Programming Languages and Compilers (CS 421)

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Based in part on slides by Mattox Beckman, as updated by Vikram Adve, Gul Agha, and Elsa L Gunter

Tuples as Values

```
//  ρ0 = {c → 4, a → 1, b → 5}
# let s = (5,"hi",3.2);;
val s : int * string * float = (5, "hi", 3.2)

//  ρ = {s → (5, "hi", 3.2), c → 4, a → 1, b → 5}
```


Pattern Matching with Tuples

```
//  ρ = {s → (5, "hi", 3.2), a → 1, b → 5, c → 4}

# let (a,b,c) = s;;      (* (a,b,c) is a pattern *)
val a : int = 5
val b : string = "hi"
val c : float = 3.2

# let (a, _, _) = s;;
val a : int = 5

# let x = 2, 9.3;;     (* tuples don't require parens in Ocaml *)
val x : int * float = (2, 9.3)
```


Nested Tuples

```
#  (*Tuples can be nested *)
# let d = ((1,4,62),("bye",15),73.95);;
val d : (int * int * int) * (string * int) * float =
  ((1, 4, 62), ("bye", 15), 73.95)

#  (*Patterns can be nested *)
# let (p, (st,_), _) = d;;
                    (* _ matches all, binds nothing *)
val p : int * int * int = (1, 4, 62)
val st : string = "bye"
```


Functions on tuples

```
# let plus_pair (n,m) = n + m;;
val plus_pair : int * int -> int = <fun>
# plus_pair (3,4);;
- : int = 7
# let twice x = (x,x);;
val twice : 'a -> 'a * 'a = <fun>
# twice 3;;
- : int * int = (3, 3)
# twice "hi";;
- : string * string = ("hi", "hi")
```


Save the Environment!

- A **closure** is a pair of an environment and an association of a sequence of variables (the input variables) with an expression (the function body), written:

  < (v1,…,vn) → exp, ρ >

- Where ρ is the environment in effect when the function is defined (for a simple function)

Closure for plus_pair

- Assume ρ_{plus_pair} was the environment just before plus_pair defined and recall
  - `let plus_pair (n,m) = n + m;;`
- Closure for fun (n,m) -> n + m:

  $<(n,m) \rightarrow n + m, \rho_{plus\_pair}>$

- Environment just after plus_pair defined:

  $\{plus\_pair \rightarrow <(n,m) \rightarrow n + m, \rho_{plus\_pair}>\} + \rho_{plus\_pair}$

Like set union! (but subtle differences, see slide 17)

Functions with more than one argument

```
# let add_three x y z = x + y + z;;
val add_three : int -> int -> int -> int = <fun>

# let t = add_three 6 3 2;;
val t : int = 11

# let add_three =
   fun x -> (fun y -> (fun z -> x + y + z));;
val add_three : int -> int -> int -> int = <fun>
```

Again, first syntactic sugar for second

Curried vs Uncurried

- Recall

```
# let add_three u v w = u + v + w;;
val add_three : int -> int -> int -> int = <fun>
```

- How does it differ from

```
# let add_triple (u,v,w) = u + v + w;;
val add_triple : int * int * int -> int = <fun>
```

- add_three is **curried**;
- add_triple is **uncurried**

Curried vs Uncurried

```
# add_three 6 3 2;;
- : int = 11
# add_triple (6,3,2);;
- : int = 11
# add_triple 5 4;;
Characters 0-10:  add_triple 5 4;;
                   ^^^^^^^^^^
This function is applied to too many arguments,
maybe you forgot a `;'
# fun x -> add_triple (5,4,x);;
: int -> int = <fun>
```


Partial application of functions

```
let add_three x y z = x + y + z;;
```

```
# let h = add_three 5 4;;
val h : int -> int = <fun>

# h 3;;
- : int = 12

# h 7;;
- : int = 16
```

