Programming Languages and Compilers (CS 421)

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Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha



Features of OCAML

- Higher order applicative language
- Call-by-value parameter passing
- Modern syntax
- Parametric polymorphism
 - Aka structural polymorphism
- Automatic garbage collection
- User-defined algebraic data types

Why learn OCAML?

- Many features not clearly in languages you have already learned
- Assumed basis for much research in programming language research
- OCAML is particularly efficient for programming tasks involving languages (eg parsing, compilers, user interfaces)
- Industrially Relevant:
 - Jane Street trades billions of dollars per day using OCaml programs
 - Major language supported at Bloomberg
- Similar languages: Microsoft F#, SML, Haskell, Scala

Session in OCAML

% ocaml
Objective Caml version 4.07.1

(* Read-eval-print loop; expressions and declarations *)
2 + 3;; (* Expression *)

-: int = 5

3 < 2;;

-: bool = false

Declarations; Sequencing of Declarations

```
# let x = 2 + 3;; (* declaration *)
val x : int = 5
# let test = 3 < 2;;
val test: bool = false
# let a = 1 let b = a + 4;; (* Sequence of dec
val a : int = 1
val b : int = 5
```

Functions

```
# let plus_two n = n + 2;;
val plus_two : int -> int = <fun>
# plus_two 17;;
- : int = 19
```

Functions

```
let plus_two n = n + 2;;
plus_two 17;;
- : int = 19
```

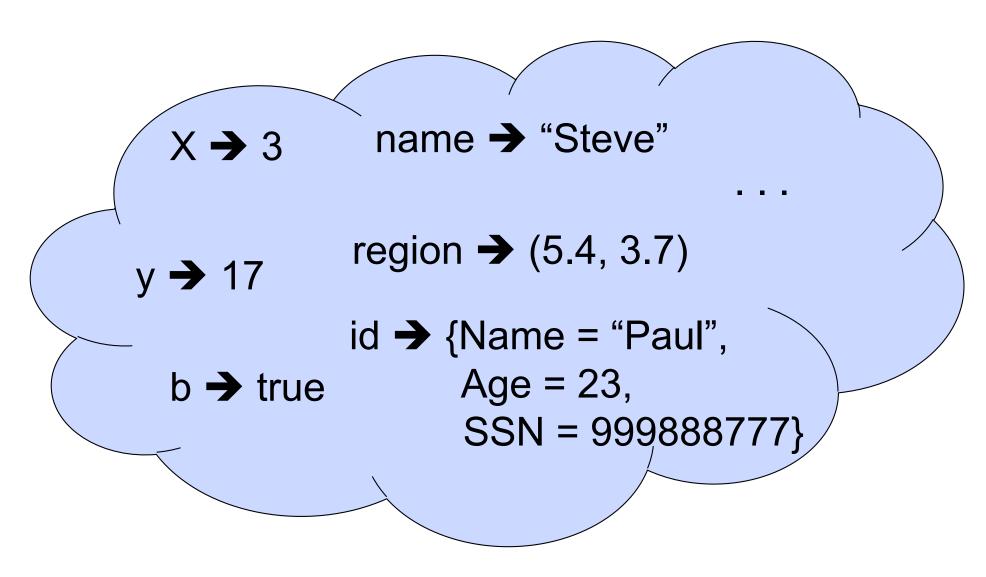
Environments

- Environments record what value is associated with a given identifier
- Central to the semantics and implementation of a language
- Notation

 $\rho = \{name_1 \rightarrow value_1, name_2 \rightarrow value_2, ...\}$ Using set notation, but describes a partial function

- Often stored as list, or stack
 - To find value start from left and take first match

