Programming Languages and Compilers (CS 421)

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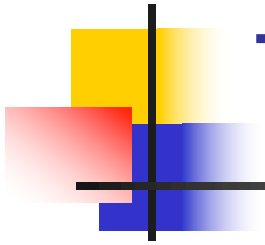
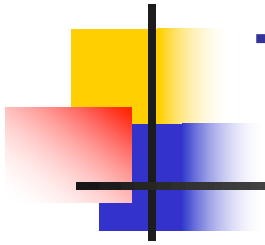
Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha



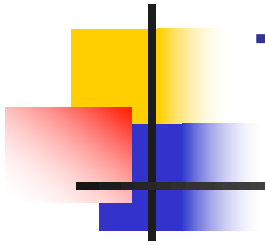
Two Problems

- Type checking
  - Question: Does exp. e have type τ in env Γ?
  - Answer: Yes / No
  - Method: Type derivation
- Typability
  - Question Does exp. e have some type in env. Γ? If so, what is it?
  - Answer: Type τ / error
  - Method: Type inference

Type Inference - Outline



- Begin by assigning a type variable as the type of the whole expression
- Decompose the expression into component expressions
- Use typing rules to generate constraints on components and whole
- Recursively find substitution that solves typing judgment of first subcomponent
- Apply substitution to next subcomponent and find substitution solving it; compose with first, etc.
- Apply comp of all substitution to orig. type var. to get answer



Type Inference - Example

- What type can we give to

```
(fun x -> fun f -> f (f x))
```

- Start with a type variable and then look at the way the term is constructed

Type Inference - Example

- First approximate:

$$\{\ \} \vdash (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha$$

- Second approximate: use fun rule

$$\frac{\{x : \beta\} \vdash (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}{\{\ \} \vdash (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}$$

- Remember constraint $\alpha \equiv (\beta \rightarrow \gamma)$

Type Inference - Example

- Third approximate: use fun rule

$$\frac{\dfrac{\{f : \delta\ ;\ x : \beta\} \vdash f\ (f\ x) : \varepsilon}{\{x : \beta\} \vdash (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}}{\{\ \} \vdash (\text{fun } x \text{ -> fun } f \text{ -> } f(f\ x)) : \alpha}$$

- $\alpha \equiv (\beta \rightarrow \gamma)$; $\gamma \equiv (\delta \rightarrow \varepsilon)$

Type Inference - Example

- Fourth approximate: use app rule

$$\dfrac{\dfrac{\dfrac{\{f{:}\delta;\ x{:}\beta\} \vdash f : \varphi \to \varepsilon \quad \{f{:}\delta;\ x{:}\beta\} \vdash f\ x : \varphi}{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \vdash (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}}{\{\ \} \vdash (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}$$

- $\alpha \equiv (\beta \to \gamma);\ \gamma \equiv (\delta \to \varepsilon)$

Type Inference - Example

- Fifth approximate: use var rule, get constraint $\delta \equiv \varphi \rightarrow \varepsilon$, Solve with same
- Apply to next sub-proof

$$\dfrac{\dfrac{\dfrac{\dfrac{\{f{:}\delta;\ x{:}\beta\} \vdash f : \varphi \rightarrow \varepsilon \quad \{f{:}\delta;\ x{:}\beta\} \vdash f\ x : \varphi}{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \vdash (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}}{\{\ \} \vdash (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}$$

- $\alpha \equiv (\beta \rightarrow \gamma);\ \gamma \equiv (\delta \rightarrow \varepsilon)$

Type Inference - Example

- Current subst: $\{\delta \equiv \varphi \rightarrow \varepsilon\}$

$$\frac{\dots \quad \{f:\varphi \rightarrow \varepsilon;\ x:\beta\} \mid\!- f\ x : \varphi}{\dfrac{\dfrac{\{f : \delta\ ;\ x : \beta\} \mid\!- (f\ (f\ x)) : \varepsilon}{\{x : \beta\} \mid\!- (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}}{\{\ \} \mid\!- (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}}$$

- $\alpha \equiv (\beta \rightarrow \gamma);\ \gamma \equiv (\delta \rightarrow \varepsilon)$

