

17.3.4

On the hereditary nature of shortest paths

You can not shortcut a shortest path

Lemma

G : directed graph with non-negative edge lengths.

$\text{dist}(s, v)$: shortest path length from s to v .

If $s = v_0 \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_k$ shortest path from s to v_k then for any

$0 \leq i < j \leq k$:

$v_i \rightarrow v_{i+1} \rightarrow \dots \rightarrow v_j$ is shortest path from v_i to v_j

Proof.

Suppose not. Then for some $0 \leq i < j \leq k$ there is a path P' from v_i to v_j of length strictly less than that of $s = v_i \rightarrow v_{i+1} \rightarrow \dots \rightarrow v_j$. Then the path

$$s = v_0 \rightarrow v_1 \rightarrow \dots \rightarrow v_i \bullet P' \bullet v_j \rightarrow v_{j+1} \rightarrow \dots \rightarrow v_k$$

is a strictly shorter path from s to v_k than $s = v_0 \rightarrow v_1 \dots \rightarrow v_k$. □

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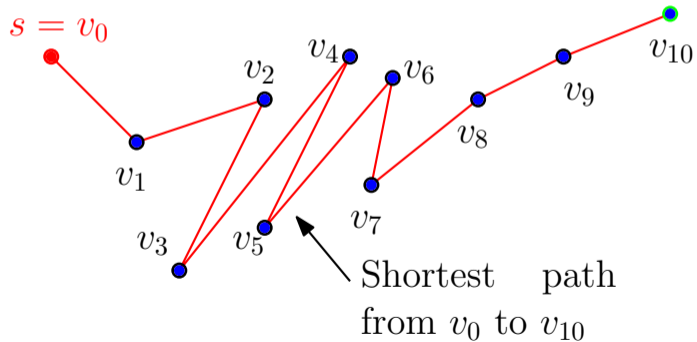
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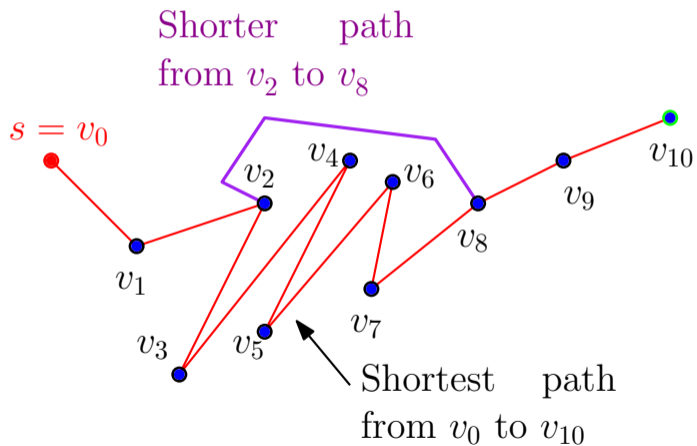
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A proof by picture

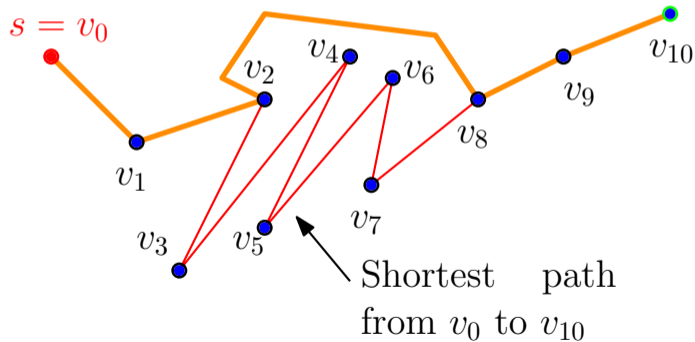


A proof by picture



A proof by picture

A shorter path
from v_0 to v_{10} .
A contradiction.



What we really need...

Corollary

G: directed graph with non-negative edge lengths.

$\text{dist}(\mathbf{s}, \mathbf{v})$: shortest path length from \mathbf{s} to \mathbf{v} .

If $\mathbf{s} = \mathbf{v}_0 \rightarrow \mathbf{v}_1 \rightarrow \mathbf{v}_2 \rightarrow \dots \rightarrow \mathbf{v}_k$ shortest path from \mathbf{s} to \mathbf{v}_k then for any $0 \leq i \leq k$:

- ① $\mathbf{s} = \mathbf{v}_0 \rightarrow \mathbf{v}_1 \rightarrow \mathbf{v}_2 \rightarrow \dots \rightarrow \mathbf{v}_i$ is shortest path from \mathbf{s} to \mathbf{v}_i
- ② $\text{dist}(\mathbf{s}, \mathbf{v}_i) \leq \text{dist}(\mathbf{s}, \mathbf{v}_k)$. *Relies on non-neg edge lengths.*

THE END

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(for now)