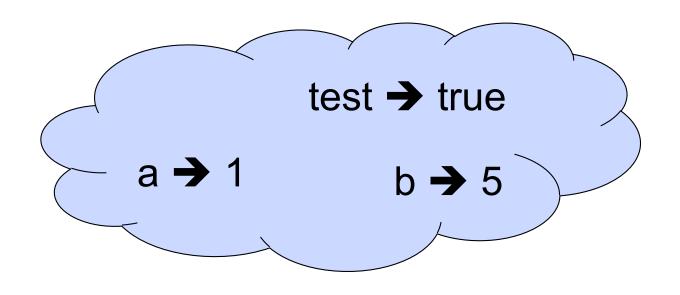
Environments



Global Variable Creation

```
# 2 + 3;; (* Expression *)
// doesn't affect the environment
# let test = 3 < 2;; (* Declaration *)
val test: bool = false
// \rho_1 = \{\text{test} \rightarrow \text{false}\}
# let a = 1 let b = a + 4;; (* Seq of dec *)
// \rho_2 = \{b \rightarrow 5, a \rightarrow 1, \text{ test} \rightarrow \text{ false}\}
```





New Bindings Hide Old

```
// \rho_2 = \{b \rightarrow 5, a \rightarrow 1, \text{ test} \rightarrow \text{false}\}
let test = 3.7;;
```

What is the environment after this declaration?

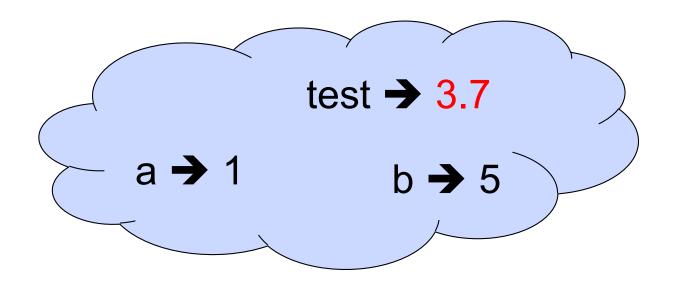
New Bindings Hide Old

```
// \rho_2 = \{b \rightarrow 5, a \rightarrow 1, \text{ test} \rightarrow \text{ false}\}
let test = 3.7;;
```

What is the environment after this declaration?

//
$$\rho_3 = \{ \text{test} \to 3.7, \, a \to 1, \, b \to 5 \}$$







Now it's your turn

You should be able to start ACT1

Local Variable Creation

```
test → 3.7
// \rho_3 = \{\text{test} \rightarrow 3.7, a \rightarrow 1, b \rightarrow 5\}^{a \rightarrow 1}
                                                                      b → 5
# let b = 5 * 4
// \rho_4 = \{b \to 20, \text{ test} \to 3.7, a \to 1\}
                                                             test → 3.7
    in 2 * b;;
-: int = 40
// \rho_5 = \rho_3 = \{\text{test} \to 3.7, \, a \to 1, \, b \to 5\}
                                                                   test → 3.7
# b;;
                                                                    b → 5
-: int = 5
```

Local let binding

```
test → 3.7
// \rho_5 = \rho_3 = \{\text{test} \rightarrow 3.7, a \rightarrow 1, b \rightarrow 5\}
                                                                              b → 5
# let c =
     let b = a + a
// \rho_6 = \{b \rightarrow 2\} + \rho_3
// = \{b \rightarrow 2, \text{ test} \rightarrow 3.7, a \rightarrow 1\}
     in b * b;;
val c : int = 4
// \rho_7 = \{c \to 4, \text{ test } \to 3.7, \text{ a } \to 1, \text{ b } \to 5\}
# b;;
-: int = 5
```

Local let binding

```
test → 3.7
// \rho_5 = \rho_3 = \{\text{test} \rightarrow 3.7, a \rightarrow 1, b \rightarrow 5\}
                                                                           b → 5
# let c =
    let b = a + a
                                                                  test → 3.7
// \rho_6 = \{b \rightarrow 2\} + \rho_3
// = \{b \rightarrow 2, \text{ test} \rightarrow 3.7, \text{ a}\}
     in b** b;*
val c : int = 4
// \rho_7 = \{c \to 4, \text{ test} \to 3.7, a \to 1, b \to 5\}
# b;;
-: int = 5
```