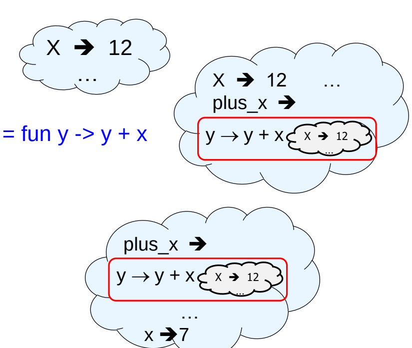
Partial application also called *sectioning*

Recall: let plus_x = fun y -> y + x

let x = 12

let plus_x = fun y -> y + x

let x = 7



Closure for plus_x

- When plus_x was defined, had environment:

  $\rho_{plus\_x} = \{\ldots, x \rightarrow 12, \ldots\}$
- Recall: let plus_x y = y + x

  is really let plus_x = fun y -> y + x
- Closure for fun y -> y + x:

  $<y \rightarrow y + x, \rho_{plus\_x}>$
- Environment just after plus_x defined:

  $\{plus\_x \rightarrow <y \rightarrow y + x, \rho_{plus\_x}>\} + \rho_{plus\_x}$

Evaluation

- Running Ocaml source:
  - Parse the program to detect each expression
  - Keep an internal environment at each time step
  - For each expression, interpret the program using the (mathematical) function **Eval**
  - Nice property of Ocaml: **everything is a declaration or an expression**!
- How does Eval (expression, environment) work:
  - Evaluation uses a starting environment ρ
  - Define the rules for evaluating declarations, constants, arithmetic expressions, function applications…

Evaluating Declarations

- Evaluation uses a starting environment ρ
- To evaluate a (simple) declaration let x = e
  - **Evaluate** expression e in ρ to value v
  - **Update** ρ with the mapping from x to v: $\{x \rightarrow v\} + \rho$

Definition of + on environments!

- **Update**: $\rho_1 + \rho_2$ has all the bindings in ρ_1 and all those in ρ_2 that are not rebound in ρ_1

It is not commutative!

```
  {x → 2, y → 3, a → "hi"}
+ {y → 100, b → 6}
= {x → 2, y → 3, a → "hi", b → 6}
```


Evaluating Declarations

- Evaluation uses a starting environment ρ
- To evaluate a (simple) declaration let x = e
  - **Evaluate** expression e in ρ to value v
  - **Update** ρ with the mapping from x to v: $\{x \rightarrow v\} + \rho$

Warm-up: we evaluate this case:

```
ρ = { x → 2 }
let y = 2*x+1;;
ρ' = { x → 2; y → 5 }
```


Evaluating Expressions **(Rules)**

- Evaluation uses an environment ρ
- **A constant** evaluates to itself
- To evaluate a **variable** x, look it up in ρ i.e., use $\rho(x)$
- To evaluate tuples, evaluate each tuple element
- To evaluate **uses of +,** _ , etc, first eval the arguments, then do the operation
- To evaluate a **local declaration**: let x = e1 in e2
  - Evaluate e1 to v, evaluate e2 using $\{x \rightarrow v\} + \rho$
- **Function application (f x) --** see next slide

Evaluation of Function Application with Closures

Function **defined** as: let f $(x_1, \ldots x_n)$= body

Function **application**: f $(e_1, \ldots, e_n)$;

Let us define Eval(f $(e_1, \ldots, e_n)$, ρ):

- In the environment ρ, evaluate the left term (f) to closure, i.e., $c = <(x_1,\ldots,x_n) \rightarrow body, \rho^*>$
- Evaluate the arguments in the application $e_1 \ldots e_n$ to their values $v_1,\ldots,v_n$ in the environment ρ
- **Call helper function App(Closure, Value) to evaluate** the function body (body) in the environment ρ^*
  - Conjoin the mapping of the arguments to values with the environment ρ^*

    $\rho' = \{x_1 \rightarrow v_1,\ldots, x_n \rightarrow v_n\} + \rho^*$
  - The App then calls Eval again for the expressions in body in the env. ρ'

Evaluation of Application of plus_x;;

