

# Programming Languages and Compilers (CS 421)

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<https://courses.engr.illinois.edu/cs421/sp2023>

Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha

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## Functions

```
# let plus_two n = n + 2;;  
val plus_two : int -> int = <fun>  
# plus_two 17;;  
- : int = 19  
# let plus_two = fun n -> n + 2;;  
val plus_two : int -> int = <fun>  
# plus_two 14;;  
- : int = 16
```

First definition syntactic sugar for second

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## Using a nameless function

```
# (fun x -> x * 3) 5;; (* An application *)  
- : int = 15  
# ((fun y -> y +. 2.0), (fun z -> z * 3));;  
(* As data *)  
- : (float -> float) * (int -> int) = (<fun>, <fun>)
```

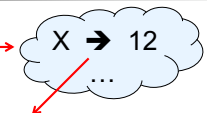
Note: in fun v -> exp(v), scope of variable is only the body exp(v)

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## Values fixed at declaration time

```
# let x = 12;;  
val x : int = 12  
# let plus_x y = y + x;;  
val plus_x : int -> int = <fun>  
# plus_x 3;;
```



What is the result?

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## Values fixed at declaration time

```
# let x = 12;;  
val x : int = 12  
# let plus_x y = y + x;;  
val plus_x : int -> int = <fun>  
# plus_x 3;;  
- : int = 15
```

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## Values fixed at declaration time

```
# let x = 7;; (* New declaration, not an update *)  
val x : int = 7  
# plus_x 3;;
```

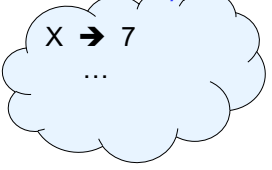
What is the result this time?

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## Values fixed at declaration time

```
# let x = 7;; (* New declaration, not an update *)  
val x : int = 7  
# plus_x 3;;
```



What is the result this time?

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## Values fixed at declaration time

```
# let x = 7;; (* New declaration, not an update *)  
val x : int = 7  
  
# plus_x 3;;  
- : int = 15
```

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## Question

- Observation: Functions are first-class values in this language
- Question: What value does the environment record for a function variable?
- Answer: a closure

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## Save the Environment!

- A *closure* is a pair of an environment and an association of a formal parameter (the input variables)\* with an expression (the function body), written:  
$$f \rightarrow \langle (v_1, \dots, v_n) \rightarrow \text{exp}, \rho_f \rangle$$
- Where  $\rho_f$  is the environment in effect when  $f$  is defined (if  $f$  is a simple function)
- \* Will come back to the "formal parameter"

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## Closure for plus\_x

- When plus\_x was defined, had environment:  
$$\rho_{\text{plus\_x}} = \{\dots, x \rightarrow 12, \dots\}$$
- Recall: `let plus_x y = y + x`  
is really `let plus_x = fun y -> y + x`
- Closure for `fun y -> y + x`:  
$$\langle y \rightarrow y + x, \rho_{\text{plus\_x}} \rangle$$
- Environment just after plus\_x defined:  
$$\{\text{plus\_x} \rightarrow \langle y \rightarrow y + x, \rho_{\text{plus\_x}} \rangle\} + \rho_{\text{plus\_x}}$$

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Now it's your turn

You should be able to complete ACT1

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## Functions with more than one argument

```
# let add_three x y z = x + y + z;;
val add_three : int -> int -> int -> int = <fun>
# let t = add_three 6 3 2;;
val t : int = 11
# let add_three =
  fun x -> (fun y -> (fun z -> x + y + z));;
val add_three : int -> int -> int -> int = <fun>
```

Again, first syntactic sugar for second

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## Functions with more than one argument

```
# let add_three x y z = x + y + z;;
val add_three : int -> int -> int -> int = <fun>
  ■ What is the value of add_three?
  ■ Let  $\rho_{\text{add\_three}}$  be the environment before the declaration
  ■ Remember:
let add_three =
  fun x -> (fun y -> (fun z -> x + y + z));;
Value: <x ->fun y -> (fun z -> x + y + z),  $\rho_{\text{add\_three}}$  >
```

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## Partial application of functions

```
let add_three x y z = x + y + z;;
```

```
# let h = add_three 5 4;;
val h : int -> int = <fun>
# h 3;;
- : int = 12
# h 7;;
- : int = 16
```