Type Inference - Example

- Current subst: $\{\delta \equiv \varphi \rightarrow \varepsilon\}$ Use App Rule

$$\frac{\{f:\varphi\rightarrow\varepsilon;\ x:\beta\}\vdash f:\zeta\rightarrow\varphi \quad \{f:\varphi\rightarrow\varepsilon;\ x:\beta\}\vdash x:\zeta}{\dots \quad \{f:\varphi \rightarrow \varepsilon;\ x:\beta\}\vdash f\ x : \varphi}$$

$$\frac{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}{\{x : \beta\} \vdash (\text{fun f -> f (f x)}) : \gamma}$$

$$\{\ \} \vdash (\text{fun x -> fun f -> f (f x)}) : \alpha$$

- $\alpha \equiv (\beta \rightarrow \gamma);\ \gamma \equiv (\delta \rightarrow \varepsilon)$

Type Inference - Example

- Current subst: $\{\delta \equiv \varphi \rightarrow \varepsilon\}$
- Var rule: Solve $\zeta \rightarrow \varphi \equiv \varphi \rightarrow \varepsilon$ Unification

$$\frac{\{f:\varphi\rightarrow\varepsilon;\ x:\beta\} \vdash f:\zeta\rightarrow\varphi \quad \{f:\varphi\rightarrow\varepsilon;\ x:\beta\} \vdash x:\zeta}{\ldots \quad \{f:\varphi \rightarrow \varepsilon;\ x:\beta\} \vdash f\ x : \varphi}$$

$$\frac{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}{\{x : \beta\} \vdash (\text{fun f -> f (f x)}) : \gamma}$$

$$\{\ \} \vdash (\text{fun x -> fun f -> f (f x)}) : \alpha$$

- $\alpha \equiv (\beta \rightarrow \gamma);\ \gamma \equiv (\delta \rightarrow \varepsilon)$

Type Inference - Example

- Current subst: $\{\zeta \equiv \varepsilon,\ \varphi \equiv \varepsilon\} \circ \{\delta \equiv \varphi \to \varepsilon\}$
- Var rule: Solve $\zeta \to \varphi \equiv \varphi \to \varepsilon$ Unification

$$\dfrac{\dfrac{\{f:\varphi\to\varepsilon;\ x:\beta\} \vdash f:\zeta\to\varphi \quad \{f:\varphi\to\varepsilon;\ x:\beta\} \vdash x:\zeta}{\ldots \quad \{f:\varphi \to \varepsilon;\ x:\beta\} \vdash f\ x : \varphi}}{\dfrac{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}{\dfrac{\{x : \beta\} \vdash (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}{\{\ \} \vdash (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}}}$$

- $\alpha \equiv (\beta \to \gamma);\ \gamma \equiv (\delta \to \varepsilon)$

Type Inference - Example

- Current subst: $\{\zeta \equiv \varepsilon,\ \varphi \equiv \varepsilon,\ \delta \equiv \varepsilon \to \varepsilon\}$
- Apply to next sub-proof

$$\dfrac{\dfrac{\dfrac{\dfrac{\ldots \qquad \{f:\varepsilon \to \varepsilon;\ x:\beta\} \mid\!- x:\varepsilon}{\ldots \quad \{f:\varphi \to \varepsilon;\ x:\beta\} \mid\!- f\ x : \varphi}}{\{f : \delta\ ;\ x : \beta\} \mid\!- (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \mid\!- (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}}{\{\ \} \mid\!- (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}$$

- $\alpha \equiv (\beta \to \gamma);\ \gamma \equiv (\delta \to \varepsilon)$

Type Inference - Example

- Current subst: $\{\zeta \equiv \varepsilon,\ \varphi \equiv \varepsilon,\ \delta \equiv \varepsilon \to \varepsilon\}$
- Var rule: $\varepsilon \equiv \beta$

$$\dfrac{\dfrac{\dfrac{\dfrac{\dots \qquad \{f:\varepsilon \to \varepsilon;\ x:\beta\} \mid\!- x:\varepsilon}{\dots \quad \{f:\varphi \to \varepsilon;\ x:\beta\} \mid\!- f\ x : \varphi}}{\{f : \delta\ ;\ x : \beta\} \mid\!- (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \mid\!- (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}}{\{\ \} \mid\!- (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}$$

- $\alpha \equiv (\beta \to \gamma);\ \gamma \equiv (\delta \to \varepsilon)$

Type Inference - Example

- Current subst: $\{\varepsilon \equiv \beta\} o \{\zeta \equiv \varepsilon,\ \varphi \equiv \varepsilon,\ \delta \equiv \varepsilon \to \varepsilon\}$
- Solves subproof; return one layer

$$\dfrac{\dfrac{\dfrac{\dfrac{}{\{f:\varepsilon \to \varepsilon;\ x:\beta\} \vdash x:\varepsilon} \quad \ldots}{\{f:\varphi \to \varepsilon;\ x:\beta\} \vdash f\ x : \varphi} \quad \ldots}{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \vdash (\text{fun f -> f (f x)}) : \gamma}$$