- Have environment:

  ρ = {plus_x → <y → y + x, ρ_{plus_x} >, … , y → 3, …}

  where ρ_{plus_x} = {x → 12, … , y → 24, …}

  - Eval (plus_x y, ρ) rewrites to
  - App (Eval(plus_x, ρ) , Eval(y, ρ)) rewrites to
  - App (<y → y + x, ρ_{plus_x} >, 3) rewrites to
  - Eval (y + x, {y → 3} +ρ_{plus_x}) rewrites to
  - Eval (3 + 12 , ρ_{plus_x}) = 15

Evaluation of Application of plus_pair

- Assume environment

ρ = {x → 3, … ,
    plus_pair →<(n,m) →n + m, ρ_{plus_pair}>} + ρ_{plus_pair}

- Eval (plus_pair (4,x), ρ)=
- App (Eval (plus_pair, ρ), Eval ((4,x), ρ)) =
- App (<(n,m) →n + m, ρ_{plus_pair}>, (4,3)) =
- Eval (n + m, {n -> 4, m -> 3} + ρ_{plus_pair}) =
- Eval (4 + 3, {n -> 4, m -> 3} + ρ_{plus_pair}) = 7

Closure question

- If we start in an empty environment, and we execute:

```
let f = fun n -> n + 5;;
(* 0 *)
let pair_map g (n,m) = (g n, g m);;
let f = pair_map f;;
let a = f (4,6);;
```

What is the environment at (* 0 *)?

Answer

ρ_{start} = {}

let f = fun n -> n + 5;;

ρ_0 = {f → <n → n + 5, { }>}

Closure question

- If we start in an empty environment, and we execute:

```
let f = fun n -> n + 5;;
let pair_map g (n,m) = (g n, g m);;
(* 1 *)
let f = pair_map f;;
let a = f (4,6);;
```

What is the environment at (* 1 *)?

Answer

```
ρ0 = {f → <n → n + 5, { }>}
let pair_map g (n,m) = (g n, g m);;

ρ1 = {
      f → <n → n + 5, { }>,
      pair_map →
          <g → (fun (n,m) -> (g n, g m)),
               {f → <n → n + 5, { }>}
          >
    }
```

Closure question

- If we start in an empty environment, and we execute:

```
let f = fun n -> n + 5;;
let pair_map g (n,m) = (g n, g m);;
let f = pair_map f;;
(* 2 *)
let a = f (4,6);;
```

What is the environment at (* 2 *)?

Evaluate pair_map f

```
ρ0 = {f → <n → n + 5, { }>}
ρ1 = {f → <n → n + 5, { }>,
      pair_map →
                  <g → (fun (n,m) -> (g n, g m)),
                   {f → <n → n + 5, { }>}>}
let f = pair_map f;;
```


Evaluate pair_map f

```
ρ0 = {f → <n → n + 5, { }>}
ρ1 = {f → <n → n + 5, { }>,
      pair_map →
                  <g → (fun (n,m) -> (g n, g m)),
                   {f → <n → n + 5, { }>}>}
let f = pair_map f;;
```

Eval(pair_map f, ρ_1) =

Evaluate pair_map f

```
ρ0 = {f → <n → n + 5, { }>}
ρ1 = {f → <n → n + 5, { }>,
      pair_map →
                  <g → (fun (n,m) -> (g n, g m)),
                   {f → <n → n + 5, { }>}>}
let f = pair_map f;;
```

Eval(pair_map f, ρ_1) =
App (<g→fun (n,m) -> (g n, g m), ρ_0>, <n → n + 5, { }>) =

Evaluate pair_map f

```
ρ0 = {f → <n → n + 5, { }>}
ρ1 = {f → <n → n + 5, { }>,
      pair_map →
                  <g → (fun (n,m) -> (g n, g m)),
                   {f → <n → n + 5, { }>}>}
let f = pair_map f;;
```

Eval(pair_map f, ρ_1) =
App (<g→fun (n,m) -> (g n, g m), ρ_0>, <n → n + 5, { }>) =

