

Midterm I Review

Feb 21, Monday 7-9 pm

(over DFA, regular, NFA, CFG (up to lec 8))
But no TM.

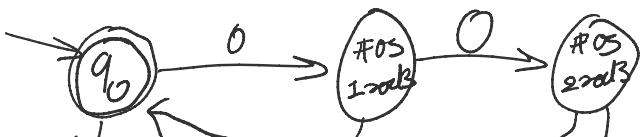
#11. All strings containing at least two 1s and at least one 0.

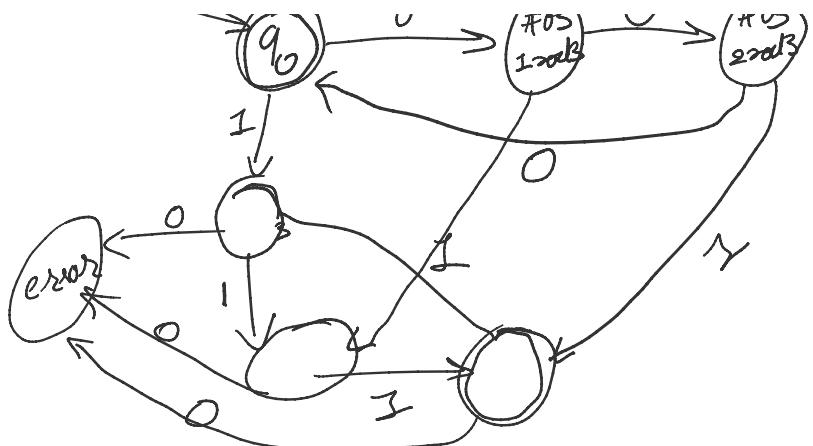
$$0 \sim 1 \sim 1 \sim \text{ OR } \underbrace{1 \dots 1}_{\text{only 1 zero}} 0 \sim 1 \sim \text{ OR}$$

$$\overbrace{0^* 1 (0+1)^* 1 (0+1)^*}^{0^*} + 1 \underbrace{0^* 1}_{1^*} (0+1)^* + 1 1^* 0 (0+1)^*$$

#26. The set of all strings in $0^* 1^*$ whose length is divisible by 3.

$$\underbrace{\epsilon, 000 \ 111,}_{(\underbrace{000})^* \cdot (\underbrace{111})^*} \underbrace{011,}_{\substack{\#0s \text{ is 0 mod 3} \\ \#1s \text{ is 1 mod 3}}} \underbrace{000 \ 001 \ 111,}_{\substack{\#0s \text{ is 1 mod 3} \\ \#1s \text{ is 2 mod 3}}} + \underbrace{(000)^* 011 (\underbrace{111})^*}_{\substack{\#0s \text{ is 1 mod 3} \\ \#1s \text{ is 2 mod 3}}} + \underbrace{(000)^* 001 (\underbrace{111})^*}_{\substack{\#0s \text{ is 2 mod 3} \\ \#1s \text{ is 1 mod 3}}}$$





#35. $L = \{0^n F_n \mid n \geq 0\}$, where F_n is the nth Fibonacci number, defined recursively as follows:

$$F_0 = 0, F_1 = 1, F_n = F_{n-1} + F_{n-2} \quad \forall n \geq 1$$

$$L = \left\{ 0^{F_n} \mid n \geq 0 \right\}$$

$$F_0 = 0$$

$$F_1 = 1$$

$$F_n = F_{n-1} + F_{n-2} \quad \forall n \geq 1$$

Fooling set

$$F = \left\{ 0^{F_i} \mid i \geq 3 \right\}$$

Given any two strings $x, y \in F$ $x \neq y$
 $x = 0^{F_i}, y = 0^{F_j}$ wlog $j > i$

$$\text{Pick } w = 0^{F_{i+1}}$$

$$\text{s.t. } xw = 0^{F_i} \cdot 0^{F_{i+1}} = 0^{F_i + F_{i+1}} = 0^{F_{i+2}} \in L$$

$$yw = 0^j \cdot 0^{F_{i+1}} = 0^{F_j + F_{i+1}}$$

$$F_j < F_j + F_{i+1} \leq F_j + F_{i+1} = F_{i+2}$$

$\left(\because j > i \right)$

$$F_0 = 0 \quad F_1 = 1 \quad \underbrace{F_2 = 1}_{i=2} \quad \underbrace{F_3 = 2}_{j=3} \quad F_4 = 3 \dots$$

#45. $L = \{www \mid w \in \{0,1\}^*\}$

$$F = \{0^i 0^i \mid i \geq 1\} \quad 000000$$

$$x, y \in F, \quad x \neq y$$

$$\text{Given, } x = 0^{i_1} 0^{i_2}$$

$$\text{Pick } z = 0^{i_1}$$

$$xz = \underbrace{0^{i_1}}_w \underbrace{0^{i_1}}_w \underbrace{0^{i_2}}_w \in L$$

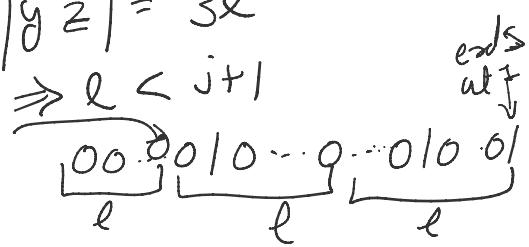
$$yz = \underbrace{0^{i_1}}_w \underbrace{0^{i_1}}_w \underbrace{0^{i_2}}_w \notin L$$

$$y = 0^j 0^j \quad j > i$$

$$|yz| = 3l$$

$$\Rightarrow l < j+1$$

ends at 0



$$L = \left\{ www \mid \frac{x, w \in \{0,1\}^*}{e} \right\} = \{0,1\}^*$$

Regular. ($\because w = e$)

#48. The set of all strings in $\{0, 1\}^*$ such that in every non-empty prefix, the number of 0s is greater than the number of 1s.