$$\{\ \} \vdash (\text{fun x -> fun f -> f (f x)}) : \alpha$$

- $\alpha \equiv (\beta \to \gamma);\ \gamma \equiv (\delta \to \varepsilon)$

Type Inference - Example

- Current subst: $\{\varepsilon \equiv \beta,\ \zeta \equiv \beta,\ \varphi \equiv \beta,\ \delta \equiv \beta \rightarrow \beta\}$
- Solves this subproof; return one layer

$$\dfrac{\dfrac{\dfrac{\dfrac{\dots \quad \{f:\varphi \rightarrow \varepsilon;\ x:\beta\} \vdash f\ x : \varphi}{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \vdash (\text{fun } f \text{ -> } f\ (f\ x)) : \gamma}}{\{\ \} \vdash (\text{fun } x \text{ -> fun } f \text{ -> } f\ (f\ x)) : \alpha}}{\dots}$$

- $\alpha \equiv (\beta \rightarrow \gamma);\ \gamma \equiv (\delta \rightarrow \varepsilon)$

Type Inference - Example

- Current subst: $\{\varepsilon \equiv \beta,\ \zeta \equiv \beta,\ \varphi \equiv \beta,\ \delta \equiv \beta \rightarrow \beta\}$
- Need to satisfy constraint $\gamma \equiv (\delta \rightarrow \varepsilon)$, given subst, becomes: $\gamma \equiv ((\beta \rightarrow \beta) \rightarrow \beta)$

$$\dfrac{\dfrac{\dfrac{\dots}{\{f : \delta\ ;\ x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \vdash (\text{fun f -> f (f x)}) : \gamma}}{\{\ \} \vdash (\text{fun x -> fun f -> f (f x)}) : \alpha}$$

- $\alpha \equiv (\beta \rightarrow \gamma);\ \gamma \equiv (\delta \rightarrow \varepsilon)$

Type Inference - Example

- Current subst:

$\{\gamma \equiv ((\beta \to \beta) \to \beta), \varepsilon \equiv \beta, \zeta \equiv \beta, \varphi \equiv \beta, \delta \equiv \beta \to \beta\}$

- Solves subproof; return one layer

$$\frac{\dfrac{\dfrac{\ldots}{\{f : \delta ; x : \beta\} \vdash (f\ (f\ x)) : \varepsilon}}{\{x : \beta\} \vdash (\text{fun } f \to f\ (f\ x)) : \gamma}}{\{\ \} \vdash (\text{fun } x \to \text{fun } f \to f\ (f\ x)) : \alpha}$$

- $\alpha \equiv (\beta \to \gamma)$; $\gamma \equiv (\delta \to \varepsilon)$

Type Inference - Example

- Current subst:

$\{\gamma \equiv ((\beta \to \beta) \to \beta), \varepsilon \equiv \beta, \zeta \equiv \beta, \varphi \equiv \beta, \delta \equiv \beta \to \beta\}$

- Need to satisfy constraint $\alpha \equiv (\beta \to \gamma)$ given subst: $\alpha \equiv (\beta \to ((\beta \to \beta) \to \beta))$

$$\dfrac{\dfrac{\ldots}{\{x : \beta\} \vdash (\text{fun f -> f (f x)}) : \gamma}}{\{\,\} \vdash (\text{fun x -> fun f -> f (f x)}) : \alpha}$$

- $\alpha \equiv (\beta \to \gamma)$;

Type Inference - Example

- Current subst:

$\{\alpha \equiv (\beta \rightarrow ((\beta \rightarrow \beta) \rightarrow \beta)),$

$\gamma \equiv ((\beta \rightarrow \beta) \rightarrow \beta), \varepsilon \equiv \beta, \zeta \equiv \beta, \varphi \equiv \beta, \delta \equiv \beta \rightarrow \beta\}$

- Solves subproof; return on layer

$$\frac{\dfrac{}{\{x : \beta\} \vdash (\text{fun f -> f (f x)}) : \gamma}}{\{\,\} \vdash (\text{fun x -> fun f -> f (f x)}) : \alpha}$$

Type Inference - Example

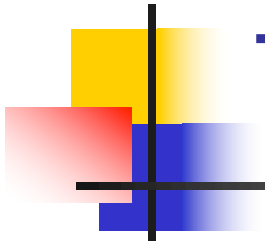
- Current subst:

$\{\alpha \equiv (\beta \rightarrow ((\beta \rightarrow \beta) \rightarrow \beta)),$

