## Lecture 9: Bottom-up parsing; ocamlyacc

- Bottom-up parsing
- ocamlyacc

## Top-down vs. bottom-up parsing

Why is top-down called "top-down"?

As we consume tokens, we build a parse tree. At any time, we are filling in the children of a particular non-terminal. As soon as we decide what production to use, we can fill in the tree. In this sense, we are building the tree from the top down.

• Example: 
$$E \rightarrow id T$$
  
 $T \rightarrow \epsilon \mid + E \mid * E$   
Input: x + y \* z

– Typeset by Foil $T_{\!E\!}\!\mathrm{X}$  –

# **Bottom-up parsing**

Bottom-up parsing works by creating small parse trees and joining them together into larger ones.

• Example: 
$$E \rightarrow id T$$
  
 $T \rightarrow \epsilon \mid + E \mid * E$ 

Input: x + y \* z

# How bottom-up parsing works

- Keep a stack of small parse trees. Based on what's in this stack, and the next input token, take one of these actions:
  - Shift: Move lookahead token to stack
  - Reduce  $A \rightarrow \alpha$ : If roots of trees on stack match  $\alpha$ , replace those trees on stack by single tree with root A.
  - Accept: Reduce when non-terminal is goal, look-ahead is eof
  - Reject
- Bottom-up parsing is also called shift-reduce parsing.

## Shift-reduce example 1

• Example:  $L \rightarrow L$ ;  $E \mid E$  $E \rightarrow id$ 

Input: x; y; z

## Shift-reduce example 2

#### • Example: $E \rightarrow E + T \mid T$ $T \rightarrow T * P \mid P$ $P \rightarrow id \mid int$

Input: x + 10 \* y

# Using ocamlyacc

- Input attribute grammar is put in file <grammar>.mly
- Execute ocamlyacc <grammar>.mly
- Produces code for parser in <grammar>.ml and interface (including type declaration for tokens) in <grammar>.mli

## Parser code

- *<grammar>.ml* defines one parsing function per entry point
- Parsing function takes a lexing function (lexer buffer to token) and a lexer buffer as arguments
- Returns semantic attribute of corresponding entry point

# **Example - expression grammar**

In this example, we will take a simple expression grammar and create a parser to parse inputs and produce abstract syntax.

Grammar:

$$\begin{array}{ll} M \rightarrow Exp \; eof \\ Exp \rightarrow Term \; | \; Term \; + \; Exp \; | \; Term \; - \; Exp \\ Term \rightarrow Factor \; | \; Factor \; * \; Term \; | \; Factor \; / \; Term \\ Factor \; \rightarrow \; id \; | \; ( \; Exp \; ) \end{array}$$

Abstract syntax:

```
(* File: expr.ml *)
type expr =
    Plus of expr * expr
    Minus of expr * expr
    Mult of expr * expr
    Div of expr * expr
    Id of string
```

## **Example - lexer**

```
(* File: exprlex.mll *)
let numeric = ['0' - '9']
let letter = ['a' - 'z' 'A' - 'Z']
rule tokenize = parse
   "+" {Plus_token}
   "-" {Minus_token}
  | "*" {Times_token}
  | "/" {Divide_token}
  | "(" {Left_parenthesis}
  | ")" {Right_parenthesis}
  | letter (letter | numeric | "_")* as id {Id_token id}
  | [', ', ' t', ' n'] \{token lexbuf\}
  l eof {EOL}
```

## **Example - parser**

```
(* File: exprparse.mly *)
%{ open Expr
%}
%token <string> Id_token
%token Left_parenthesis Right_parenthesis
%token Times_token Divide_token
%token Plus_token Minus_token
%token EOL
%start main
%type <expr> main
%%
```

# Example - parser (exprparse.mly)

expr:

term	{\$1}
term Plus_token expr	{Plus(\$1,\$3)}
term Minus_token expr	{Minus(\$1,\$3)}

term:

factor {\$1}
| factor Times\_token term {Mult(\$1,\$3)}
| factor Divide\_token term {Div(\$1,\$3)}

factor:

Id\_token {Id \$1}
| Left\_parenthesis expr Right\_parenthesis {\$2}

main:

| expr EOL {\$1}

## **Example - using parser**

```
# #use "expr.ml";;
...
# #use "exprparse.ml";;
...
# #use "exprlex.ml";;
...
# let test s =
let lexbuf = Lexing.from_string(s^"\n") in
main tokenize lexbuf;;
# test "a + b";;
- : expr = Plus(Id "a",Id "b")
```

# ocamlyacc Input

#### File format:

```
%{
	<header>
%}
	<declarations>
%%
	<rules>
%%
	<trailer>
```

## ocamlyacc <header>

- Contains arbitrary Ocaml code
- Typcially used to give types and functions needed for the semantic actions of rules and to give specialized error recovery
- May be omitted
- < footer> similar. Possibly used to call parser

## **ocamlyacc** <declarations>

- %token symbol ... symbol
  Declare given symbols as tokens
- %token <type> symbol ... symbol
  Declare given symbols as token constructors, taking an argument of type type
- %start symbol ... symbol
  Declare given symbols as entry points; functions of same names in <grammar>.ml

## **ocamlyacc** <declarations>

- %type <type> symbol ... symbol
  Specify type of attributes for given symbols. Mandatory for start symbol
- %left symbol ... symbol
- %right symbol ... symbol
- %nonassoc symbol ... symbol
   Associate precedences and associativities to given symbols.
   Same line, same precedence; earlier line, lower precedence
   (broadest scope)

## ocamlyacc <rules>

```
nonterminal:
symbol ... symbol { semantic_action }
| ...
| symbol ... symbol { semantic_action }
;
```

- Semantic actions are arbitrary Ocaml expressions
- Must be of same type as declared (or inferred) for nonterminal
- Access values semantic attributes of symbols by position: \$1 for first symbol, \$2 for second, etc.

# Friday's class

- Big question: how to choose whether to shift or reduce.
- ocamlyacc uses a method called LALR(1) to construct tables which say what action to take.
- There are times when there is no good way to make this decision. (ocamlyacc will reject grammar and give an error message.) In bottom-up parsing, these are called *conflicts*. There are two types: shift/reduce and reduce/reduce.
- As with top-down parsing, these problems can sometimes be resolved by modifying the grammar.
- On Friday, will discuss these conflicts and give some advice on how to resolve them.

# **MP 6**

- MP 6 starts with a grammar embedded in an incomplete ocamlyacc specification. You will need to finish the spec:
  - Remove "extended BNF" productions ocamlyacc cannot handle them
  - Resolve grammar conflicts
  - Fill in actions so as to produce ASTs.