

21.3.2

NFAs/DFAs and Universality

DFA Universality

A DFA M is **universal** if it accepts every string.
That is, $L(M) = \Sigma^*$, the set of all strings.

Problem 21.2 (DFA universality).

Input: A DFA M .

Goal: *Is M universal?*

How do we solve **DFA Universality**?

We check if M has any reachable non-final state.

DFA Universality

A DFA M is **universal** if it accepts every string.
That is, $L(M) = \Sigma^*$, the set of all strings.

Problem 21.2 (DFA universality).

Input: A DFA M .

Goal: *Is M universal?*

How do we solve **DFA Universality**?

We check if M has any reachable non-final state.

DFA Universality

A DFA M is **universal** if it accepts every string.
That is, $L(M) = \Sigma^*$, the set of all strings.

Problem 21.2 (DFA universality).

Input: A DFA M .

Goal: *Is M universal?*

How do we solve **DFA Universality**?

We check if M has any reachable non-final state.

DFA Universality

A DFA M is **universal** if it accepts every string.
That is, $L(M) = \Sigma^*$, the set of all strings.

Problem 21.2 (DFA universality).

Input: A DFA M .

Goal: *Is M universal?*

How do we solve **DFA Universality**?

We check if M has any reachable non-final state.

NFA Universality

An **NFA** N is said to be **universal** if it accepts every string. That is, $L(N) = \Sigma^*$, the set of all strings.

Problem 21.3 (NFA universality).

Input: A **NFA** M .

Goal: *Is M universal?*

How do we solve **NFA Universality**?

Reduce it to **DFA Universality**?

Given an **NFA** N , convert it to an equivalent **DFA** M , and use the **DFA Universality** Algorithm.

The reduction takes **exponential time**!

NFA Universality is known to be PSPACE-Complete and we do not expect a polynomial-time algorithm.

NFA Universality

An **NFA** N is said to be **universal** if it accepts every string. That is, $L(N) = \Sigma^*$, the set of all strings.

Problem 21.3 (NFA universality).

Input: A **NFA** M .

Goal: *Is M universal?*

How do we solve **NFA Universality**?

Reduce it to **DFA Universality**?

Given an **NFA** N , convert it to an equivalent **DFA** M , and use the **DFA Universality** Algorithm.

The reduction takes **exponential time**!

NFA Universality is known to be PSPACE-Complete and we do not expect a polynomial-time algorithm.

NFA Universality

An **NFA** N is said to be **universal** if it accepts every string. That is, $L(N) = \Sigma^*$, the set of all strings.

Problem 21.3 (NFA universality).

Input: A **NFA** M .

Goal: *Is M universal?*

How do we solve **NFA Universality**?

Reduce it to **DFA Universality**?

Given an **NFA** N , convert it to an equivalent **DFA** M , and use the **DFA Universality** Algorithm.

The reduction takes **exponential time**!

NFA Universality is known to be PSPACE-Complete and we do not expect a polynomial-time algorithm.

NFA Universality

An **NFA** N is said to be **universal** if it accepts every string. That is, $L(N) = \Sigma^*$, the set of all strings.

Problem 21.3 (NFA universality).

Input: A **NFA** M .

Goal: *Is M universal?*

How do we solve **NFA Universality**?

Reduce it to **DFA Universality**?

Given an **NFA** N , convert it to an equivalent **DFA** M , and use the **DFA Universality** Algorithm.

The reduction takes **exponential time**!

NFA Universality is known to be PSPACE-Complete and we do not expect a polynomial-time algorithm.

THE END

...

(for now)