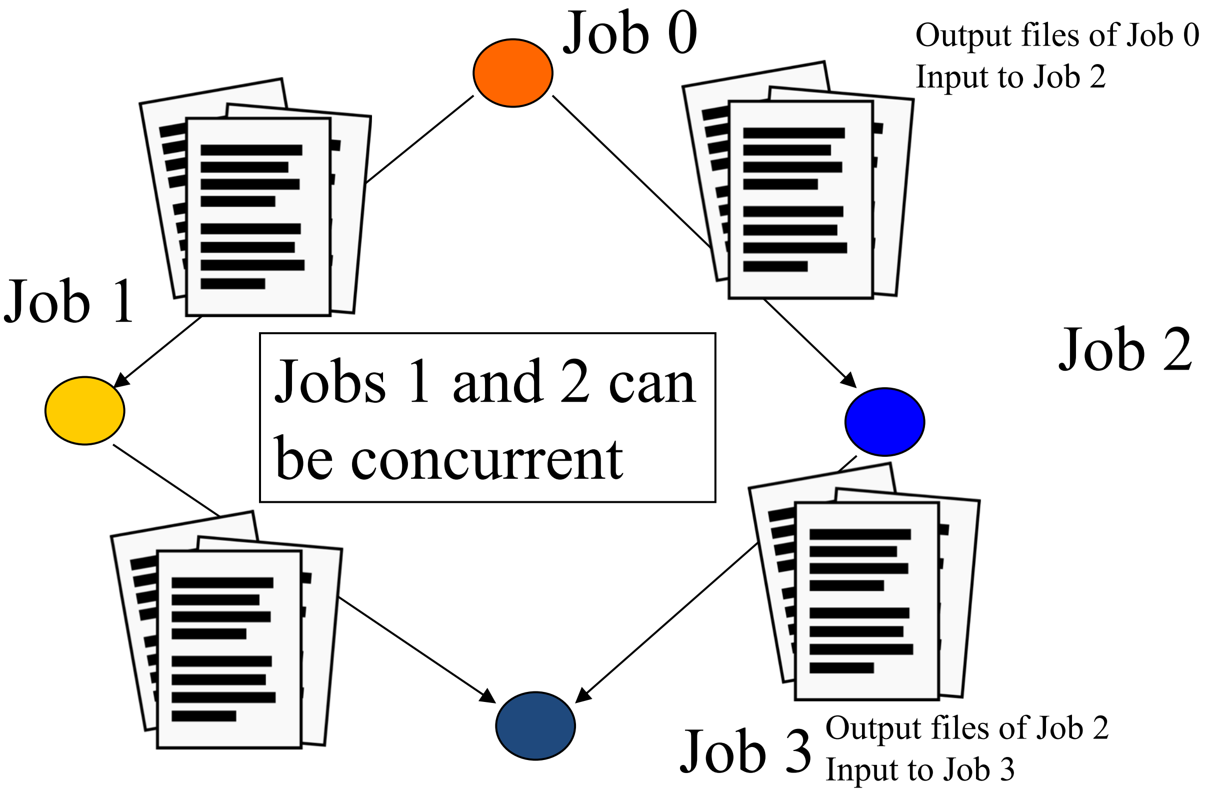


Wisconsin

MIT

NCSA

An Application Coded by a Physicist



Job 0

Output files of Job 0
Input to Job 2

Job 1

Jobs 1 and 2 can be concurrent

Job 2

Job 3

Output files of Job 2
Input to Job 3

An Application Coded by a Physicist

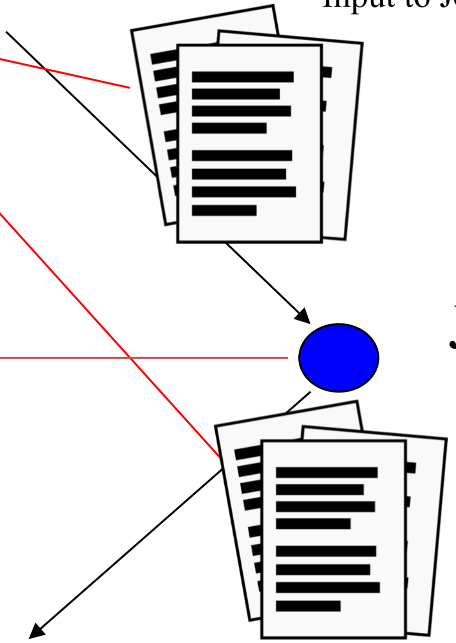
Several GBs

May take several hours/days