Eval(fun (n,m)->(g n, g m), {g→<n→n + 5, { }>}+ρ_0) =
<(n,m) →(g n, g m), {g→<n→n + 5, { }>}+ρ_0> =
<(n,m) →(g n, g m), {g→<n→n + 5, { }>, f→<n→n + 5, { }>}

Answer

```
ρ0 = {f → <n → n + 5, { }>}
ρ1 = {f → <n → n + 5, { }>,
      pair_map →
                  <g → (fun (n,m) -> (g n, g m)),
                   {f → <n → n + 5, { }>}>}
let f = pair_map f;;
ρ2 = {f → <(n,m) →(g n, g m),
                  {g → <n → n + 5, { }>,
                   f → <n → n + 5, { }>}>,
         pair_map → <g →  fun (n,m) -> (g n, g m),
                           {f → <n → n + 5, { }>}
                    >
      }
```

Closure question

- If we start in an empty environment, and we execute:

```
let f = fun n -> n + 5;;
let pair_map g (n,m) = (g n, g m);;
let f = pair_map f;;
let a = f (4,6);;
(* 3 *)
```

What is the environment at (* 3 *)?

Final Evalution?

```
ρ2 = {f → <(n,m) →(g n, g m),
                {g → <n → n + 5, { }>,
                 f → <n → n + 5, { }>}>,
         pair_map → <g →  fun (n,m) -> (g n, g m),
                         {f → <n → n + 5, { }>}
                    >
     }
let a = f (4,6);;
```


Evaluate f (4,6);;

```
ρ2 = {f → <(n,m) →(g n, g m),
                {g → <n → n + 5, { }>,
                 f → <n → n + 5, { }>}>,
         pair_map → <g →  fun (n,m) -> (g n, g m),
                         {f → <n → n + 5, { }>}
                    >
     }
let a = f (4,6);;

Eval(f (4,6), ρ2) =
```


Evaluate f (4,6);;

```
ρ2 = {f → <(n,m) →(g n, g m),
                {g → <n → n + 5, { }>,
                 f → <n → n + 5, { }>}>,
         pair_map → <g →  fun (n,m) -> (g n, g m),
                         {f → <n → n + 5, { }>}
                    >
     }
let a = f (4,6);;

Eval(f (4,6), ρ2) =
App(<(n,m) →(g n, g m), {g → <n → n + 5, { }>,
                         f → <n → n + 5, { }>}>
    ,
    (4,6)) =
```


Evaluate f (4,6);;

```
App(<(n,m) →(g n, g m), {g → <n → n + 5, { }>,
                          f → <n → n + 5, { }>}>,
    (4,6)) =

Eval((g n, g m), {n → 4, m → 6} +
                 {g → <n → n + 5, { }>,
                  f → <n → n + 5, { }>}) =
(App(<n → n + 5, { }>, 4),
 App (<n → n + 5, { }>, 6)) =
```


Evaluate f (4,6);;

```
(App(<n → n + 5, { }>, 4),
 App (<n → n + 5, { }>, 6)) =

(Eval(n + 5, {n → 4} + { }),
 Eval(n + 5, {n → 6} + { })) =

(Eval(4 + 5, {n → 4} + { }),
 Eval(6 + 5, {n → 6} + { }))  = (9, 11)
```

Finally:

```
ρ3 = {a -> (9, 11)} + ρ2
```

Functions as arguments

```
# let thrice f x = f (f (f x));;
val thrice : ('a -> 'a) -> 'a -> 'a = <fun>

# let g = thrice plus_two;;   (* plus_two x is x+2 *)
val g : int -> int = <fun>

# g 4;;
- : int = 10

# thrice (fun s -> "Hi! " ^ s) "Good-bye!";;
- : string = "Hi! Hi! Hi! Good-bye!"
```