$\gamma \equiv ((\beta \rightarrow \beta) \rightarrow \beta), \varepsilon \equiv \beta, \zeta \equiv \beta, \varphi \equiv \beta, \delta \equiv \beta \rightarrow \beta\}$

- Done: $\alpha \equiv (\beta \rightarrow ((\beta \rightarrow \beta) \rightarrow \beta))$

{ } |- (fun x -> fun f -> f (f x)) : α



Type Inference Algorithm

Let infer (Γ, *e*, τ) = σ

- Γ is a typing environment (giving polymorphic types to expression variables)
- *e* is an expression
- τ is a type (with type variables),
- σ is a substitution of types for type variables
- Idea: σ is substitution solving the constraints on type variables necessary for Γ |- *e* : τ
- Should have $\sigma(\Gamma)$ |- *e* : $\sigma(\tau)$ valid

Type Inference Algorithm

infer (Γ, *exp*, τ) =

- Case *exp* of
  - Var *v* --> return Unify{$\tau \equiv$ freshInstance(Γ(*v*))}
    - Replace all quantified type vars by fresh ones
  - Const *c* --> return Unify{$\tau \equiv$ freshInstance φ } where Γ |- *c* : φ by the constant rules
  - fun *x* -> *e* -->
    - Let α, β be fresh variables
    - Let σ = infer ({x: α} + Γ, *e*, β)
    - Return Unify({$\sigma(\tau) \equiv \sigma(\alpha \rightarrow \beta)$}) $\circ$ σ

Example of inference with Var Rule

Instance {'a -> 'w} ('w a fresh variable)

{x: All 'a. ('a * 'b) list, y:All. 'b}|- x : (int * string) list

freshInstance(All 'a. ('a * 'b) list) = ('w * 'b) list

Unify {((int*string)list = ('w * 'b) list)} ={'w -> int, 'b -> string}

After substitution:

Instance {'a -> int}

{x:All 'a. ('a * string) list, y:All. string}|- x:(int * string) list

Type Inference Algorithm (cont)

- Case *exp* of
  - App (e_1 e_2) -->
    - Let α be a fresh variable
    - Let $\sigma_1 = \text{infer}(\Gamma, e_1, \alpha \rightarrow \tau)$
    - Let $\sigma_2 = \text{infer}(\sigma_1(\Gamma), e_2, \sigma_1(\alpha))$
    - Return $\sigma_2 \circ \sigma_1$

Type Inference Algorithm (cont)

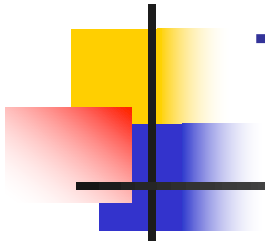
- Case *exp* of
  - If e_1 then e_2 else e_3 -->
    - Let σ_1 = infer(Γ, e_1, bool)
    - Let σ_2 = infer($\sigma_1\Gamma$, e_2, $\sigma_1(\tau)$)
    - Let σ_3 = infer($\sigma_2 o \sigma_1(\Gamma)$,$e_3$,$\sigma_2 o \sigma_1(\tau)$)
    - Return $\sigma_3 \circ \sigma_2 \circ \sigma_1$

Type Inference Algorithm (cont)

- Case *exp* of
  - let $x = e_1$ in e_2 -->
    - Let α be a fresh variable
    - Let $\sigma_1 = \text{infer}(\Gamma, e_1, \alpha)$
    - Let $\sigma_2 =$ $\text{infer}(\{x:\text{GEN}(\sigma_1(\Gamma), \sigma_1(\alpha))\} + \sigma_1(\Gamma), e_2, \sigma_1(\tau))$
    - Return $\sigma_2 \circ \sigma_1$

Type Inference Algorithm (cont)

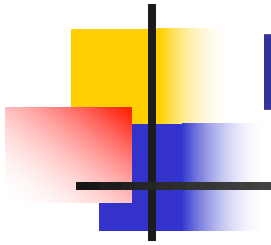
- Case *exp* of
  - let rec $x = e_1$ in e_2 -->
    - Let α be a fresh variable
    - Let σ_1 = infer({x: α} + Γ, e_1, α)
    - Let σ_2 = infer({x:GEN($\sigma_1(\Gamma)$,$\sigma_1(\alpha)$)} + $\sigma_1(\Gamma)$}, e_2, $\sigma_1(\tau)$)
    - Return $\sigma_2 \circ \sigma_1$



Type Inference Algorithm (cont)

- To infer a type, introduce type_of
- Let α be a fresh variable
- type_of (Γ, *e*) =
  - Let σ = infer (Γ, *e*, α)
  - Return σ (α)

