

# Programming Languages and Compilers (CS 421)

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Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha



# Terms

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- A function is in **Direct Style** when it returns its result back to the caller.
- A function is in **Continuation Passing Style** when it, and every function call in it, passes its result to another function.
- Instead of returning the result to the caller, we pass it forward to another function giving the computation after the call.



# Simple Functions Taking Continuations

---

- Given a primitive operation, can convert it to pass its result forward to a continuation

- Examples:

```
# let subk (x, y) k = k(x - y);;
```

```
val subk : int * int -> (int -> 'a) -> 'a = <fun>
```

```
# let eqk (x, y) k = k(x = y);;
```

```
val eqk : 'a * 'a -> (bool -> 'b) -> 'b = <fun>
```

```
# let timesk (x, y) k = k(x * y);;
```

```
val timesk : int * int -> (int -> 'a) -> 'a = <fun>
```



# Nesting Continuations

---

```
# let add_triple (x, y, z) = (x + y) + z;;  
val add_triple : int * int * int -> int = <fun>  
# let add_triple (x,y,z)=let p = x + y in p + z;;  
val add_triple : int * int * int -> int = <fun>  
# let add_triple_k (x, y, z) k =  
    addk (x, y) (fun p -> addk (p, z) k);;  
val add_triple_k: int * int * int -> (int -> 'a) ->  
'a = <fun>
```



## add\_three: a different order

---

- `# let add_triple (x, y, z) = x + (y + z);;`
- How do we write `add_triple_k` to use a different order?
  
- `let add_triple_k (x, y, z) k =`



## add\_three: a different order

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- `# let add_triple (x, y, z) = x + (y + z);;`
- How do we write `add_triple_k` to use a different order?
- `let add_triple_k (x, y, z) k =  
 addk (y,z) (fun r -> addk(x,r) k)`



## add\_three: a different order

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- `# let add_triple (x, y, z) = x + (y + z);;`
- How do we write `add_triple_k` to use a different order?
- `let add_triple_k (x, y, z) k =`  
    `addk (y,z) (fun r -> addk(x,r) k)`



## add\_three: a different order

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- `# let add_triple (x, y, z) = x + (y + z);;`
- How do we write `add_triple_k` to use a different order?
- `let add_triple_k (x, y, z) k =`  
    `addk (y,z) (fun r -> addk(x,r) k)`





# Recursive Functions

---

## ■ Recall:

```
# let rec factorial n =
```

```
  if n = 0 then 1 else n * factorial (n - 1);;
```

```
val factorial : int -> int = <fun>
```

```
# factorial 5;;
```

```
- : int = 120
```



# Recursive Functions

---

```
# let rec factorial n =  
  let b = (n = 0) in (* First computation *)  
  if b then 1 (* Returned value *)  
  else let s = n - 1 in (* Second computation *)  
        let r = factorial s in (* Third computation *)  
        n * r (* Returned value *) ;;  
val factorial : int -> int = <fun>  
# factorial 5;;  
- : int = 120
```



# Recursive Functions

---

```
# let rec factorialk n k =  
  eqk (n, 0)  
  (fun b -> (* First computation *)  
    if b then k 1 (* Passed value *)  
    else subk (n, 1) (* Second computation *)  
    (fun s -> factorialk s (* Third computation *)  
      (fun r -> timesk (n, r) k))) (* Passed value *)  
val factorialk : int -> (int -> 'a) -> 'a = <fun>  
# factorialk 5 report;;  
120  
- : unit = ()
```



# Recursive Functions

---

- To make recursive call, must build intermediate continuation to
  - take recursive value:  $r$
  - build it to final result:  $n * r$
  - And pass it to final continuation:
    - $\text{times}(n, r) k = k(n * r)$



# Recursive Functions

---

```
# let rec factorialk n k =  
  eqk (n, 0)  
  (fun b -> (* First computation *)  
    if b then k 1 (* Passed value *)  
    else subk (n, 1) (* Second computation *)  
    (fun s -> factorialk s (* Third computation *)  
      (fun r -> timesk (n, r) k))) (* Passed value *)  
val factorialk : int -> (int -> 'a) -> 'a = <fun>  
# factorialk 5 report;;  
120  
- : unit = ()
```



## Example: CPS for length

---

```
let rec length list = match list with [] -> 0  
  | (a :: bs) -> 1 + length bs
```

What is the let-expanded version of this?



