

Lecture 3

Rootfinding

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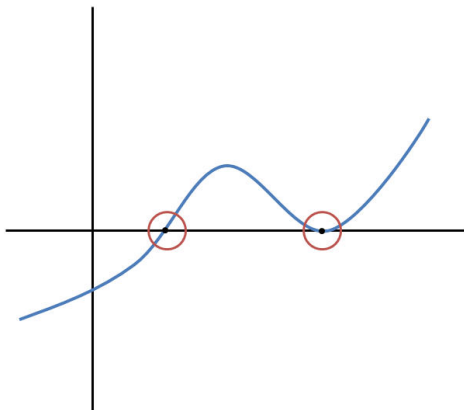
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Root Finding

Given a function $f(x)$, find x so that $f(x) = 0$



Rootfinding

Goals:

- Find roots to equations
 - Compare usability of different methods
 - Compare convergence properties of different methods
- 1 bracketing methods
 - 2 Bisection Method
 - 3 Newton's Method
 - 4 Secant Method
 - 5 (opt) fixed point iterations
 - 6 (opt) special Case: Roots of Polynomials

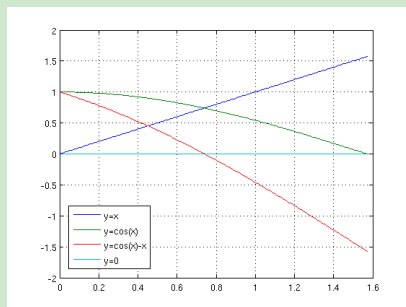


Roots of $f(x)$

- Any single valued equation $g(x) = h(x)$ can be written as $f(x) = g(x) - h(x) = 0$

Example

- Find x so that $\cos(x) = x$
- That is, find where $f(x) = \cos(x) - x = 0$



Analyze your Application

- Is the function complicated to evaluate?
 - lots of expressions?
 - singularities?
 - simplify? polynomial?
- How accurate does our root need to be?
- How fast/robust should our method be?

!

From this, you can pick the right method...



Basic Root Finding Strategy

- 1 Plot the function
 - ▶ Get an initial guess
 - ▶ Identify problematic parts
- 2 Start with the initial guess and iterate

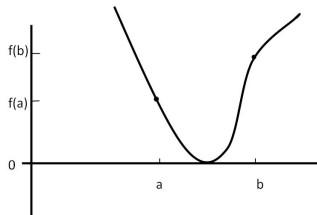
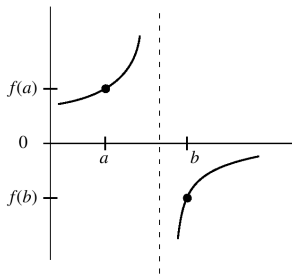
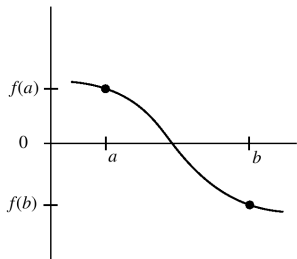
Iteration

We need to study some iterations.

- iteratively finding a root to an equation
- iteratively finding the solution to an algebraic system
- iteratively finding solutions to Ordinary Differential Equations (ODEs)
- ...

Bracket Basics

- A root x is *bracketed* on $[a, b]$ if $f(a)$ and $f(b)$ have opposite sign.
- Changing signs does not guarantee bracketed, however: singularity



Bracket Algorithm

given: $f(x)$, x_{min} , x_{max} , n

Listing 1: Bracket Algorithm

```
1
2
3 dx = (x_max - x_min)/n
4 x_left = x_min
5 i=0
6
7 while i < n:
8     i = i + 1
9     x_right = x_left + dx
10    if (f(x) changes sign in [x_left,x_right]):
11        save [x_left,x_right]# as an interval with a root
12    x_left = x_right
```



Testing Sign

$$f(a) \times f(b) < 0$$

Should we use?

```
fa = myfunc(a);
```

```
fb = myfunc(b);
```

```
if(fa*fb<0)
```

```
    (save)
```

```
end
```



Better Sign Test

!

Nope. Underflow...

sign()

Use Python's sign function

```
import numpy as np
fa = myfunc(a);
fb = myfunc(b);

if np.sign(fa) != np.sign(fb):
    (save)
```



Moving forward...

Bracketing is fine. But we need to find the actual root:

- Bisection
- Newton's Method
- Secant Method
- Fixed Point Iteration

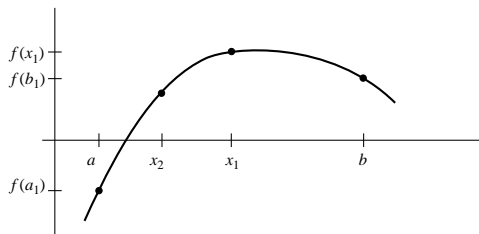
Process:

- 1 Implement the bracket algorithm to get a visual and brackets
- 2 search brackets with these methods



Bisection

Given $f : \mathbb{R} \rightarrow \mathbb{R}$ and $f \in C([a, b])$ and $\text{sign}(f(a)) \neq \text{sign}(f(b))$ by the Intermediate Value Theorem we know we have a bracketed root on the interval $[a, b]$. Bisection Method: halve the interval while continuing to bracket the root.