- any one*
60. All strings that satisfy ~~all~~ of the following:
- the number of 0s is even
 - the number of 1s is divisible by 3
 - the total length is divisible by 5

DFA $M = (\emptyset, \Sigma, q_0, \delta, A)$

where
 $\emptyset = \{(i, j, k) \mid i \in \{0, 1\}, j \in \{0, 1, 2\}, k \in \{0, 1, 2, 3, 4\}\}$
 length words
 $\#0's \equiv 0 \pmod{2}$
 $\#1's \equiv 0 \pmod{3}$

$$q_0 = (0, 0, 0)$$

$$\delta((i, j, k), 0) = ((i+1) \pmod{2}, j, (k+1) \pmod{5})$$

$$\delta((i, j, k), 1) = (i, (j+1) \pmod{3}, (k+1) \pmod{5})$$

$$A = \{(0, 0, 0)\}$$

$$A' = \{(0, j, k) \mid j \in \{0, 1, 2\}, k \in \{0, 1, 2, 3, 4\}\} \cup$$

$$\{(i, 0, k) \mid \dots\}$$

$$\{(i, j, 0) \mid \dots\}$$

$\} \cup \{ \}$

89. $\{w^{\#0^{\#(0,w)}} \mid w \in \{0, 1\}^*\}$ (write CFG)

11011 # 0

01011 # 00

1001011 # 000

S → 0 S 0 | 1 S | #

71. $\text{Even}(L) = \{ \text{evens}(w) \mid w \in L \}$

If L is regular then $\text{Even}(L)$ is regular.

$\text{evens}(w) = \epsilon$ if $w = \epsilon$

$= \text{odds}(w)$ if $w = ax$, $a \in \Sigma$
 $x \in \Gamma$

$\cdot \text{odds}(w) = \epsilon$ if $w = \epsilon$

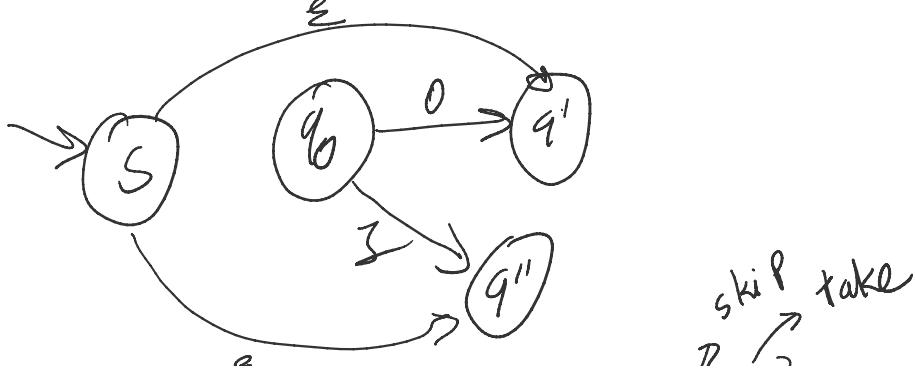
$= a \text{evens}(x)$ if $w = ax$

$\text{evens}(00 \underset{n}{\underset{\pi}{|}} 10 \underset{n}{\underset{\pi}{|}} 0) = 000$

$\text{odds}(00 \underset{n}{\underset{\pi}{|}} 10 \underset{n}{\underset{\pi}{|}} 0) = 011$

Let $M = (Q, \Sigma, q_0, \delta, A)$ be DFA for L .

(construct NFA_M^A for $\text{EVEN}(L)$) $M' = (Q', \Sigma, S, \delta', A')$



$$e \xrightarrow{\quad} \tilde{q}(q') \xrightarrow{\text{skip}} q' \xrightarrow{\text{take}} q$$

$$\mathcal{Q}' = \mathcal{Q} \times \{\tilde{q}, t\}$$

$$s = (q_0, e)$$

$$s'((q, \tilde{e}), e) = \left\{ \left(s(q, 0), \underline{t} \right), \left(s(q, 1), \underline{t} \right) \right\}$$

$$s'((q, \underline{t}), 0) = \left\{ (s(q, 0), e) \right\}$$

take

$$s'((q, t), 1) = \left\{ (s(q, 1), e) \right\}$$

$$A' = \underbrace{\left\{ (q, e) \mid q \in A \right\}}_{\text{odd length}} \cup \underbrace{\left\{ (q, t) \mid q \in A \right\}}_{\text{even length}}.$$

$$\text{even } (0 \underset{e}{\overset{\vee}{|}} 0) = 1$$

$$\text{even } (00 \underset{e}{\overset{\vee}{|}} 11) = 1$$