Higher Order Functions

- A function is *higher-order* if it takes a function as an argument or returns one as a result
- Example:

```
# let compose f g = fun x -> f (g x);;
val compose : ('a -> 'b) -> ('c -> 'a) -> 'c -> 'b
  = <fun>
```

- The type ('a -> 'b) -> ('c -> 'a) -> 'c -> 'b is a higher order type because of ('a -> 'b) and ('c -> 'a) and -> 'c -> 'b

Thrice

- Recall:

```
# let thrice f x = f (f (f x));;
val thrice : ('a -> 'a) -> 'a -> 'a = <fun>
```

- How do you write thrice with compose?

```
# let thrice f = compose f (compose f f);;
val thrice : ('a -> 'a) -> 'a -> 'a = <fun>
```


Lambda Lifting

```
# (+)
- : int -> int -> int = <fun>

# let add_two = (+) (print_string "test\n"; 2);;

# let add2 =      (* lambda lifted *)
    fun x -> (+) (print_string "test\n"; 2) x;;
```


Lambda Lifting

- You must remember the rules for evaluation when you use partial application

```
# let add_two = (+) (print_string "test\n"; 2);;
test
val add_two : int -> int = <fun>

# let add2 =      (* lambda lifted *)
    fun x -> (+) (print_string "test\n"; 2) x;;
val add2 : int -> int = <fun>
```


Lambda Lifting

```
# thrice add_two 5;;
- : int = 11

# thrice add2 5;;
test
test
test
- : int = 11
```

- Lambda lifting delayed the evaluation of the argument to (+) until the second argument was supplied

Reminder: Pattern Matching with Tuples

```
# let (a,b,c) = s;;       (* (a,b,c) is a pattern *)
val a : int = 5
val b : string = "hi"
val c : float = 3.2

# let (a, _, _) = s;;
val a : int = 5

#  (*Patterns can be nested *)
# let (p, (st,_), _) = d;;
                  (* _ matches all, binds nothing *)
val p : int * int * int = (1, 4, 62)
val st : string = "bye"
```


Match Expressions

```
# let triple_to_pair triple =

  match triple with
    (0, x, y) -> (x, y)
  | (x, 0, y) -> (x, y)
  | (x, y, _) -> (x, y)
```

- Each clause: pattern on left, expression on right
- Each x, y has scope of only its clause
- Use first matching clause

```
val triple_to_pair : int * int * int -> int * int
  = <fun>
```


Recursive Functions

```
# let rec factorial n =
    if n = 0 then 1
    else n * factorial (n - 1);;
  val factorial : int -> int = <fun>

# factorial 5;;
- : int = 120

# (* rec  is needed for recursive function
  declarations *)
```


Recursion Example

Compute n^2 recursively using:
$$n^2 = (2 * n - 1) + (n - 1)^2$$

```
# let rec nthsq n =          (* rec for recursion *)
   match n with      (* pattern matching for cases *)
     0 -> 0                  (* base case *)
   | n -> (2 * n -1)         (* recursive case *)
           + nthsq (n -1);;  (* recursive call *)

val nthsq : int -> int = <fun>

# nthsq 3;;
- : int = 9
```

Structure of recursion similar to inductive proof

Recursion and Induction

```
# let rec nthsq n =
    match n with
              0 -> 0   (*Base case!*)
            | n -> (2 * n - 1) + nthsq (n - 1) ;;
```

- Base case is the last case; it stops the computation
- Recursive call must be to arguments that are somehow smaller - must progress to base case
- **if or match must contain the base case (!!!)**
  - Failure of selecting base case **will** cause **non-termination**
  - But the program will crash because it exhausts the stack!

