# Programming Languages and Compilers (CS 421)

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https://courses.engr.illinois.edu/cs421/fa2017/CS421A

Based in part on slides by Mattox Beckman, as updated by Vikram Adve, Gul Agha, and Elsa L Gunter

#### Structural Recursion

- Functions on recursive datatypes (eg lists) tend to be recursive
- Recursion over recursive datatypes generally by structural recursion
  - Recursive calls made to components of structure of the same recursive type
  - Base cases of recursive types stop the recursion of the function

# Structural Recursion: List Example

- Nil case [] is base case
- Cons case recurses on component list xs

#### Forward Recursion

- In Structural Recursion, split input into components and (eventually) recurse on components
- Forward Recursion form of Structural Recursion
- In forward recursion, first call the function recursively on all recursive components, and then build final result from partial results
- Wait until whole structure has been traversed to start building answer

# Forward Recursion: Examples

```
# let rec double up list =
    match list with
         [ ] -> [ ]
       | (x :: xs) -> (x :: x :: double_up xs);;
val double up : 'a list -> 'a list = <fun>
# let rec poor rev list =
  match list with
       [] -> []
     | (x::xs) -> poor_rev xs @ [x];;
val poor rev : 'a list -> 'a list = <fun>
```

How do you write length with forward recursion?

let rec length 1 =

How do you write length with forward recursion?

#### **Functions Over Lists**

```
# let fib5_2 = double_up fib5;;
val fib5_2 : int list = [8; 8; 5; 5; 3; 3; 2; 2;
1; 1; 1; 1]
```

#### **Functions Over Lists**

```
# poor_rev silly;;
- : string list = ["there"; "there"; "hi"; "hi"]
```

#### Your Turn

Write a function odd\_count fr: int list -> int such that it returns the number of odd integers found in the input list. The function is required to use (only) forward recursion (no other form of recursion).

```
# let rec odd_count_fr l =
```

```
# odd_count_fr [1;2;3];;
- : int = 2
```

# An Important Optimization

Normal call

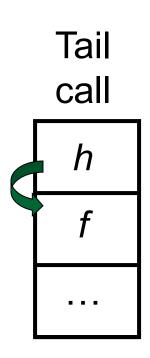
h

g

f

- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if f calls g and g calls h, but calling h is the last thing g does (a tail call)?
  - let f x = (g x) + 1
  - let g x = h (x+1)
  - let h x = ...

# An Important Optimization



- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if f calls g and g calls h, but calling h is the last thing g does (a tail call)?
- Then h can return directly to f instead of g

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#### Tail Recursion

- A recursive program is tail recursive if all recursive calls are tail calls
- Tail recursive programs may be optimized to be implemented as loops, thus removing the function call overhead for the recursive calls
- Tail recursion generally requires extra "accumulator" arguments to pass partial results
  - May require an auxiliary function

# Example of Tail Recursion

```
# let rec prod l =
    match l with [] -> 1
    (x :: rem) \rightarrow x * prod rem;;
val prod : int list -> int = <fun>
# let prod list =
    let rec prod aux 1 acc =
        match l with [] -> acc
        (y :: rest) -> prod_aux rest (acc * y)
(* Uses associativity of multiplication *)
    in prod aux list 1;;
 val prod : int list -> int = <fun>
```

How do you write length with tail recursion?

```
let length 1 =
```

How do you write length with tail recursion?

```
let length l =
  let rec length_aux list n =
```

in

How do you write length with tail recursion?

```
let length l =
    let rec length_aux list n =
    match list with [] ->
     (a :: bs) ->
in
```

How do you write length with tail recursion?

```
let length l =
    let rec length_aux list n =
    match list with [] -> n
     (a :: bs) ->
in
```

How do you write length with tail recursion?