- Need an algorithm for Unif



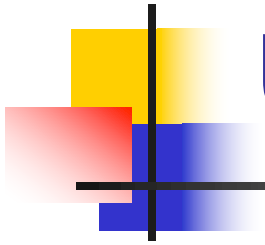
Background for Unification

- Terms made from constructors and variables (for the simple first order case)
- Constructors may be applied to arguments (other terms) to make new terms
- Variables and constructors with no arguments are base cases
- Constructors applied to different number of arguments (arity) considered different
- Substitution of terms for variables

Simple Implementation Background

```
type term = Variable of string
                    | Const of (string * term list)
let x = Variable "'a";;   let tm = Const ("2",[]);;

let rec subst var_name residue term =
    match term with Variable name ->
        if var_name = name then residue else term
      | Const (c, tys) ->
        Const (c, List.map (subst var_name residue)
                           tys);;
```



Unification Problem

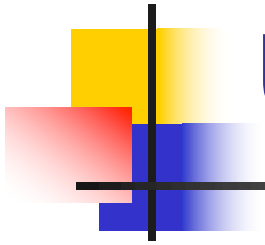
Given a set of pairs of terms ("equations")

$$\{(s_1, t_1), (s_2, t_2), \ldots, (s_n, t_n)\}$$

(the *unification problem*) does there exist a substitution σ (the *unification solution*) of terms for variables such that

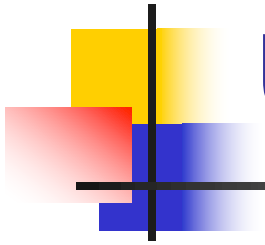
$$\sigma(s_i) = \sigma(t_i),$$

for all $i = 1, \ldots, n$?



Uses for Unification

- Type Inference and type checking
- Pattern matching as in OCaml
  - Can use a simplified version of algorithm
- Logic Programming - Prolog
- Simple parsing



Unification Algorithm

- Let $S = \{(s_1 = t_1), (s_2 = t_2), \ldots, (s_n = t_n)\}$ be a unification problem.

- Case $S = \{\ \}$: $Unif(S)$ = Identity function (i.e., no substitution)

- Case $S = \{(s, t)\} \cup S'$: Four main steps

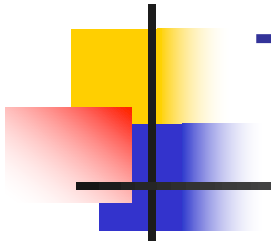
Unification Algorithm

- Delete: if $s = t$ (they are the same term) then $Unif(S) = Unif(S')$
- Decompose: if $s = f(q_1, \ldots, q_m)$ and $t = f(r_1, \ldots, r_m)$ (same f, same m!), then

  $Unif(S) = Unif(\{(q_1, r_1), \ldots, (q_m, r_m)\} \cup S')$
- Orient: if $t = x$ is a variable, and s is not a variable, $Unif(S) = Unif(\{(x = s)\} \cup S')$

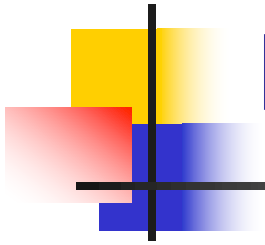
Unification Algorithm

- Eliminate: if $s = x$ is a variable, and x does not occur in t (the occurs check), then
  - Let $\varphi = \{x \rightarrow t\}$
    - $\text{Unif}(S) = \text{Unif}(\varphi(S')) \circ \{x \rightarrow t\}$
  - Let $\psi = \text{Unif}(\varphi(S'))$
  - $\text{Unif}(S) = \{x \rightarrow \psi(t)\} \circ \psi$
    - Note: $\{x \rightarrow a\} \circ \{y \rightarrow b\} = \{y \rightarrow (\{x \rightarrow a\}(b))\} \circ \{x \rightarrow a\}$ if y not in a



Tricks for Efficient Unification

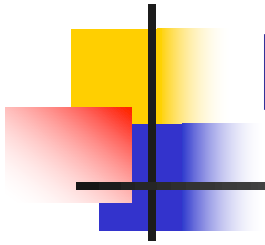
- Don't return substitution, rather do it incrementally
- Make substitution be constant time
  - Requires implementation of terms to use mutable structures (or possibly lazy structures)
  - We won't discuss these



Example

- x,y,z variables, f,g constructors

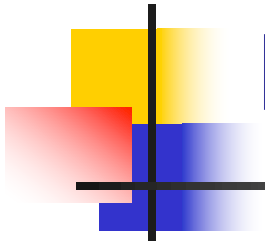
- Unify {(f(x) = f(g(f(z),y))), (g(y,y) = x)} = ?