Lists

- First example of a recursive datatype (aka algebraic datatype)
- Unlike tuples, lists are homogeneous in type (all elements same type)

Lists

- List can take one of two forms:
  - **Empty list**, written [ ]
  - **Non-empty list**, written x :: xs
    - x is head element,
    - xs is tail list, :: called "cons"
- How we typically write them (syntactic sugar):
  - **[x]** == x :: [ ]
  - **[x1; x2; …; xn]** == x1 :: x2 :: … :: xn :: [ ]

Lists

```
# let fib5 = [8;5;3;2;1;1];;
val fib5 : int list = [8; 5; 3; 2; 1; 1]

# let fib6 = 13 :: fib5;;
val fib6 : int list = [13; 8; 5; 3; 2; 1; 1]

# (8::5::3::2::1::1::[ ]) = fib5;;
- : bool = true

# fib5 @ fib6;;
- : int list =
        [8; 5; 3; 2; 1; 1; 13; 8; 5; 3; 2; 1; 1]
```


Lists are Homogeneous

```
# let bad_list = [1; 3.2; 7];;
Characters 19-22:
  let bad_list = [1; 3.2; 7];;
                     ^^^
```

This expression has type float but is here used with type int

Question

- Which one of these lists is invalid?

1. [2; 3; 4; 6]
2. [2,3; 4,5; 6,7]
3. [(2.3,4); (3.2,5); (6,7.2)] ← **3 is invalid because of the last pair**
4. [["hi"; "there"]; ["wahcha"]; [ ]; ["doin"]]

Functions Over Lists

```
# let rec double_up list =
    match list with
        [ ] -> [ ]  (* pattern before ->,
                       expression after *)
      | (x :: xs) -> (x :: x :: double_up xs);;
val double_up : 'a list -> 'a list = <fun>

(* fib5 = [8;5;3;2;1;1] *)
# let fib5_2 = double_up fib5;;
val fib5_2 : int list = [8; 8; 5; 5; 3; 3; 2; 2;
  1; 1; 1; 1]
```


Functions Over Lists

```
# let silly = double_up ["hi"; "there"];;
val silly : string list = ["hi"; "hi"; "there";
  "there"]

# let rec poor_rev list =
  match list
  with [] -> []
     | (x::xs) -> poor_rev xs @ [x];;
val poor_rev : 'a list -> 'a list = <fun>

# poor_rev silly;;
- : string list = ["there"; "there"; "hi"; "hi"]
```

Question: Length of list

- Problem: write code for the length of the list
  - How to start?

```
let length l =
```


Question: Length of list

- Problem: write code for the length of the list
  - How to start?

```
let rec length l =
    match l with
```


Question: Length of list

- Problem: write code for the length of the list
  - What patterns should we match against?

```
let rec length l =
    match l with
```


Question: Length of list

- Problem: write code for the length of the list
  - What patterns should we match against?

```
let rec length l =
    match l with [] ->
     | (a :: bs) ->
```


Question: Length of list

- Problem: write code for the length of the list
  - What result do we give when l is empty?

```
let rec length l =
    match l with [] -> 0
     | (a :: bs) ->
```


Question: Length of list

- Problem: write code for the length of the list
  - What result do we give when l is not empty?

```
let rec length l =
    match l with [] -> 0
     | (a :: bs) ->
```

Question: Length of list

- Problem: write code for the length of the list
  - What result do we give when l is not empty?

```
let rec length l =
    match l with [] -> 0
     | (a :: bs) -> 1 + length bs
```


Same Length

- How can we efficiently answer if two lists have the same length?

Tactics:

- First list is empty: then true if second list is empty else false
- First list in not empty: then if second list empty return false, or otherwise compare whether the sublists (after the first element) have the same length

Same Length

- How can we efficiently answer if two lists have the same length?

```
let rec same_length list1 list2 =
  match list1 with
    [] -> (
          match list2 with [] -> true
                       | (y::ys) -> false
        )
  | (x::xs) -> (
          match list2 with [] -> false
                       | (y::ys) -> same_length xs ys
        )
```


Functions Over Lists

```
# let rec map f list =
  match list with
    [] -> []
  | (h::t) -> (f h) :: (map f t);;
val map : ('a -> 'b) -> 'a list -> 'b list = <fun>

# map plus_two fib5;;
- : int list = [10; 7; 5; 4; 3; 3]

# map (fun x -> x - 1) fib6;;
: int list = [12; 7; 4; 2; 1; 0; 0]
```