#### Your Turn

Write a function odd\_count\_tr : int list -> int such that it returns the number of odd integers found in the input list. The function is required to use (only) tail recursion (no other form of recursion).

```
# let rec odd_count_tr l =
```

```
# odd_count_tr [1;2;3];;
- : int = 2
```

# Mapping Recursion

 One common form of structural recursion applies a function to each element in the structure

```
# let rec doubleList list = match list with
      [ ] -> [ ]
      | x::xs -> 2 * x :: doubleList xs;;
val doubleList : int list -> int list = <fun>
# doubleList [2;3;4];;
- : int list = [4; 6; 8]
```

# Mapping Functions Over Lists

```
# let rec map f list =
  match list with
    [] -> []
  | (h::t) -> (f h) :: (map f t);;
val map : ('a \rightarrow 'b)\rightarrow 'a list\rightarrow 'b list = \langle fun \rangle
# map plus_two fib5;;
-: int list = [10; 7; 5; 4; 3; 3]
# map (fun x \rightarrow x - 1) fib6;;
: int list = [12; 7; 4; 2; 1; 0; 0]
```

# Mapping Recursion

 Can use the higher-order recursive map function instead of direct recursion

```
# let doubleList list =
    List.map (fun x -> 2 * x) list;;
val doubleList : int list -> int list = <fun>
# doubleList [2;3;4];;
- : int list = [4; 6; 8]
```

Same function, but no rec

#### Your turn now

#### Write a function

```
make_app : (('a -> 'b) * 'a) list -> 'b list
```

that takes a list of function – input pairs and gives the result of applying each function to its argument. Use map, no explicit recursion.

```
let make_app lst =
```

# Folding Recursion

 Another common form "folds" an operation over the elements of the structure

```
# let rec multList list = match list with
    [ ] -> 1
    | x::xs -> x * multList xs;;
val multList : int list -> int = <fun>
# multList [2;4;6];;
- : int = 48
```

Computes (2 \* (4 \* (6 \* 1)))

#### How are the following functions similar?

```
# let rec sumlist list = match list with
     | -> 0
   | x::xs \rightarrow x + sumlist xs;;
# sumlist [2;3;4];;
- : int = 9
# let rec prodlist list = match list with
     [ ] -> 1
   | x::xs -> x * prodlist xs;;
# prodlist [2;3;4];;
-: int = 24
```

```
How are the following functions similar?
# let rec sumlist list = match list with
    | x::xs \rightarrow x + sumlist xs;;
# sumlist [2;3;4];;
                            Base Case
- : int = 9
# let rec prodlist list = match list with
     [ ] -> 1
    x::xs -> x * prodlist xs;;
# prodlist [2;3;4];;
-: int = 24
```

```
How are the following functions similar?
# let rec sumlist list = match list with
      [ ] -> 0
    \mid x::xs \rightarrow x + |sumlist xs|;
# sumlist [2;3;4];;
                            Recursive Call
- : int = 9
# let rec prodlist list/= match list with
      [ ] -> 1
    | x::xs -> x * prodlist xs;
# prodlist [2;3;4];;
- : int = 24
```

```
How are the following functions similar?
# let rec sumlist list = match list with
     [ ] -> 0
   | x::xs \rightarrow x + sumlist xs;
# sumlist [2;3;4];;
                          Head Element
- : int = 9
# let rec prodlist list = match list with
     x::xs -> x * prodlist xs;
# prodlist [2;3;4];;
- : int = 24
```

```
How are the following functions similar?
# let rec sumlist list = match list with
     [ ] -> 0
   x::xs \rightarrow x + sumlist xs;
# sumlist [2;3;4];;
                           Combining Operator
- : int = 9
# let rec prodlist list = match list with
     [ ] -> 1
   x::xs -> x * prodlist xs;
# prodlist [2;3;4];;
-: int = 24
```

## Recursing over lists

```
# let rec fold right f list b =
  match list with
    [] -> b
  (x :: xs) -> f x (fold_right f xs b);;
# fold right
    (fun val init -> val + init)
    [1; 2; 3]
    0;;
- : int = 6
```

