# Lecture 10: LR parsing and resolving conflicts

- What are conflicts
- Example 1: a simple, unambiguous grammar
  - ocamlyacc output
  - Using parse trees to understand conflict
  - Fixing conflict
- Example 2: ambiguous grammar for conditional expressions
  - Eliminating conflicts using %prec declarations

### Conflicts

- ocamlyacc generates tables saying what action to take at each point in parse
  - Method is called "LALR(1)"
  - "LR(1)" is a similar, but somewhat more powerful, method
     will often use "LR(1)" and "LALR(1)" as synonyms.
- Not every grammar can be parsed using this method.
  - Problem is *always* that ocamlyacc cannot decide on the proper action in some cases
  - "Shift/reduce conflict" cannot decide whether to shift or reduce
  - "Reduce/reduce conflict" knows to reduce, but can't decide which production to use

### Example 1

- $\begin{array}{cccc} \bullet & A & \to & B \ , \ int \\ & B & \to \ id \mid id \ , \ B \end{array}$
- Grammar is unambiguous, but consider these two inputs:
  - x,y,10
  - x,y,z,10
- Both lead to an identical stack/lookahead configuration, but the correct action in one case is shift and in the other is reduce.
- Look at s-r parse, and at two parse trees.

#### Presented to ocamlyacc:

%token int id comma
%start A
%type <int> A
%%
A: B comma int {0}
B: id {0}
| id comma B {0}

#### Using "ocamlyacc -v", file simple.output contains:

3: shift/reduce conflict (shift 6, reduce 2) on comma
state 3
B : id . (2)
B : id . comma B (3)

comma shift 6

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#### One way to fix grammar:

 $\begin{array}{l} A \ \rightarrow \ B \ int \\ b \ \rightarrow \ id \ , \ | \ id \ , \ B \end{array}$ 

Another way to fix grammar:

 $\begin{array}{l} A \ \rightarrow \ B \ , \ int \\ b \ \rightarrow \ id \ \mid B \ , \ id \end{array}$ 

### **Example 2**

#### Ambiguous grammar for conditional expressions:

- $CondExpr \rightarrow id | CondExpr || CondExpr |$
- Consider this input:
  - x || y && z
- Leads to a stack/lookahead configuration in which shifting and reducing both work, but produce different parse trees.
- Look at s-r parse, and at two parse trees.

ocamlyacc -v output contains:

10: shift/reduce conflict (shift 7, reduce 2) on and 10: shift/reduce conflict (shift 8, reduce 2) on or state 10 CondExpr : CondExpr . or CondExpr (2) CondExpr : CondExpr or CondExpr . (2)

CondExpr : CondExpr . and CondExpr (3)

and shift 7 or shift 8 \$end reduce 2

One way to resolve conflict: fix grammar.

#### Use "stratified grammar," as for arithmetic expressions:

 $\begin{array}{l} CondExpr \rightarrow CondTerm \mid CondExpr \mid \mid CondTerm \\ CondTerm \rightarrow CondPrimary \mid CondTerm \&\& \ CondPrimary \\ CondPrimary \rightarrow id \mid ! \ CondPrimary \end{array}$ 

Another way to resolve conflict: precedence declarations.

- Suppose  $t_1$  is the topmost terminal symbol on the stack, and  $t_2$  is the lookahead symbol. Then:
  - If  $t_1$ ,  $t_2$  appear in the same %left declaration, then reduce
  - If  $t_1$ ,  $t_2$  appear in the same %right declaration, then shift
  - If  $t_1$  appears in a declaration before  $t_2$ , shift
  - If  $t_1$  appears in a declaration after  $t_2$ , reduce

#### Use the ambiguous grammar, but add these declarations:

%left or %left and

•  $x \parallel y \&\& z$  is now handled correctly.

#### However, ocamlyacc still reports conflicts. Verbose output:

```
6: shift/reduce conflict (shift 7, reduce 4) on and
6: shift/reduce conflict (shift 8, reduce 4) on or
state 6
CondExpr : CondExpr . or CondExpr (2)
CondExpr : CondExpr . and CondExpr (3)
CondExpr : not CondExpr . (4)
and shift 7
or shift 8
```

\$end reduce 4

Problem is that we didn't resolve ambiguity involving "!".
Add "%nonassoc not" after above two lines.

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### More on conflicts

- Posted supplementary notes discuss four grammars that have conflicts, and how to resolve them.
- All are relevant to the current MP.