Describe deterministic finite-state automata that accept each of the following languages over the alphabet $\Sigma = \{0, 1\}$. Describe briefly what each state in your DFAs \textit{means}.

1. All strings containing the substring \texttt{000}.

\textbf{Solution:}

\begin{center}
\begin{tikzpicture}[node distance=2cm, thick, main/.style = {draw, circle}]
  \node[main] (s) {$s$};
  \node[main] (a) [right of=s] {$a$};
  \node[main] (b) [right of=a] {$b$};
  \node[main] (c) [right of=b] {$c$};
  \draw[->] (s) -- node {$\theta$} (a);
  \draw[->] (a) -- node {$1$} (b);
  \draw[->] (b) -- node {$\theta$} (c);
  \draw[->] (c) -- node {$\theta, 1$} (a);
  \draw[->] (c) -- node {$1$} (s);
\end{tikzpicture}
\end{center}

- $s$: We didn't just read a $0$
- $a$: We've read one $0$ since the last $1$ or the start of the string.
- $b$: We've read two $0$s since the last $1$ or the start of the string.
- $c$: We've read the substring \texttt{000}.

\textbf{2. All strings \textit{not} containing the substring \texttt{000}.}

\textbf{Solution:}

\begin{center}
\begin{tikzpicture}[node distance=2cm, thick, main/.style = {draw, circle}]
  \node[main] (s) {$s$};
  \node[main] (a) [right of=s] {$a$};
  \node[main] (b) [right of=a] {$b$};
  \node[main] (c) [right of=b] {$c$};
  \draw[->] (s) -- node {$\theta$} (a);
  \draw[->] (a) -- node {$1$} (b);
  \draw[->] (b) -- node {$\theta$} (c);
  \draw[->] (c) -- node {$\theta, 1$} (a);
  \draw[->] (c) -- node {$1$} (s);
\end{tikzpicture}
\end{center}

- $s$: We didn't just read a $0$
- $a$: We've read one $0$ since the last $1$ or the start of the string.
- $b$: We've read two $0$s since the last $1$ or the start of the string.
- $c$: We've read the substring \texttt{000}.

(Yes, these are the same states as in problem 1.)
3. All strings in which every run of 0s has length at least 3.

Solution:

4. All strings in which no substring 000 appears before a 1.
   (Equivalently: All strings in which every substring 000 appears after every 1.)

Solution: A string is in this language if and only if it does not contain the substring 0001.
5. All strings containing at least three 0s.

**Solution:**

```
  1  
 s  ->  a  ->  b  ->  c
  0  1  0
      0, 1
```

- **s**: We've read no 0s.
- **a**: We've read one 0.
- **b**: We've read two 0s.
- **c**: We've read at least three 0s; accept.

6. Every string except 000. [Hint: Don't try to be clever.]

**Solution:**

```
  1
 s  ->  a  ->  b  ->  c  ->  d
  0  1  0  0, 1
      0, 1
```

- **s**: We haven’t read anything yet
- **a**: Input so far is 0.
- **b**: Input so far is 00.
- **c**: Input so far is 000.
- **d**: Input is not 000; accept.
Work on these later:

7. All strings $w$ such that in every prefix of $w$, the number of $0$s and $1$s differ by at most 1.

Solution: This is the same as the set of strings that alternate between $0$s and $1$s.

8. All strings containing at least two $0$s and at least one $1$.

Solution:
9. All strings \( w \) such that in every prefix of \( w \), the number of \( 0 \)s and \( 1 \)s differ by at most 2.

Solution:

![Diagram](image)

The fail state \( \infty \) indicates that we have read some prefix where the number of \( 0 \)s and \( 1 \)s differ by more than 2. Each of the other states states \(-2, -1, 0, 1, 2\) indicates the number of \( 1 \)s minus the number of \( 0 \)s of the prefix read so far.

*10. All strings in which the substring \( 000 \) appears an even number of times.
   (For example, \( 0001000 \) and \( 0000 \) are in this language, but \( 00000 \) is not.)

Solution:

![Diagram](image)

Each state is labeled with an integer from 0 to 3, indicating how many consecutive \( 0 \)s have just been read, and a letter \( e \) or \( o \), indicating whether we have read an even or odd number of \( 000 \) substrings.