Example

- x,y,z variables, f,g constructors
- S = {(f(x) = f(g(f(z),y))), (g(y,y) = x)} is nonempty

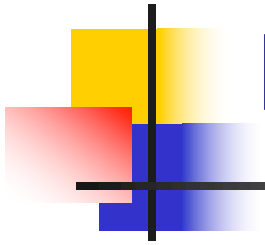
- Unify {(f(x) = f(g(f(z),y))), (g(y,y) = x)} = ?



Example

- x,y,z variables, f,g constructors
- Pick a pair: (g(y,y) = x)

- Unify {(f(x) = f(g(f(z),y))), (g(y,y) = x)} = ?

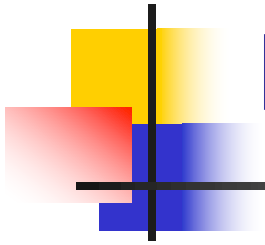


Example

- x,y,z variables, f,g constructors
- Pick a pair: (g(y,y)) = x)
- Orient: (x = g(y,y))

- Unify {(f(x) = f(g(f(z),y))), (g(y,y) = x)} =
  Unify {(f(x) = f(g(f(z),y))), (x = g(y,y))}

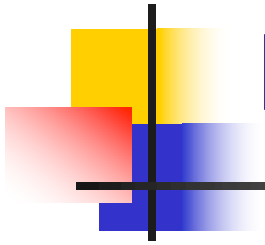
by Orient



Example

- x,y,z variables, f,g constructors

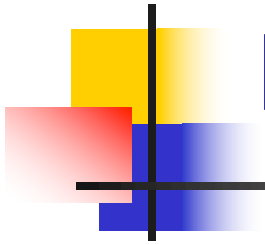
- Unify {(f(x) = f(g(f(z),y))), (x = g(y,y))} = ?



Example

- x,y,z variables, f,g constructors
- {(f(x) = f(g(f(z),y))), (x = g(y,y))} is non-empty

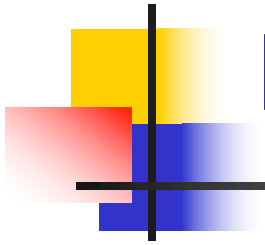
- Unify {(f(x) = f(g(f(z),y))), (x = g(y,y))} = ?



Example

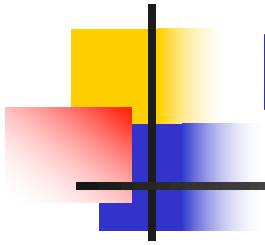
- x,y,z variables, f,g constructors
- Pick a pair: (x = g(y,y))

- Unify {(f(x) = f(g(f(z),y))), (x = g(y,y))} = ?



Example

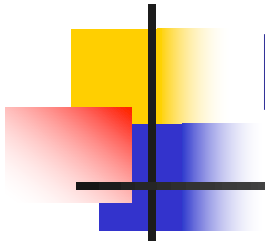
- x,y,z variables, f,g constructors
- Pick a pair: (x = g(y,y))
- Eliminate x with substitution {x→ g(y,y)}
  - Check: x not in g(y,y)
- Unify {(f(x) = f(g(f(z),y))), (x = g(y,y))} = ?



Example

- x,y,z variables, f,g constructors
- Pick a pair: (x = g(y,y))
- Eliminate x with substitution {x→ g(y,y)}

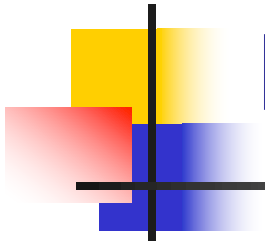
- Unify {(f(x) = f(g(f(z),y))), (x = g(y,y))} =
  Unify {(f(g(y,y)) = f(g(f(z),y)))}
  o {x→ g(y,y)}



Example

- x,y,z variables, f,g constructors

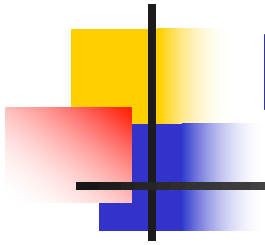
- Unify {(f(g(y,y)) = f(g(f(z),y)))}
  o {x→ g(y,y)} = ?