4 stages of a job

- Init
- Stage in
- Execute
- Stage out
- Publish

Computation Intensive,
so Massively Parallel

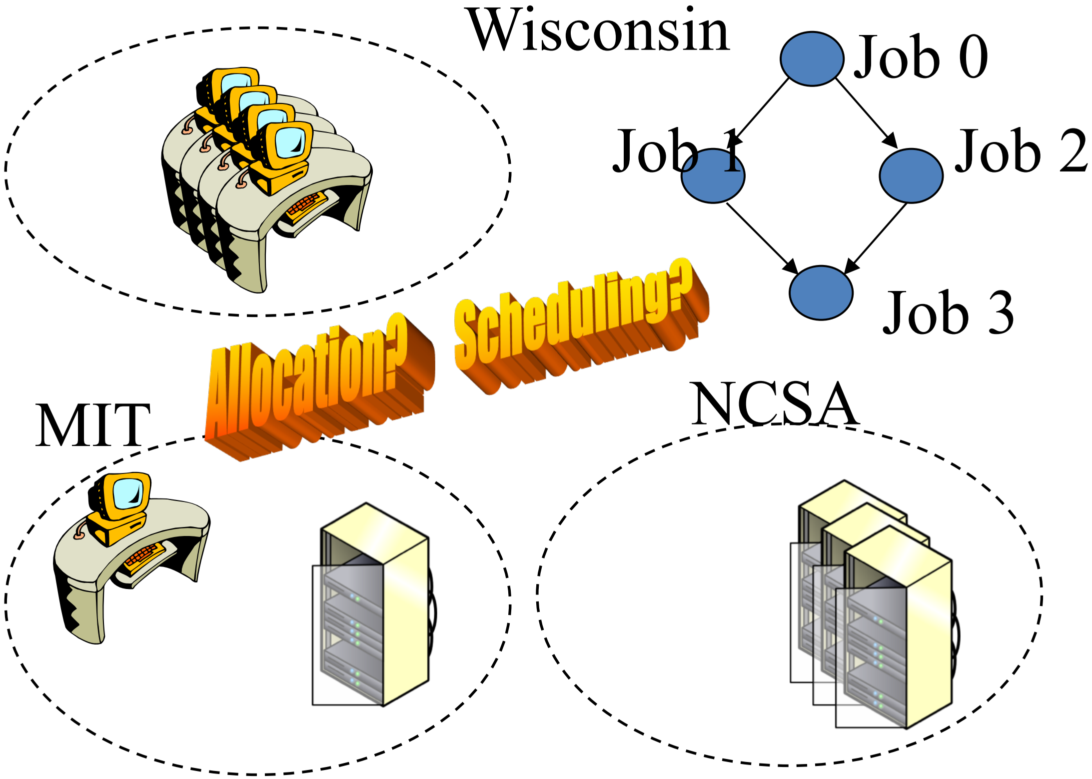


Output files of Job 0
Input to Job 2

Job 2

Output files of Job 2
Input to Job 3

Scheduling Problem



Wisconsin

Job 0

Job 1

Job 2

Job 3

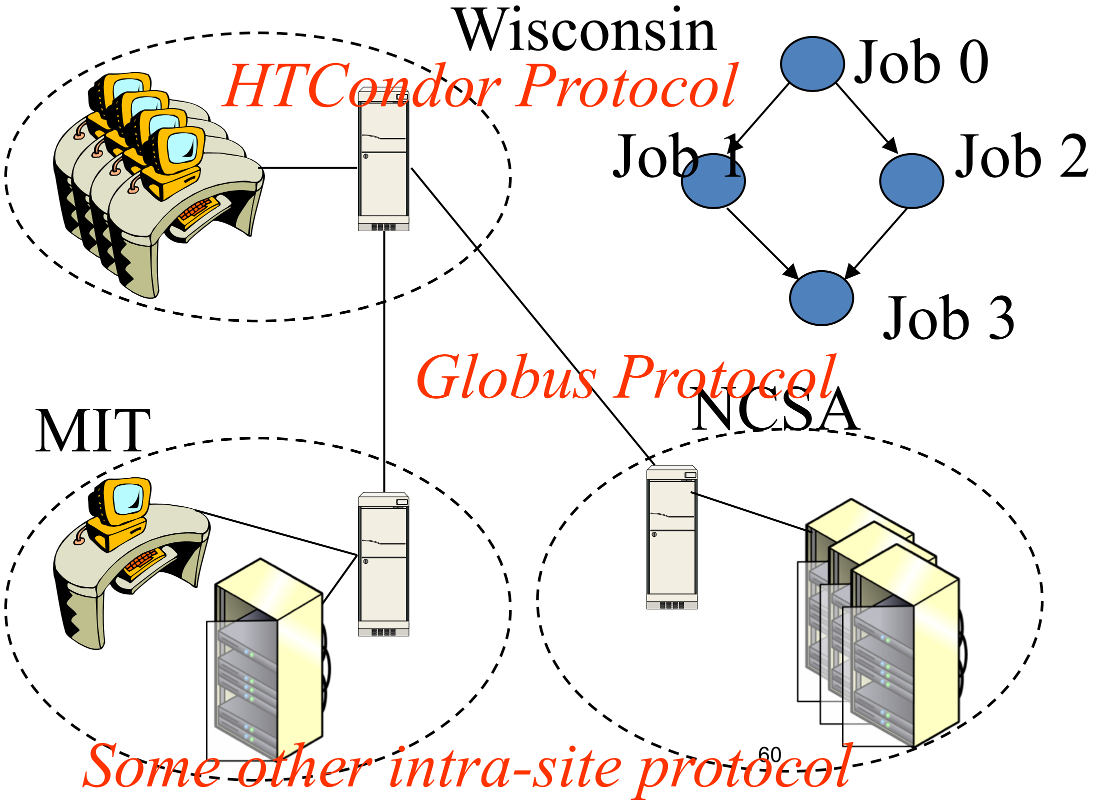
Allocation?

Scheduling?

