

11.2

Multiplication using Divide and Conquer

Divide and Conquer

Assume n is a power of 2 for simplicity and numbers are in decimal.

Split each number into two numbers with equal number of digits

- ① $b = b_{n-1}b_{n-2} \dots b_0$ and $c = c_{n-1}c_{n-2} \dots c_0$
- ② $b = b_{n-1} \dots b_{n/2} \mathbf{0} \dots \mathbf{0} + b_{n/2-1} \dots b_0$
- ③ $b(x) = b_L x + b_R$, where $x = 10^{n/2}$, $b_L = b_{n-1} \dots b_{n/2}$ and $b_R = b_{n/2-1} \dots b_0$
- ④ Similarly $c(x) = c_L x + c_R$ where $c_L = c_{n-1} \dots c_{n/2}$ and $c_R = c_{n/2-1} \dots c_0$

Example

$$\begin{aligned} 1234 \times 5678 &= (12x + 34) \times (56x + 78) && \text{for } x = 100. \\ &= 12 \cdot 56 \cdot x^2 + (12 \cdot 78 + 34 \cdot 56)x + 34 \cdot 78. \end{aligned}$$

$$\begin{aligned} 1234 \times 5678 &= (100 \times 12 + 34) \times (100 \times 56 + 78) \\ &= 10000 \times 12 \times 56 \\ &\quad + 100 \times (12 \times 78 + 34 \times 56) \\ &\quad + 34 \times 78 \end{aligned}$$

Divide and Conquer for multiplication

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① $b = b_{n-1}b_{n-2} \dots b_0$ and $c = c_{n-1}c_{n-2} \dots c_0$

② $b \equiv b(x) = b_Lx + b_R$

where $x = 10^{n/2}$, $b_L = b_{n-1} \dots b_{n/2}$ and $b_R = b_{n/2-1} \dots b_0$

③ $c \equiv c(x) = c_Lx + c_R$ where $c_L = c_{n-1} \dots c_{n/2}$ and $c_R = c_{n/2-1} \dots c_0$

Therefore, for $x = 10^{n/2}$, we have

$$\begin{aligned}bc &= b(x)c(x) = (b_Lx + b_R)(c_Lx + c_R) \\ &= b_Lc_Lx^2 + (b_Lc_R + b_Rc_L)x + b_Rc_R \\ &= 10^n b_Lc_L + 10^{n/2}(b_Lc_R + b_Rc_L) + b_Rc_R\end{aligned}$$

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Time Analysis

$$bc = 10^n b_L c_L + 10^{n/2} (b_L c_R + b_R c_L) + b_R c_R$$

4 recursive multiplications of number of size $n/2$ each plus 4 additions and left shifts (adding enough 0's to the right)

$$T(n) = 4T(n/2) + O(n) \quad T(1) = O(1)$$

$T(n) = \Theta(n^2)$. No better than grade school multiplication!

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THE END

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(for now)