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let rec length list = match list with [] -> 0  
  | (a :: bs) -> let r1 = length bs in 1 + r1
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## Example: CPS for length

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```
#let rec length list = match list with [] -> 0  
  | (a :: bs) -> let r1 = length bs in 1 + r1
```

What is the CSP version of this?





# Example: CPS for length

---

```
#let rec length list = match list with [] -> 0
  | (a :: bs) -> let r1 = length bs in 1 + r1
```

What is the CSP version of this?

```
#let rec lengthk list k = match list with [ ] -> k 0
  | x :: xs -> lengthk xs (fun r -> addk (r,1) k);;
```

```
val lengthk : 'a list -> (int -> 'b) -> 'b = <fun>
```

```
# lengthk [2;4;6;8] report;;
```

```
4
```

```
- : unit = ()
```



# CPS for sum

---

```
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> x + sum xs ;;
val sum : int list -> int = <fun>
```



# CPS for sum

---

```
# let rec sum list = match list with [ ] -> 0  
  | x :: xs -> x + sum xs ;;
```

```
val sum : int list -> int = <fun>
```

```
# let rec sum list = match list with [ ] -> 0  
  | x :: xs -> let r1 = sum xs in x + r1;;
```



# CPS for sum

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# let rec sum list = match list with [ ] -> 0  
  | x :: xs -> x + sum xs ;;
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# let rec sum list = match list with [ ] -> 0  
  | x :: xs -> let r1 = sum xs in x + r1;;
```

```
val sum : int list -> int = <fun>
```

```
# let rec sumk list k = match list with [ ] -> k 0  
  | x :: xs -> sumk xs (fun r1 -> addk x r1 k);;
```



# CPS for sum

---

```
# let rec sum list = match list with [ ] -> 0
```

```
  | x :: xs -> x + sum xs ;;
```

```
val sum : int list -> int = <fun>
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# let rec sum list = match list with [ ] -> 0
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```
  | x :: xs -> let r1 = sum xs in x + r1;;
```

```
val sum : int list -> int = <fun>
```

```
# let rec sumk list k = match list with [ ] -> k 0
```

```
  | x :: xs -> sumk xs (fun r1 -> addk (x, r1) k);;
```

```
val sumk : int list -> (int -> 'a) -> 'a = <fun>
```

```
# sumk [2;4;6;8] report;;
```

```
20
```

```
- : unit = ()
```



# CPS for Higher Order Functions

---

- In CPS, every procedure / function takes a continuation to receive its result
- Procedures passed as arguments take continuations
- Procedures returned as results take continuations
- CPS version of higher-order functions must expect input procedures to take continuations



## Example: all

---

```
#let rec all (p, l) = match l with [] -> true
```

```
  | (x :: xs) -> let b = p x in
```

```
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

- What is the CPS version of this?



## Example: all

---

```
#let rec all (p, l) = match l with [] -> true  
  | (x :: xs) -> let b = p x in  
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

- What is the CPS version of this?

```
#let rec allk (pk, l) k =
```





## Example: all

---

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

- What is the CPS version of this?

```
#let rec allk (pk, l) k = match l with [] -> true
```



## Example: all

---

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

- What is the CPS version of this?

```
#let rec allk (pk, l) k = match l with [] -> k true
```



## Example: all

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```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

- What is the CPS version of this?

```
#let rec allk (pk, l) k = match l with [] -> k true
  | (x :: xs) ->
```



## Example: all

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```
#let rec all (p, l) = match l with [] -> true  
  | (x :: xs) -> let b = p x in  
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

- What is the CPS version of this?

```
#let rec allk (pk, l) k = match l with [] -> k true  
  | (x :: xs) -> pk x
```



## Example: all

---

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

- What is the CPS version of this?

```
#let rec allk (pk, l) k = match l with [] -> k true
  | (x :: xs) -> pk x
    (fun b -> if b then
    ) else
```



## Example: all

---

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
```

```
val all : ('a -> bool) -> 'a list -> bool = <fun>
```

■ What is the CPS version of this?