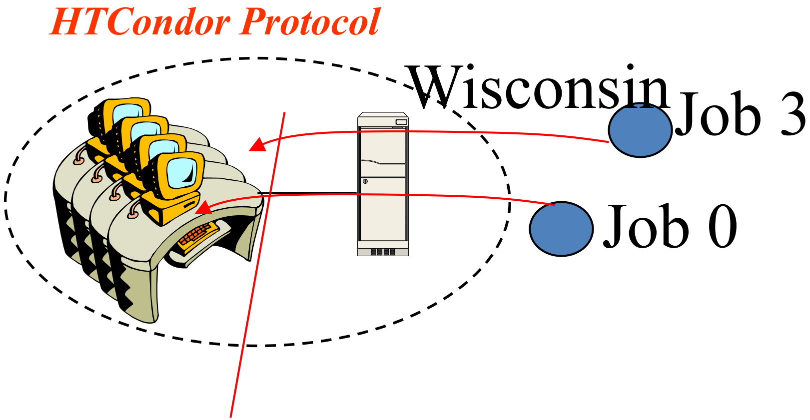
MIT

NCSA

2-level Scheduling Infrastructure



Intra-site Protocol



Internal Allocation & Scheduling

Monitoring

Distribution and Publishing of Files

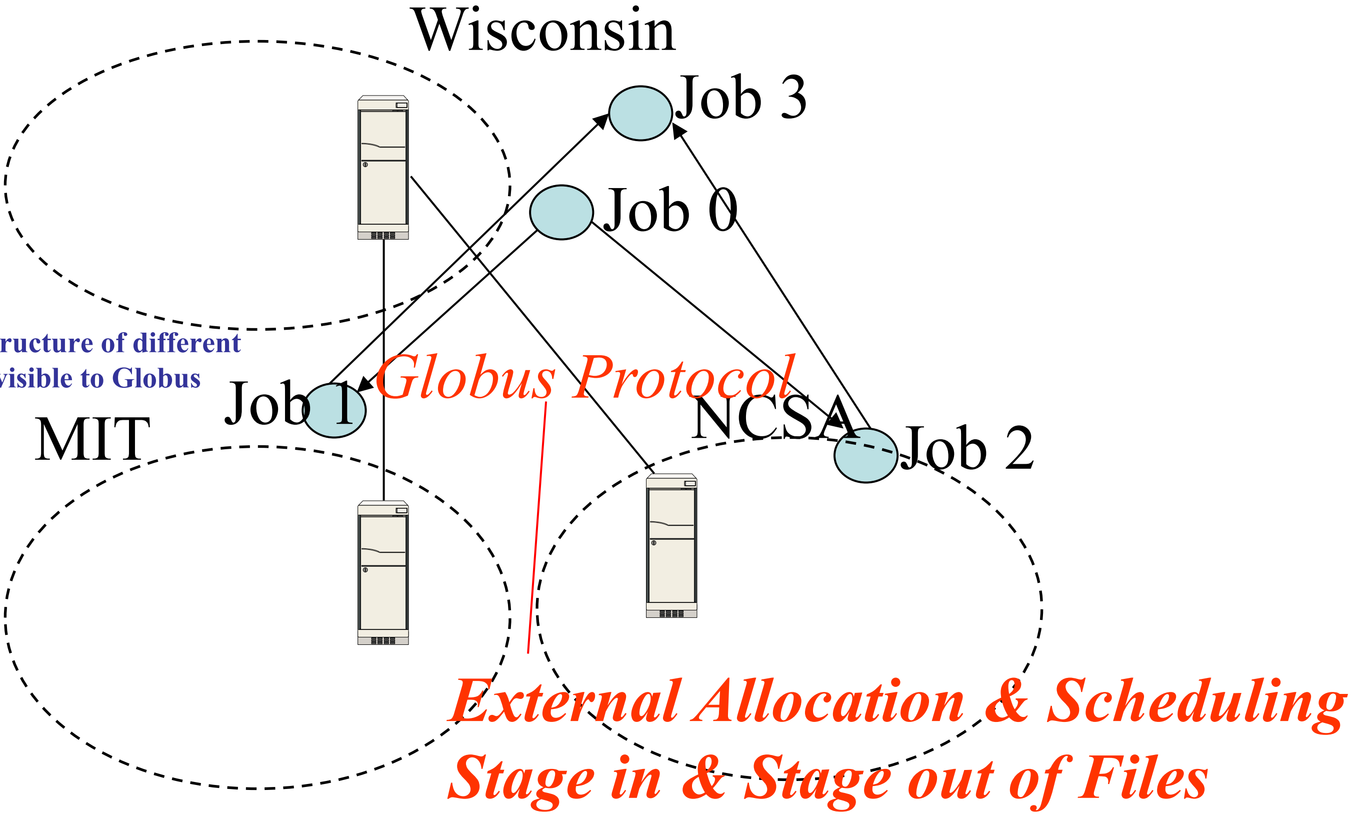
Condor (now HTCondor)

- High-throughput computing system from U. Wisconsin Madison
- Belongs to a class of "Cycle-scavenging" systems
  - SETI@Home and Folding@Home are other systems in this category

Such systems

- Run on a lot of workstations
- When workstation is free, ask site's central server (or Globus) for tasks
- If user hits a keystroke or mouse click, stop task
  - Either kill task or ask server to reschedule task
- Can also run on dedicated machines

Inter-site Protocol



Internal structure of different sites invisible to Globus

External Allocation & Scheduling
Stage in & Stage out of Files

Globus

- Globus Alliance involves universities, national US research labs, and some companies
- Standardized several things, especially software tools
- Separately, but related: Open Grid Forum
- Globus Alliance has developed the Globus Toolkit

http://toolkit.globus.org/toolkit/

Globus Toolkit

- Open-source
- Consists of several components
  - GridFTP: Wide-area transfer of bulk data
  - GRAM5 (Grid Resource Allocation Manager): submit, locate, cancel, and manage jobs
    - Not a scheduler
    - Globus communicates with the schedulers in intra-site protocols like HTCondor or Portable Batch System (PBS)
  - RLS (Replica Location Service): Naming service that translates from a file/dir name to a target location (or another file/dir name)
  - Libraries like XIO to provide a standard API for all Grid IO functionalities
  - Grid Security Infrastructure (GSI)

Security Issues

- Important in Grids because they are *federated,* i.e., no single entity controls the entire infrastructure

- Single sign-on: collective job set should require once-only user authentication
- Mapping to local security mechanisms: some sites use Kerberos, others using Unix
- Delegation: credentials to access resources inherited by subcomputations, e.g., job 0 to job 1
- Community authorization: e.g., third-party authentication

- These are also important in clouds, but less so because clouds are typically run under a central control
- In clouds the focus is on failures, scale, on-demand nature

Summary

- Grid computing focuses on computation-intensive computing (HPC)
- Though often federated, architecture and key concepts have a lot in common with that of clouds
- Are Grids/HPC converging towards clouds?
  - E.g., Compare OpenStack and Globus

Announcements

- MP1: Due this Sunday, demos Monday
  - VMs distributed: see Piazza
  - Demo signup sheet: soon on Piazza
  - Demo details: see Piazza
    - Make sure you print individual and total linecounts
- Check Piazza often! It’s where all the announcements are at!