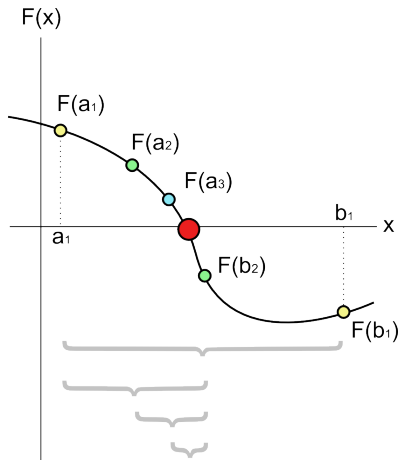
Bisection (2)

For the bracket interval $[a, b]$ the midpoint is

$$x_m = \frac{1}{2}(a + b)$$

idea:

- 1 split bracket in half
- 2 select the bracket that has the root
- 3 goto step 1



Bisection Algorithm

```
1 import numpy as np
2 from scipy import optimize
3 import pprint
4
5 def bisection(f,a1,b1,tol):
6     a = a1
7     b = b1
8     sfb = np.sign(f(b))
9     k=1
10    print('   k   a           b           x_mid       f(x_mid)   width')
11    while b - a > tol:
12        x = (a+b)/2.
13        y = f(x)
14        sfx = np.sign(y)
15        w = np.abs(b-a)
16        print('%5d %10.6f %10.8f %10.8f %11.8f %11.8f' % (k,a,b,x,y,w))
17        if sfx == 0.:
18            a = x
19            b = x
20            break
21        elif sfx == sfb:
22            b = x
23        else:
24            a = x
25        k = k + 1
26
27
28 def f(x):
29     return x - x**(1./3.) - 2
30
31 if __name__ == "__main__":
32     bisection(f,3.,4.,1.e-3)
```



Bisection Example

Solve with bisection:

$$x - x^{1/3} - 2 = 0 \text{ solution from Matlab: } 3.521379706804568$$

k	a	b	x_mid	f(x_mid)	width
1	3.000000	4.000000	3.500000	-0.01829449	1.000000
2	3.500000	4.000000	3.750000	0.19638375	0.500000
3	3.500000	3.750000	3.625000	0.08884159	0.250000
4	3.500000	3.625000	3.562500	0.03522131	0.125000
5	3.500000	3.562500	3.531250	0.00845016	0.062500
6	3.500000	3.531250	3.515625	-0.00492550	0.031250
7	3.515625	3.531250	3.5234375	0.00176150	0.015625
8	3.515625	3.5234375	3.51953125	-0.00158221	0.0078125
9	3.519531	3.5234375	3.52148438	0.00008959	0.00390625
10	3.519531	3.52148438	3.52050781	-0.00074632	0.00195312



Analysis of Bisection

Let $\delta_n = x_{b_n} - x_{a_n}$ be the size of the bracketing interval $[x_{a_n}, x_{b_n}]$ with x_n the middle of the n^{th} stage of bisection. If r is the bracketed root then

$$|x_n - r| \leq \frac{1}{2} \delta_n \text{ where}$$

$$\delta_1 = b - a = \text{initial bracketing interval}$$

$$\delta_2 = \frac{1}{2} \delta_1$$

$$\delta_3 = \frac{1}{2} \delta_2 = \frac{1}{4} \delta_1$$

$$\vdots$$

$$\delta_n = \left(\frac{1}{2}\right)^{n-1} \delta_1 \quad \text{thus}$$

$$|x_n - r| \leq \left(\frac{1}{2}\right)^n \delta_1$$



Analysis of Bisection

$$\frac{\delta_{n+1}}{\delta_1} = \left(\frac{1}{2}\right)^n = 2^{-n} \quad \text{or} \quad n = \log_2\left(\frac{\delta_1}{\delta_{n+1}}\right)$$

n	$\frac{\delta_{n+1}}{\delta_1}$	function evaluations
5	3.1×10^{-2}	7
10	9.8×10^{-4}	12
20	9.5×10^{-7}	22
30	9.3×10^{-10}	32
40	9.1×10^{-13}	42
50	8.9×10^{-16}	52

Remember the game Twenty questions?

Convergence Criteria

An automatic root-finding procedure needs to monitor progress toward the root and stop when current guess is close enough to the desired root.