## Recursing over lists

```
# let rec fold right f list b =
  match list with
    [] -> b
  (x :: xs) \rightarrow f x (fold right f xs b);;
# fold right
    (fun s -> fun () -> print string s)
    ["hi"; "there"]
    ();;
therehi-: unit = ()
```

## Folding Recursion

- multList folds to the right
- Same as:

```
# let multList list =
    List.fold right
       (fun x -> fun p -> x * p)
       list 1;;
val multList : int list -> int = <fun>
# multList [2;4;6];;
-: int = 48
```

## **Encoding Recursion with Fold**

```
# let rec append list1 list2 = match list1 with
  [ ] -> list2 | x::xs -> x :: append xs list2;;
val append : 'a list -> 'a list -> \'a list = <fun>
              Operation | Recursive call
  Base Case
# let append list1 list2 = 
   fold_right (fun x y -> x :: y) list1 list2;;
val append : 'a list -> 'a list -> 'a list = <fun>
# append [1;2;3] [4;5;6];;
 -: int list = [1; 2; 3; 4; 5; 6]
```

How do you write length with fold\_right, but no explicit recursion?

How do you write length with fold\_right, but no explicit recursion?

How do you write length with fold\_right, but no explicit recursion?

Can you write fold\_right (or fold\_left) with just map? How, or why not?

#### Iterating over lists

```
# let rec fold left f a list =
  match list with
   || -> a
  (x :: xs) -> fold left f (f a x) xs;;
val fold left: ('a -> 'b -> 'a) -> 'a -> 'b list
  -> 'a = <fun>
# fold left
   (fun () -> print_string)
   ["hi"; "there"];;
hithere- : unit = ()
```

# Encoding Tail Recursion with fold\_left

```
# let prod list = let rec prod_aux l acc =
        match 1 with
          -> acc
        | (y :: rest) -> prod_aux rest (acc * y)
       in prod aux list_1;;
    Init Acc Value
                    Recursive Call
                                       Operation
# let prod list
   List.fold left (fun acc y -> acc * y) 1 list;;
# prod [4;5;6];;
 -: int = 120
```

How do you write length with fold\_left, but no explicit recursion?

How do you write length with fold\_left, but no explicit recursion?

# **Folding**

```
# let rec fold left f a list = match list with
     [] -> a
   (x :: xs) -> fold left f (f a x) xs;;
fold_left f a [x_1; x_2;...;x_n] = f(...(f (f a <math>x_1) x_2)...)x_n
# let rec fold_right f list b = match list with
     [] -> b
   | (x :: xs) -> f x (fold_right f xs b);;
fold_right f [x_1; x_2;...;x_n] b = f x_1(f x_2 (...(f x_n b)...))
```

#### Recall

What is its running time?

# Quadratic Time

- Each step of the recursion takes time proportional to input
- Each step of the recursion makes only one recursive call.
- List example:

## Comparison

- poor\_rev [1,2,3] =
- (poor\_rev [2,3]) @ [1] =
- ((poor\_rev [3]) @ [2]) @ [1] =
- (((poor\_rev []) @ [3]) @ [2]) @ [1] =
- (([] @ [3]) @ [2]) @ [1]) =
- ([3] @ [2]) @ [1] =
- **(**3:: ([ ] @ [2])) @ [1] =
- [3,2] @ [1] =
- **3** :: ([2] @ [1]) =
- 3 :: (2:: ([ ] @ [1])) = [3, 2, 1]

#### Tail Recursion - Example

```
# let rec rev aux list revlist =
  match list with
    [ ] -> revlist
  x :: xs -> rev_aux xs (x::revlist);;
val rev aux : 'a list -> 'a list -> 'a list =
  <fun>
# let rev list = rev aux list [ ];;
val rev : 'a list -> 'a list = <fun>
```

What is its running time?