Example

- x,y,z variables, f,g constructors
- {(f(g(y,y)) = f(g(f(z),y)))} is non-empty

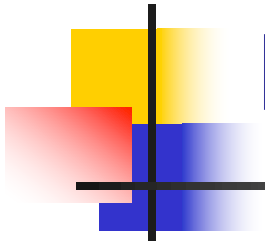
- Unify {(f(g(y,y)) = f(g(f(z),y)))}
  o {x→ g(y,y)} = ?



Example

- x,y,z variables, f,g constructors
- Pick a pair: (f(g(y,y)) = f(g(f(z),y)))

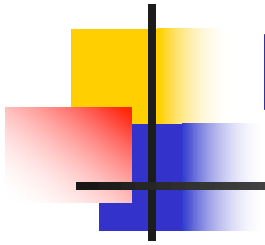
- Unify {(f(g(y,y)) = f(g(f(z),y)))}
  o {x→ g(y,y)} = ?



Example

- x,y,z variables, f,g constructors
- Pick a pair: (f(g(y,y)) = f(g(f(z),y)))
- Decompose:(f(g(y,y)) = f(g(f(z),y))) becomes {(g(y,y) = g(f(z),y))}

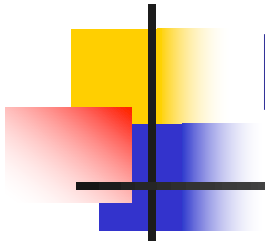
- Unify {(f(g(y,y)) = f(g(f(z),y)))}
  o {x→ g(y,y)} =
  Unify {(g(y,y) = g(f(z),y))} o {x→ g(y,y)}



Example

- x,y,z variables, f,g constructors
- {(g(y,y) = g(f(z),y))} is non-empty

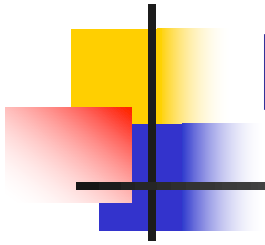
- Unify {(g(y,y) = g(f(z),y))}
  o {x→ g(y,y)} = ?



Example

- x,y,z variables, f,g constructors
- Pick a pair: (g(y,y) = g(f(z),y))

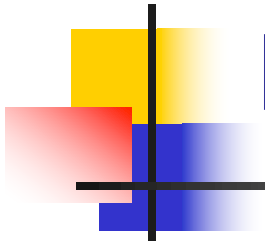
- Unify {(g(y,y) = g(f(z),y))}
  o {x→ g(y,y)} = ?



Example

- x,y,z variables, f,g constructors
- Pick a pair: (f(g(y,y)) = f(g(f(z),y)))
- Decompose: (g(y,y)) = g(f(z),y)) becomes {(y = f(z)); (y = y)}

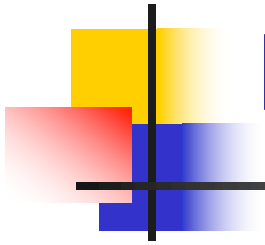
- Unify {(g(y,y) = g(f(z),y))} o {x→ g(y,y)} = Unify {(y = f(z)); (y = y)} o {x→ g(y,y)}



Example

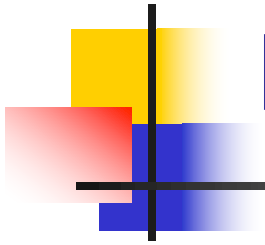
- x,y,z variables, f,g constructors

- Unify {(y = f(z)); (y = y)} o {x→ g(y,y)} = ?



Example

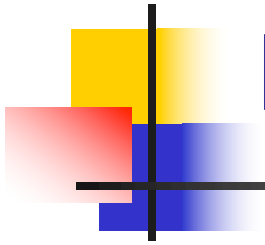
- x,y,z variables, f,g constructors
- {(y = f(z)); (y = y)} o {x→ g(y,y) is non-empty
- Unify {(y = f(z)); (y = y)} o {x→ g(y,y)} = ?



Example

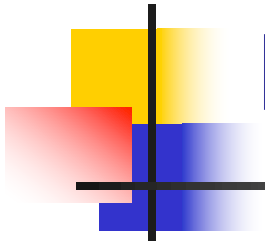
- x,y,z variables, f,g constructors
- Pick a pair: (y = f(z))

- Unify {(y = f(z)); (y = y)} o {x→ g(y,y)} = ?