- Convergence checking will avoid searching to unnecessary accuracy.
- Check how closeness of successive approximations

$$|x_k - x_{k-1}| < \delta_x$$

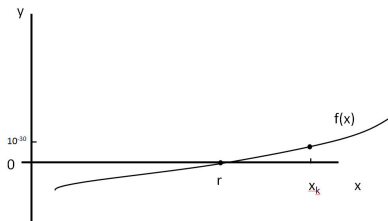
- Check how close $f(x)$ is to zero at the current guess.

$$|f(x_k)| < \delta_f$$

- Which one you use depends on the problem being solved



Convergence Criteria on x versus $f(x)$



Is x_k a sufficient approximation of a root at r ? What if $r = 1$ and $x_k = 100$?

Alternative view

We have two views for finding roots

- Find r such that $f(r) = 0$
- Compute $r = f^{-1}(0)$

The two views give us two ways to determine errors.



Condition Number of Problem

Given a function $G : \mathbb{R} \rightarrow \mathbb{R}$, suppose we wish to compute $y = G(x)$. How sensitive is the solution to changes in x ? We can measure this sensitivity in two ways:

- Absolute Condition Number = $\lim_{h \rightarrow 0} \frac{|G(x+h) - G(x)|}{|h|}$
- Relative Condition Number = $\lim_{h \rightarrow 0} \frac{\frac{|G(x+h) - G(x)|}{|G(x)|}}{\frac{|h|}{|x|}}$

Condition numbers much greater than one mean that the problem is inherently sensitive.



Condition Number Example

Given the problem of finding a root of a function $f : \mathbb{R} \rightarrow \mathbb{R}$, consider the absolute condition number applied to the problem of computing $f^{-1}(0)$.

$$\begin{aligned} \text{Absolute Condition Number} &= \lim_{h \rightarrow 0} \frac{|f^{-1}(0+h) - f^{-1}(0)|}{|h|} \\ &= \left. \frac{df^{-1}(y)}{dy} \right|_{y=0} \quad \text{and from Calculus} \\ &= \frac{1}{\left. \frac{df(x)}{dx} \right|_{x=r}} \end{aligned}$$

We conclude that the root finding problem is inherently sensitive to change if

$$\left| \frac{df(r)}{dx} \right| \approx 0.$$



Condition Number Example

Given the problem of finding a root of a function $f : \mathbb{R} \rightarrow \mathbb{R}$, consider the absolute condition number applied to the problem of computing $f(r)$ where r is a root of f .

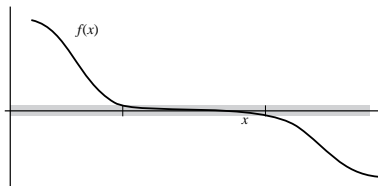
$$\begin{aligned} \text{Absolute Condition Number} &= \lim_{h \rightarrow 0} \frac{|f(r+h) - f(r)|}{|h|} \\ &= \left. \frac{df(x)}{dx} \right|_{x=r} \end{aligned}$$

We conclude that the root finding problem is inherently sensitive to change if $\left| \frac{df(r)}{dx} \right| \gg 1$.

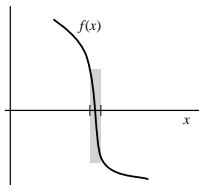


Convergence Criteria Compared

If $f'(x)$ is small near the root, it is easy to satisfy tolerance on $f(x)$ for a large range of Δx .



If $f'(x)$ is large near the root, it is possible to satisfy the tolerance on Δx when $|f(x)|$ is still large.



Convergence rate of a root finding iteration

- Let $e_n = x^* - x_n$ be the error.
- In general, a sequence is said to converge with rate if r is the largest real for which the limit below is finite.

$$\lim_{n \rightarrow \infty} \frac{|e_{n+1}|}{|e_n|^r} = C$$

Special Cases:

- If $r = 1$ and $C = 1$, then the rate is *sublinear*
- If $r = 1$ and $C < 1$, then the rate is *linear*
- If $r > 1$ (i.e. $r = 1$ and $C = 0$), then the rate is *superlinear*
- If $r = 2$ and $C > 0$, then the rate is *quadratic*



Convergence rate of the bisection method

When the bisection method "converges" it can be shown that,

Bisection Method

The bisection method converges with rate $r = 1$ and $C = 0.5$.