#### Comparison

- rev [1,2,3] =
- rev\_aux [1,2,3] [] =
- rev\_aux [2,3] [1] =
- rev\_aux [3] [2, I] =
- rev\_aux [] [3,2,1] = [3,2,1]

# Folding - Tail Recursion

```
# let rec rev aux list revlist =
  match list with
    [ ] -> revlist
  | x :: xs -> rev_aux xs (x::revlist);;
# let rev list = rev_aux list [ ];;
# let rev list =
        fold left
          (fun 1 -> fun x -> x :: 1) (* comb op *)
                              (* accumulator cell *)
            list
```

# **Folding**

- Can replace recursion by fold\_right in any forward primitive recursive definition
  - Primitive recursive means it only recurses on immediate subcomponents of recursive data structure
- Can replace recursion by fold\_left in any tail primitive recursive definition

## Example of Tail Recursion

```
# let rec app fl x =
    match fl with [] \rightarrow x
     | (f :: rem_fs) -> f (app rem_fs x);;
val app : ('a -> 'a) list -> 'a -> 'a = \langle fun \rangle
# let app fs x =
    let rec app aux fl acc =
        match fl with [] -> acc
        (f :: rem fs) -> app aux rem fs
                                      (fun z -> acc (f z))
    in app_aux fs (fun y -> y) x;;
val app : ('a -> 'a) list -> 'a -> 'a = <fun>
```

# Continuation Passing Style

- A programming technique for all forms of "non-local" control flow:
  - non-local jumps
  - exceptions
  - general conversion of non-tail calls to tail calls
- Essentially it's a higher-order function version of GOTO

#### **Continuations**

- Idea: Use functions to represent the control flow of a program
- Method: Each procedure takes a function as an argument to which to pass its result; outer procedure "returns" no result
- Function receiving the result called a continuation
- Continuation acts as "accumulator" for work still to be done

# Continuation Passing Style

 Writing procedures so that they take a continuation to which to give (pass) the result, and return no result, is called continuation passing style (CPS)

## Example

Simple reporting continuation:

Simple function using a continuation:

# Example of Tail Recursion & CSP

```
\# let app fs x =
    let rec app_aux fl acc=
        match fl with
          [] -> acc
        (f:: rem_fs) -> app_aux rem_fs
                                   (fun z -> acc (f z))
    in app_aux fs (fun y -> y) x;;
val app : ('a -> 'a) list -> 'a -> 'a = <fun>
# let rec appk fl x k =
    match fl with
      [] \rightarrow k x
    | (f :: rem_fs) -> appk rem_fs x (fun z -> k (f z));
hval appk : ('a -> 'a) list -> 'a -> ('a -> 'b) -> 'b
  9/11/2018
                                                       62
```

## Example of Tail Recursion & CSP

```
# let rec appk fl x k =
    match fl with
    [] -> k x
    | (f :: rem_fs) -> appk rem_fs x (fun z -> k (f z));;
# appk [(fun x->x+1); (fun x -> x*5)] 2 (fun x->x);;
- : int = 11
```

# Continuation Passing Style

 A compilation technique to implement non-local control flow, especially useful in interpreters.

 A formalization of non-local control flow in denotational semantics

 Possible intermediate state in compiling functional code

#### **Terms**

- A function is in **Direct Style** when it returns its result back to the caller.
- A Tail Call occurs when a function returns the result of another function call without any more computations (e.g. tail recursion)
- A function is in Continuation Passing Style when it passes its result to another function.
  - Instead of returning the result to the caller, we pass it forward to another function.

# Continuation Passing Style

 A compilation technique to implement non-local control flow, especially useful in interpreters.

 A formalization of non-local control flow in denotational semantics

 Possible intermediate state in compiling functional code

## Example

Simple reporting continuation:

Simple function using a continuation:

#### Simple Functions Taking Continuations

- Given a primitive operation, can convert it to pass its result forward to a continuation
- Examples:

# Nesting Continuations

```
# let add three x y z = x + y + z;;
val add three : int -> int -> int -> int = <fun>
# let add three x y z = let p = x + y in p + z;
val add three : int -> int -> int -> int = <fun>
# let add_three_k x y z k =
   addk x y (fun p -> addk p z k );;
val add three k : int -> int -> int -> (int -> 'a
  -> 'a = <fun>
```