Example

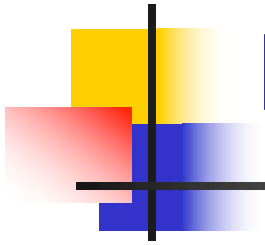
- x,y,z variables, f,g constructors
- Pick a pair: (y = f(z))
- Eliminate y with {y → f(z)}
- Unify {(y = f(z)); (y = y)} o {x→ g(y,y)} =
  Unify {(f(z) = f(z))}
    o ({y → f(z)} o {x→ g(y,y)})=
  Unify {(f(z) = f(z))}
    o {y → f(z); x→ g(f(z), f(z))}



Example

- x,y,z variables, f,g constructors

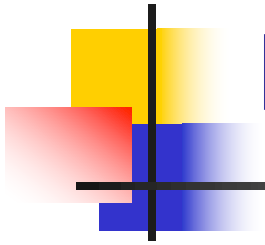
- Unify {(f(z) = f(z))}
  - o {y $\rightarrow$ f(z); x$\rightarrow$ g(f(z), f(z))} = ?



Example

- x,y,z variables, f,g constructors
- {(f(z) = f(z))} is non-empty

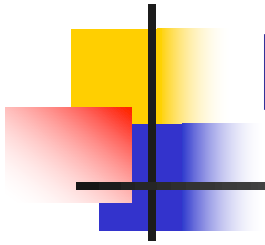
- Unify {(f(z) = f(z))}
  - o {y $\rightarrow$ f(z); x$\rightarrow$ g(f(z), f(z))} = ?



Example

- x,y,z variables, f,g constructors
- Pick a pair: (f(z) = f(z))

- Unify {(f(z) = f(z))}
  - o {y $\rightarrow$ f(z); x$\rightarrow$ g(f(z), f(z))} = ?

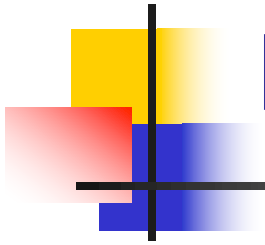


Example

- x,y,z variables, f,g constructors
- Pick a pair: (f(z) = f(z))
- Delete
- Unify {(f(z) = f(z))}

  o {y → f(z); x→ g(f(z), f(z))} =

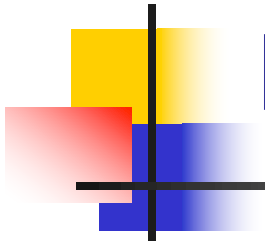
  Unify {} o {y → f(z); x→ g(f(z), f(z))}



Example

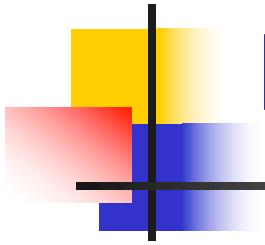
- x,y,z variables, f,g constructors

- Unify {} o {y $\rightarrow$ f(z); x$\rightarrow$ g(f(z), f(z))} = ?



Example

- x,y,z variables, f,g constructors
- {} is empty
- Unify {} = identity function
- Unify {} o {y $\rightarrow$ f(z); x$\rightarrow$ g(f(z), f(z))} =
  {y $\rightarrow$ f(z); x$\rightarrow$ g(f(z), f(z))}



Example

- Unify {(f(x) = f(g(f(z),y))), (g(y,y) = x)} =
  {y → f(z); x→ g(f(z), f(z))}

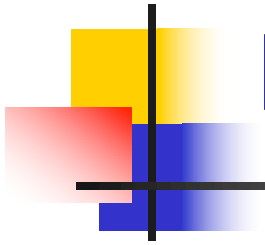
f(x) = f(g(f(z), y))

→ f(g(f(z), f(z))) = f(g(f(z), f(z)))

g(y , y) = x

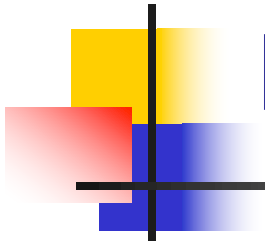
→ g(f(z),f(z)) = g(f(z), f(z))

Example of Failure: Decompose



- Unify{(f(x,g(y)) = f(h(y),x))}
- Decompose: (f(x,g(y)) = f(h(y),x))
- = Unify {(x = h(y)), (g(y) = x)}
- Orient: (g(y) = x)
- = Unify {(x = h(y)), (x = g(y))}
- Eliminate: (x = h(y))
- Unify {(h(y) = g(y))} o {x $\rightarrow$ h(y)}
- No rule to apply! Decompose fails!

Example of Failure: Occurs Check



- Unify{(f(x,g(x)) = f(h(x),x))}
- Decompose: (f(x,g(x)) = f(h(x),x))
- = Unify {(x = h(x)), (g(x) = x)}
- Orient: (g(x) = x)
- = Unify {(x = h(x)), (x = g(x))}
- No rules apply.