```
#let rec allk (pk, l) k = match l with [] -> k true
  | (x :: xs) -> pk x
    (fun b -> if b then allk (pk, xs) k else k
false)
```

```
val allk : ('a -> (bool -> 'b) -> 'b) * 'a list ->
(bool -> 'b) -> 'b = <fun>
```



# Terminology: Review

---

- A function is in **Direct Style** when it returns its result back to the caller.
- A function is in **Continuation Passing Style** when it, and every function call in it, passes its result to another function.
- A **Tail Call** occurs when a function returns the result of another function call without any more computations (eg tail recursion)
- Instead of returning the result to the caller, we pass it forward to another function giving the computation after the call.



# CPS Transformation

---

- Step 1: Add continuation argument to any function definition:
  - $\text{let } f \text{ arg} = e \Rightarrow \text{let } f \text{ arg } k = e$
  - Idea: Every function takes an extra parameter saying where the result goes
- Step 2: A simple expression in tail position should be passed to a continuation instead of returned:
  - $\text{return } a \Rightarrow k \ a$
  - Assuming  $a$  is a constant or variable.
  - “Simple” = “No available function calls.”





# CPS Transformation

---

- Step 3: Pass the current continuation to every function call in tail position
  - $\text{return } f \text{ arg} \Rightarrow f \text{ arg } k$
  - The function “isn’ t going to return,” so we need to tell it where to put the result.



# CPS Transformation

---

- Step 4: Each function call not in tail position needs to be converted to take a new continuation (containing the old continuation as appropriate)
  - $\text{return op (f arg)} \Rightarrow \text{f arg (fun r -> k(op r))}$
  - op represents a primitive operation
  - $\text{return g(f arg)} \Rightarrow \text{f arg (fun r-> g r k)}$



# Example

---

## Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

## After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)
```



# Example

---

## Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

## After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  k false (* rule 2 *)  
  
  k true (* rule 2 *)
```



# Example

---

## Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

## After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  
  k false (* rule 2 *)  
  
  k true (* rule 2 *)  
  memk (y, xs) k (* rule 3 *)
```



# Example

---

## Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
| x :: xs ->  
  if (x = y)  
  then true  
  else mem(y,xs);;
```

## After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  k false (* rule 2 *)  
  
  eqk (x, y)  
  (fun b -> b (* rule 4 *))  
  k true (* rule 2 *)  
  memk (y, xs) (* rule 3 *)
```



# Example

---

## Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
| x :: xs ->  
  if (x = y)  
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let rec memk (y,lst) k =  
  (* rule 1 *)  
  k false (* rule 2 *)  
  
  eqk (x, y)  
  (fun b -> if b (* rule 4 *)  
  then k true (* rule 2 *)  
  else memk (y, xs) (* rule 3 *))
```



# Example

---

## Before:

```
let rec mem (y,lst) =
```

```
  match lst with
```

```
    [ ] -> false
```

```
  | x :: xs ->
```

```
    if (x = y)
```

```
      then true
```

```
      else mem(y,xs);;
```

## After:

```
let rec memk (y,lst) k =  
      (* rule 1 *)
```

```
  match lst with
```

```
    | [ ] -> k false (* rule 2 *)
```

```
    | x :: xs ->
```

```
      eqk (x, y)
```

```
        (fun b -> if b (* rule 4 *))
```

```
      then k true (* rule 2 *)
```

```
        else memk (y, xs) k (* rule 3 *)
```





# Example

---

## Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
| x :: xs ->  
  if (x = y)  
  then true  
  else mem(y,xs);;
```

## After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  match lst with  
  | [ ] -> k false (* rule 2 *)  
  | x :: xs ->  
    eqk (x, y)  
    (fun b -> if b (* rule 4 *)  
  then k true (* rule 2 *)  
    else memk (y, xs) k (* rule 3 *)
```



# Example

---

## Before:

```
let rec add_list lst =  
  match lst with  
  | [] -> 0  
  | 0 :: xs -> add_list xs  
  | x :: xs -> (+) x  
    (add_list xs);;
```

## After:

```
let rec add_listk lst k =  
  (* rule 1 *)  
  match lst with  
  | [] -> k 0 (* rule 2 *)  
  | 0 :: xs -> add_listk xs k  
    (* rule 3 *)  
  | x :: xs -> add_listk xs  
    (fun r -> k ((+) x r));;  
  (* rule 4 *)
```