Iterating over lists

```
# let rec fold_left f a list =
  match list with
    [] -> a
  | (x :: xs) -> fold_left f (f a x) xs;;
val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list
  -> 'a = <fun>

# fold_left
   (fun () -> print_string)
   ()
   ["hi"; "there"];;
hithere- : unit = ()
```


Iterating over lists

```
# let rec fold_right f list b =
  match list with
    [] -> b
  | (x :: xs) -> f x (fold_right f xs b);;
val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b
  -> 'b = <fun>

# fold_right
    (fun s -> fun () -> print_string s)
    ["hi"; "there"]
    ();;
therehi- : unit = ()
```

Structural Recursion

- **Functions on recursive datatypes (eg lists) tend to be recursive**
- Recursion over recursive datatypes generally by **structural recursion**
  - Recursive calls made to components of structure of the same recursive type
  - Base cases of recursive types stop the recursion of the function

Structural Recursion : List Example

```
# let rec length list =
  match list with
     [] -> 0                    (* Nil case *)
   | x :: xs -> 1 + length xs;;  (* Cons case *)
val length : 'a list -> int = <fun>

# length [5; 4; 3; 2];;
- : int = 4
```

- Nil case [ ] is base case
- Cons case recurses on component list xs

Forward Recursion

- In **Structural Recursion**, split input into components and (eventually) recurse
- **Forward Recursion** is a form of Structural Recursion
- In forward recursion, first call the function recursively on all recursive components, and then build final result from partial results
- Wait until whole structure has been traversed to start building answer

Forward Recursion: Examples

```
# let rec double_up list =
    match list
    with [ ] -> [ ]
       | (x :: xs) -> (x :: x :: double_up xs);;
val double_up : 'a list -> 'a list = <fun>

# let rec poor_rev list =
  match list
  with [] -> []
     | (x::xs) -> poor_rev xs @ [x];;
val poor_rev : 'a list -> 'a list = <fun>
```


Encoding Recursion with Fold

```
# let rec append list1 list2 = match list1 with
  [ ] -> list2 | x::xs -> x :: append xs list2;;
val append : 'a list -> 'a list -> 'a list = <fun>

# append [1;2;3] [4;5;6];;
 - : int list = [1; 2; 3; 4; 5; 6]

# let append_alt list1 list2 =
   fold_right (fun x y -> x :: y) list1 list2;;
val append_alt : 'a list -> 'a list -> 'a list = <fun>
```


Mapping Recursion

- One common form of structural recursion applies a function to each element in the structure

```
# let rec doubleList list = match list
   with [ ] -> [ ]
   | x::xs -> 2 * x :: doubleList xs;;
val doubleList : int list -> int list = <fun>

# doubleList [2;3;4];;
- : int list = [4; 6; 8]
```

Mapping Recursion

- Can use the higher-order recursive map function instead of direct recursion

```
# let doubleList list =
    List.map (fun x -> 2 * x) list;;
val doubleList : int list -> int list = <fun>

# doubleList [2;3;4];;
- : int list = [4; 6; 8]
```

- Same function, but no recursion

Folding Recursion

- Another common form "folds" an operation over the elements of the structure

```
# let rec multList list = match list
  with [ ] -> 1
  | x::xs -> x * multList xs;;
val multList : int list -> int = <fun>

# multList [2;4;6];;
- : int = 48
```

- Computes (2 * (4 * (6 * 1)))

Folding Recursion

- multList folds to the right
- Same as:

```
# let multList list =
    List.fold_right
    (fun x -> fun p -> x * p)
    list 1;;
val multList : int list -> int = <fun>

# multList [2;4;6];;
- : int = 48
```


How long will it take?

