

# Hoare Logic 2

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# Hoare Triple

P { ...code... } Q

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$$P[e/x] \{ x := e \} P$$

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$$\frac{P \{ C_1 \} R \quad R \{ C_2 \} Q}{P \{ C_1; C_2 \} Q}$$

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$$\frac{P \wedge b \{ C_1 \} Q \quad P \wedge \neg b \{ C_2 \} Q}{P \{ \text{if } b \text{ then } C_1 \text{ else } C_2 \} Q}$$

# While Rule

$$\frac{P \wedge b \{ C \} P}{P \{ \text{While } b C \} P \wedge \neg b}$$

(P is a **loop invariant**)

# Rule of Consequence

$$\frac{P \rightarrow P' \quad P' \{ C \} Q' \quad Q' \rightarrow Q}{P \{ C \} Q}$$

# Sample Proofs

- sum of n
- fibonacci
- list append
- list reverse
- termination

# Sum of n

$x = 0 \ \& \ y = 0$

$P \equiv x = 1 + \dots + y \wedge y \leq n$

{

While  $y < n$

$y := y + 1;$

$x := x + y$

}

$x = 1 + \dots + n$

$$x = 0 \wedge y = 0 \rightarrow x = 1 + \dots + y \wedge y \leq n \quad \checkmark$$

$$x = 1 + \dots + y \wedge y \leq n \wedge \neg(y < n) \rightarrow x = 1 + \dots + n \quad \checkmark$$

$$x = 1 + \dots + y \wedge y \leq n \wedge y < n \rightarrow ? \quad \checkmark$$

$$x + y + 1 = (1 + \dots + (y + 1)) \wedge y + 1 \leq n$$

$\uparrow \{y := y + 1\} \quad x + y = 1 + \dots + y \wedge y \leq n$   
 $\uparrow \{x := x + y\} \quad x = 1 + \dots + y \wedge y \leq n$

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$$? \quad \{y := y + 1; x := x + y\} \quad x = 1 + \dots + y \wedge y \leq n$$

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$$x = 1 + \dots + y \wedge y \leq n \wedge y < n \quad \{y := y + 1; x := x + y\} \quad x = 1 + \dots + y \wedge y \leq n$$

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$$x = 1 + \dots + y \wedge y \leq n \quad \{\text{While } y < n \dots\} \quad x = 1 + \dots + y \wedge y \leq n \wedge \neg(y < n)$$

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$$x = 0 \wedge y = 0 \quad \{\text{While } \dots\} \quad x = 1 + \dots + n$$

# Fibonacci

```
x = 0 & y = 1 & z = 1 & 1 ≤ n
{
    While z < n      P ≡ y = fib z ∧ x = fib (z-1)
        y := x + y;      ∧ z ≤ n
        x := y - x;
        z := z + 1
}
y = fib n
```

$x = 0 \wedge y = 1 \wedge z = 0 \wedge 1 \leq n \rightarrow y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n$	✓
$y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge \neg(z < n) \rightarrow y = \text{fib } n$	✓
$y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge z < n \rightarrow ?$	✓

$x+y = \text{fib } (z+1) \wedge x+y-x = \text{fib } (z+1-1) \wedge z + 1 \leq n$

$$\begin{array}{ll} \{y := x + y\} & y = \text{fib } (z+1) \wedge y-x = \text{fib } (z+1-1) \wedge z + 1 \leq n \\ \{x := y - x\} & y = \text{fib } (z+1) \wedge x = \text{fib } (z+1-1) \wedge z + 1 \leq n \\ \{z := z + 1\} & y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \end{array}$$

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?       $\{y := x + y; x := y - x; z := z + 1\} \quad y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n$

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$y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge z < n \quad \{y := x + y; x := y - x; z := z + 1\} \quad y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n$

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$y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \quad \{\text{While } z < n \dots\} \quad y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge \neg(z < n)$

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$x = 0 \wedge y = 1 \wedge z = 0 \wedge 1 \leq n \quad \{\text{While } \dots\} \quad y = \text{fib } n$

# List length

$x = \text{lst} \ \& \ y = 0$

$P \equiv \text{len lst} = y + \text{len } x$

{

While  $x \neq []$

$x := \text{tl } x;$

$y := y + 1$

}

$y = \text{len lst}$

$$x = \text{lst} \wedge y = 0 \rightarrow \text{len lst} = y + \text{len } x \quad \checkmark$$

$$\text{len lst} = y + \text{len } x \wedge \neg(x \neq []) \rightarrow y = \text{len lst} \quad \checkmark$$

$$\text{len lst} = y + \text{len } x \wedge x \neq [] \rightarrow ? \quad \checkmark$$

$$\boxed{\text{len lst} = y + 1 + \text{len}(tl \ x)}$$

$$\begin{array}{ll} \uparrow \{x := tl \ x\} & \text{len lst} = y + 1 + \text{len } x \\ \{y := y + 1\} & \text{len lst} = y + \text{len } x \end{array}$$

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$$? \quad \{x := tl \ x; y := y + 1\} \quad \text{len lst} = y + \text{len } x$$

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$$\text{len lst} = y + \text{len } x \wedge x \neq [] \quad \{x := tl \ x; y := y + 1\} \quad \text{len lst} = y + \text{len } x$$

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$$\text{len lst} = y + \text{len } x \quad \{\text{While } x \neq [] \dots\} \quad \text{len lst} = y + \text{len } x \wedge \neg(x \neq [])$$

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$$x = \text{lst} \wedge y = 0 \quad \{\text{While } \dots\} \quad y = \text{len lst}$$

# List reverse

$x = \text{lst} \ \& \ y = [] \quad P \equiv \text{lst} = \text{rev } y @ x$

{

While  $x \neq []$

$y := \text{hd } x :: y;$

$x := \text{tl } x$

}

$y = \text{rev lst}$

$$x = \text{lst} \wedge y = [] \rightarrow \text{lst} = \text{rev } y @ x \quad \checkmark$$

$$\text{lst} = \text{rev } y @ x \wedge \neg(x \neq []) \rightarrow y = \text{rev lst} \quad \checkmark$$

$$\text{lst} = \text{rev } y @ x \wedge x \neq [] \rightarrow ? \quad \checkmark$$

$\text{lst} = \text{rev} (\text{hd } x @ y) @ (\text{tl } x)$   
 $\uparrow \{y := \text{hd } x @ y\} \quad \text{lst} = \text{rev } y @ (\text{tl } x)$   
 $\{x := \text{tl } x\} \quad \text{lst} = \text{rev } y @ x$

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$$? \quad \{y := \text{hd } x @ y; x := \text{tl } x\} \quad \text{lst} = \text{rev } y @ x$$


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$$\text{lst} = \text{rev } y @ x \wedge x \neq [] \quad \{y := \text{hd } x @ y; x := \text{tl } x\} \quad \text{lst} = \text{rev } y @ x$$


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$$\text{lst} = \text{rev } y @ x \quad \{\text{While } x \neq [] \dots\} \quad \text{lst} = \text{rev } y @ x \wedge \neg(x \neq [])$$


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$$x = \text{lst} \wedge y = [] \quad \{\text{While } \dots\} \quad y = \text{rev lst}$$