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## Partial application of functions

```
let add_three x y z = x + y + z;;
```

```
# let h = add_three 5 4;;
val h : int -> int = <fun>
# h 3;;
- : int = 12
# h 7;;
- : int = 16
```

- Partial application also called *sectioning*

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## Functions as arguments

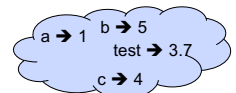
```
# let thrice f x = f (f (f x));;
val thrice : ('a -> 'a) -> 'a -> 'a = <fun>
# let g = thrice plus_two;;
val g : int -> int = <fun>
# g 4;;
- : int = 10
# thrice (fun s -> "Hi! " ^ s) "Good-bye!";;
- : string = "Hi! Hi! Hi! Good-bye!"
```

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## Tuples as Values

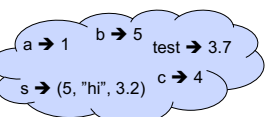
```
//  $\rho_7 = \{c \rightarrow 4, \text{test} \rightarrow 3.7, a \rightarrow 1, b \rightarrow 5\}$ 
```



```
# let s = (5, "hi", 3.2);;
```

```
val s : int * string * float = (5, "hi", 3.2)
```

```
//  $\rho_8 = \{s \rightarrow (5, \text{"hi"}, 3.2), c \rightarrow 4, \text{test} \rightarrow 3.7, a \rightarrow 1, b \rightarrow 5\}$ 
```

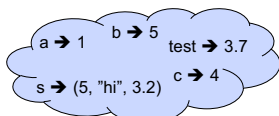


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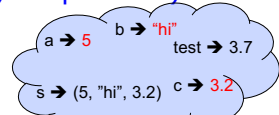
## Pattern Matching with Tuples

```
/ ρ8 = {s → (5, "hi", 3.2),
        c → 4, test → 3.7,
        a → 1, b → 5}
```



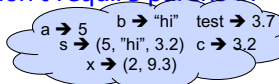
```
# let (a,b,c) = s;; (* (a,b,c) is a pattern *)
```

```
val a : int = 5
val b : string = "hi"
val c : float = 3.2
```



```
# let x = 2, 9.3;; (* tuples don't require parens in Ocaml *)
```

```
val x : int * float = (2, 9.3)
```



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## Nested Tuples

```
# (*Tuples can be nested *)
```

```
let d = ((1,4,62),("bye",15),73.95);;
```

```
val d : (int * int * int) * (string * int) * float =
  ((1, 4, 62), ("bye", 15), 73.95)
```

```
# (*Patterns can be nested *)
```

```
let (p,(st,_,_) = d;; (* _ matches all, binds nothing *)
```

```
val p : int * int * int = (1, 4, 62)
```

```
val st : string = "bye"
```

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## Functions on tuples

```
# let plus_pair (n,m) = n + m;;
val plus_pair : int * int -> int = <fun>
# plus_pair (3,4);;
- : int = 7
# let double x = (x,x);;
val double : 'a -> 'a * 'a = <fun>
# double 3;;
- : int * int = (3, 3)
# double "hi";;
- : string * string = ("hi", "hi")
```

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## Match Expressions

```
# let triple_to_pair triple =
```

```
  match triple
```

```
  with (0, x, y) -> (x, y)
```

```
      | (x, 0, y) -> (x, y)
```

```
      | (x, y, _) -> (x, y);;
```

```
val triple_to_pair : int * int * int -> int * int =
  <fun>
```

- Each clause: pattern on left, expression on right
- Each x, y has scope of only its clause
- Use first matching clause

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## Closure for plus\_pair

- Assume  $\rho_{\text{plus\_pair}}$  was the environment just before `plus_pair` defined
- Closure for `plus_pair`:
 
$$\langle (n,m) \rightarrow n + m, \rho_{\text{plus\_pair}} \rangle$$
- Environment just after `plus_pair` defined:
 
$$\{\text{plus\_pair} \rightarrow \langle (n,m) \rightarrow n + m, \rho_{\text{plus\_pair}} \rangle\} \\ + \rho_{\text{plus\_pair}}$$

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## Save the Environment!

- A *closure* is a pair of an environment and an association of a pattern (e.g.  $(v_1, \dots, v_n)$  giving the input variables) with an expression (the function body), written:
 
$$\langle (v_1, \dots, v_n) \rightarrow \text{exp}, \rho \rangle$$
- Where  $\rho$  is the environment in effect when the function is defined (for a simple function)

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