Common big-O times:

- Constant time $O(1)$
  - input size doesn't matter
- Linear time $O(n)$
  - 2x input size ⇒ 2x time
- Quadratic time $O(n^2)$
  - 3x input size ⇒ 9x time
- Exponential time $O(2^n)$
  - Input size n+1 ⇒ 2x time

Linear Time

- Expect most list operations to take linear time $O(n)$
- Each step of the recursion can be done in constant time
- Each step makes only one recursive call
- List example: multList, append
- Integer example: factorial

Quadratic Time

- Each step of the recursion takes time proportional to input
- Each step of the recursion makes only one recursive call.
- List example:

```
# let rec poor_rev list =
  match list
  with [] -> []
     | (x::xs) -> poor_rev xs @ [x];;
val poor_rev : 'a list -> 'a list = <fun>
```

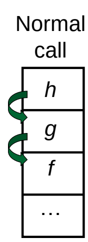
Exponential running time

- Hideous running times on input of any size
- Each step of recursion takes constant time
- Each recursion makes two recursive calls
- Easy to write naïve code that is exponential for functions that can be linear

Exponential running time

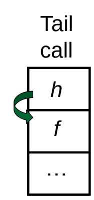
```
# let rec naiveFib n = match n
  with 0 -> 0
  | 1 -> 1
  | _ -> naiveFib (n-1) + naiveFib (n-2);;
val naiveFib : int -> int = <fun>
```


An Important Optimization



- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if *f* calls *g* and *g* calls *h*, but calling *h* is the last thing *g* does (a *tail call*)?

An Important Optimization



- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if *f* calls *g* and *g* calls *h*, but calling *h* is the last thing *g* does (a *tail call*)?
- Then *h* can return directly to *f* instead of *g*

Tail Recursion

- A recursive program is tail recursive if all recursive calls are tail calls
- Tail recursive programs may be optimized to be implemented as loops, thus removing the function call overhead for the recursive calls
- Tail recursion generally requires extra "accumulator" arguments to pass partial results
  - May require an auxiliary function

Tail Recursion - Example

```
# let rec rev_aux list revlist =
  match list with [ ] -> revlist
  | x :: xs -> rev_aux xs (x::revlist);;
val rev_aux : 'a list -> 'a list -> 'a list =
  <fun>

# let rev list = rev_aux list [ ];;
val rev : 'a list -> 'a list = <fun>
```

- What is its running time?

Folding Functions over Lists

How are the following functions similar?

```
# let rec sumlist list = match list with
  [ ] -> 0 | x::xs -> x + sumlist xs;;
val sumlist : int list -> int = <fun>

# sumlist [2;3;4];;
- : int = 9

# let rec prodlist list = match list with
  [ ] -> 1 | x::xs -> x * prodlist xs;;
val prodlist : int list -> int = <fun>

# prodlist [2;3;4];;
- : int = 24
```

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Folding

```
# let rec fold_left f a list = match list
  with [] -> a | (x :: xs) -> fold_left f (f a x) xs;;
val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a
  = <fun>
```

fold_left f a [x_1; x_2;…;x_n] = f(…(f (f a x_1) x_2)…)x_n

```
# let rec fold_right f list b = match list
  with [ ] -> b | (x :: xs) -> f x (fold_right f xs b);;
val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b
  = <fun>
```

fold_right f [x_1; x_2;…;x_n] b = f x_1(f x_2 (…(f x_n b)…))

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Folding - Forward Recursion

```
# let sumlist list = fold_right (+) list 0;;
val sumlist : int list -> int = <fun>

# sumlist [2;3;4];;
- : int = 9

# let prodlist list = fold_right ( * ) list 1;;
val prodlist : int list -> int = <fun>

# prodlist [2;3;4];;
- : int = 24
```

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Folding - Tail Recursion

```
- # let rev list =
-         fold_left
-           (fun l -> fun x -> x :: l) //comb op
            []             //accumulator cell
            list
```

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Folding

- Can replace recursion by fold_right in any forward primitive recursive definition
  - Primitive recursive means it only recurses on immediate subcomponents of recursive data structure
- Can replace recursion by fold_left in any tail primitive recursive definition

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