Local let binding

```
test → 3.7
// \rho_5 = \rho_3 = \{\text{test} \rightarrow 3.7, a \rightarrow 1, b \rightarrow 5\}
                                                                         b → 5
# let c =
    let b = a + a
                                                                test → 3.7
// \rho_6 = \{b \rightarrow 2\} + \rho_3
// = \{b \rightarrow 2, \text{ test}\}
    in b** b;*
                                                                          test → 3.7
val c : int = 4
                                                                 c → 4
                                                                         b → 5
// \rho_7 = \{c \to 4, \text{ test} \to 3.7, a \to 1, b \to 5\}
# b;;
-: int = 5
```

Functions

```
# let plus_two n = n + 2;;
val plus_two : int -> int = <fun>
# plus_two 17;;
- : int = 19
```

Functions

```
let plus_two n = n + 2;;
plus_two 17;;
- : int = 19
```



Nameless Functions (aka Lambda Terms)





35

Functions

```
# let plus two n = n + 2;;
val plus_two : int -> int = <fun>
# plus two 17;;
-: int = 19
# let plus_two = fun n \rightarrow n + 2;;
val plus_two : int -> int = <fun>
# plus_two 14;;
-: int = 16
```

First definition syntactic sugar for second

Using a nameless function

```
# (fun x -> x * 3) 5;; (* An application *)
- : int = 15
# ((fun y -> y +. 2.0), (fun z -> z * 3));;
   (* As data *)
- : (float -> float) * (int -> int) = (<fun>,
        <fun>)
```

Note: in fun $v \rightarrow \exp(v)$, scope of variable is only the body $\exp(v)$

Values fixed at declaration time

What is the result?

Values fixed at declaration time

```
# let x = 12;;
val x : int = 12
# let plus_x y = y + x;;
val plus_x : int -> int = <fun>
# plus_x 3;;
- : int = 15
```

Values fixed at declaration time

```
# let x = 7;; (* New declaration, not an
    update *)
val x : int = 7

# plus_x 3;;
```

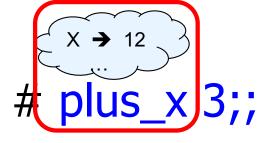
What is the result this time?



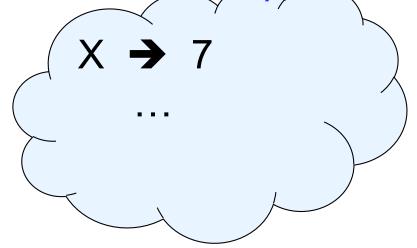
Values fixed at declaration time

```
update *)
```

$$val x : int = 7$$



let x = 7; (* New declaration, not an



What is the result this time?

Values fixed at declaration time

```
# let x = 7;; (* New declaration, not an
    update *)
val x : int = 7

# plus_x 3;;
- : int = 15
```



Observation: Functions are first-class values in this language

• Question: What value does the environment record for a function variable?

Answer: a closure

Save the Environment!

A closure is a pair of an environment and an association of a formal parameter (the input variables)* with an expression (the function body), written:

$$f \rightarrow \langle (v1,...,vn) \rightarrow exp, \rho_f \rangle$$

- Where p_f is the environment in effect when f is defined (if f is a simple function)
- * Will come back to the "formal parameter"

Closure for plus_x

When plus_x was defined, had environment:

$$\rho_{\text{plus } X} = \{..., X \rightarrow 12, ...\}$$

- Recall: let plus_x y = y + x is really let plus_x = fun y -> y + x
- Closure for fun y -> y + x:

$$\langle y \rightarrow y + x, \rho_{plus} \rangle$$

Environment just after plus_x defined:

$$\{plus_x \rightarrow \langle y \rightarrow y + x, \rho_{plus_x} \rangle\} + \rho_{plus_x}$$



Now it's your turn

You should be able to complete ACT1

125 minutes