

CS 421 Lecture 25: Parallel programming and functional languages

- Lecture outline
 - Why parallel programming is hard
 - Why functional programming helps
 - Two case studies
 - Google's MapReduce
 - F#'s asynchronous workflows

Why parallel programming is hard

- Dependencies
- Race conditions
- Deadlock

Granularity of parallelism

- Instruction-level parallelism
 - “Fine grained”
- Higher-level parallelism
 - “Coarse grained”

Approaches to parallel programming

- Automatic parallelization, *i.e.*, parallelizing compilers
- Manual parallelization – low-level
 - MPI, OpenMP
- Manual parallelization – high-level
 - Languages incorporate abstract models of parallelism
 - Libraries implement models of parallelism

Problem

- Automatic parallelization may not know how to take advantage of high-level information
 - Must rely only on logical dependencies to parallelize code
- Manual parallelization has a scalability problem
 - A few threads – OK
 - What about multicore processors? 8, 16, ..., 1024 cores?
 - Also, hard to keep track of all dependencies – high bug potential

Why functional languages help

- Reduce number of dependencies
 - Makes both automatic and manual methods easier
- *E.g.*, in application of map function, applications of function to each element are usually independent:

Why functional languages help

“Due to the absence of side-effects in a purely functional program, it is relatively easy to partition programs so that sub-programs can be executed in parallel: any computation which is needed to produce the result of the program may be run as a separate task. ...

“Higher-order functions (functions which act on functions) can also introduce program-specific control structures, which may be exploited by suitable parallel implementations.”

- *Kevin Hammond, www-fp.dcs.stand.ac.uk/~kh/papers/pasco94/pasco94.html*

Why functional languages help

- Consider imperative and functional implementations of quicksort

- Imperative

```
quicksort(a, lo, hi):  
    p = choose pivot, move to a[lo]  
    partition (a, lo+1, hi, pivot)  
    quicksort(a, lo+1, (lo+hi)/2)  
    quicksort(a, (lo+hi)/2+1, hi)
```

- Functional

```
quicksort(lis):  
    p = choose pivot, remove from lis  
    l = quicksort (filter (< p) lis)  
    u = quicksort (filter (>= p) lis)  
    l @ [p] @ u
```


Two case studies

- Google's MapReduce
 - Parallelism in processing large amounts of data from multiple processors in a data center
 - Library-based model of parallelism
- Microsoft's F# w/ asynchronous workflows
 - Programming model for parallelism in functional language

Google's MapReduce

- Used to access data from Google's data centers.
- Inspired by map and reduce (fold) operations:
 - Divide calculation into two parts:
 - map – apply function to data independently on a set of processors
 - reduce – combine results of map operations
- Available to public in “Hadoop” implementation
- More info: Dean & Ghemawat, “MapReduce: Simplified data processing in large clusters”

Google's MapReduce

- User defines (usually in C++) functions map and reduce:

```
map: string * string -> (string * string) list
```

```
reduce: string * (string list) -> string list
```

- **map** is executed on a collection of processors, producing a list of (key, value) pairs on each
- The underlying MapReduce library combines these pairs, groups and sorts by key, then calls **reduce** for each key, giving all the values associated with that key. It returns the combined list of all values returned from these calls.

Word-counting

- `map (string docname, string doccontents):`
 for each word `w` in `doccontents`:
 `emit (w, "1")`
- `reduce (string word, list<string> counts):`
 `int result = 0`
 for each `n` in `counts`:
 `result := parseInt(n)`
 `emit([""+result])`
- User also supplies mapreduce specification object telling system how to get started (*e.g.*, document names to apply map to)

F#'s asynchronous workflows

- F# a .NET implementation of (a variant of) OCaml.
- “Asynchronous workflows” is a way to turn ordinary programs into parallel programs.
- Based on language feature called “computation expressions”
- Underlying implementation uses “Task Parallel Library”
- Video (PDC 2008) - <http://channel9.msdn.com/pdc2008/TL11/>

How asynchronous workflows work

- “Computation expressions,” are an F# feature, inspired by the Haskell “monad” feature, which allows for a kind of reflection.
- Computation expressions allow certain language constructs to be re-interpreted using user-supplied semantics. The Async library is a workflow.

Computation expressions

- `seq { ... yield e ... }` executes “... yield e ...” and gathers the values of `e` into a list.
- Within “...”, can use limited number of constructs:
 - `use var=expr in expr`
 - `let var=expr in expr`
 - `expr; expr`
 - `yield expr, ...`
- “seq” is not a keyword, but the name of an object that says how to interpret these language constructs.

Computation expressions

- General form of computation expression:
name { ... expression as above ... }
 - name must be bound to an object of a class that implements these operations:
 - Bind: $\alpha \text{ comp} * (\alpha \rightarrow \beta \text{ comp}) \rightarrow \beta \text{ comp}$
 - Delay: $(\text{unit} \rightarrow \alpha \text{ comp}) \rightarrow \alpha \text{ comp}$
 - Let: $\alpha \text{ comp} * (\alpha \rightarrow \alpha \text{ comp}) \rightarrow \alpha \text{ comp}$
 - Return: $\alpha \rightarrow \alpha \text{ comp}$
- where comp is any type constructor you want (*e.g.*, list).

Computation expressions (cont.)

- The definitions of the above operators are used to interpret the syntax within the computation expression.

E.g.,

```
c { let n1 = f in1
    let n2 = g in2
    let sum = n1+n2
    yield sum }
```

would translate (statically) to

```
c.Delay(fun () ->
  c.Bind(f in1, (fun n1 ->
    c.Bind(f in2, (fun n2 ->
      c.Let(n1+n2, (fun sum -> c.Return sum))))))))))
```

Asynchronous workflows

- Asynchronous workflows are an application of computation expressions.

The Async module implements these operations (among others) using the Async type constructor:

Bind: $\alpha \text{ Async} * (\alpha \rightarrow \beta \text{ Async}) \rightarrow \beta \text{ Async}$

Return: $\alpha \rightarrow \alpha \text{ Async}$

plus these methods:

- Run: $\alpha \text{ Async} * \text{int} * \text{bool} \rightarrow \alpha$
- Parallel: $(\alpha \text{ Async}) \text{ list} \rightarrow (\alpha \text{ list}) \text{ Async}$
- Spawn: $\text{unit Async} \rightarrow \text{unit}$

Reminders

- Final review session tomorrow
- Still have lecture on Thursday
 - But not on the exam
- Unit projects due this week