Example

Convergence Rate

- 1 $10^{-2}, 10^{-3}, 10^{-4}, 10^{-5} \dots$
- 2 $10^{-2}, 10^{-4}, 10^{-6}, 10^{-8} \dots$
- 3 $10^{-2}, 10^{-3}, 10^{-5}, 10^{-8} \dots$
- 4 $10^{-2}, 10^{-4}, 10^{-8}, 10^{-16} \dots$
- 5 $10^{-2}, 10^{-6}, 10^{-18}, \dots$



Example

Convergence Rate

- 1 $10^{-2}, 10^{-3}, 10^{-4}, 10^{-5} \dots$ (linear with $C = 10^{-1}$)
 - 2 $10^{-2}, 10^{-4}, 10^{-6}, 10^{-8} \dots$ (linear with $C = 10^{-2}$)
 - 3 $10^{-2}, 10^{-3}, 10^{-5}, 10^{-8} \dots$ (superlinear, not quadratic)
 - 4 $10^{-2}, 10^{-4}, 10^{-8}, 10^{-16} \dots$ (quadratic)
 - 5 $10^{-2}, 10^{-6}, 10^{-18}, \dots$ (cubic)
- Linear: Adds equal number of digits of accuracy at each step
 - Quadratic: Doubles the number of digits at each step

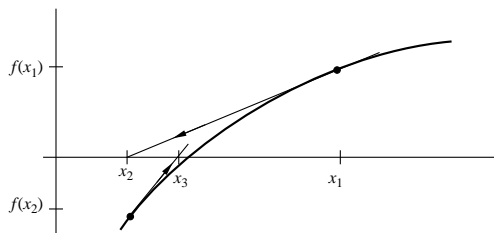


Performing Division

- Ever wondered how a computer process performs division?
- “Long” division requires lookup, subtraction, shifts
- Generates one digit and a time. Can we do better?

To answer this, we need to look at faster methods than bisection

Newton's Method



For a current guess x_k , use $f(x_k)$ and the slope $f'(x_k)$ to predict where $f(x)$ crosses the x axis.

Newton's Method

Expand $f(x)$ in Taylor Series around x_k

$$f(x_k + \Delta x) = f(x_k) + \Delta x \left. \frac{df}{dx} \right|_{x_k} + \frac{(\Delta x)^2}{2} \left. \frac{d^2f}{dx^2} \right|_{x_k} + \dots$$

Substitute $\Delta x = x_{k+1} - x_k$
and neglect 2nd order terms to get

$$f(x_{k+1}) \approx f(x_k) + (x_{k+1} - x_k) f'(x_k)$$

where

$$f'(x_k) = \left. \frac{df}{dx} \right|_{x_k}$$



Newton's Method

Goal is to find x such that $f(x) = 0$.

Set $f(x_{k+1}) = 0$ and solve for x_{k+1}

$$0 = f(x_k) + (x_{k+1} - x_k)f'(x_k)$$

or, solving for x_{k+1}

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$



Newton's Method Algorithm

```
1 initialize:  $x_1 = \dots$   
2 for  $k = 2, 3, \dots$   
3    $x_k = x_{k-1} - f(x_{k-1})/f'(x_{k-1})$   
4   if converged, stop  
5 end
```



Newton's Method Example

Solve:

$$x - x^{1/3} - 2 = 0$$

First derivative is

$$f'(x) = 1 - \frac{1}{3}x^{-2/3}$$

The iteration formula is

$$x_{k+1} = x_k - \frac{x_k - x_k^{1/3} - 2}{1 - \frac{1}{3}x_k^{-2/3}}$$



Newton's Method Example

```
1 import numpy as np
2 from scipy import optimize
3 import pprint
4
5 def newton(f,fp, x,tol):
6     k = 1
7     print('    k           x_k           fp(x_k)     f(x_k)')
8     print('%5d %22.20f %11.8f %11.8g' % (k,x,fp(x),f(x)))
9     k = k + 1
10    while np.abs( f(x) ) > tol:
11        x = x - f(x)/fp(x)
12        print('%5d %22.20f %11.8f %11.8g' % (k,x,fp(x),f(x)))
13        k = k + 1
14
15
16 def f(x):
17     return x - x**(1./3.) - 2.
18
19 def fp(x):
20     return 1. - x**(-2./3.)/3.
21
22
23
24 if __name__ == "__main__":
25     , newton(f, fp, 3., 1.e-25)
```



Newton's Method Example

$$x_{k+1} = x_k - \frac{x_k - x_k^{1/3} - 2}{1 - \frac{1}{3}x_k^{-2/3}}$$

The approximate true root = 3.52137970680457046412926

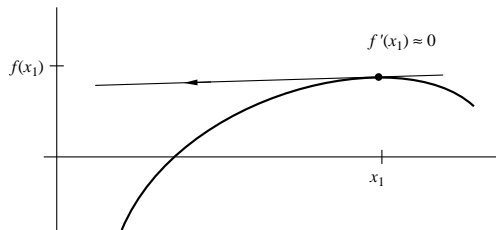
k	x_k	fp(x_k)	f'(x_k)
1	3.00000000000000000000	0.83975005	-0.44224957
2	3.52664429313903271535	0.85612976	0.0045067918
3	3.52138014739732829739	0.85598641	3.7714141e-07
4	3.52137970680457090822	0.85598640	2.6645353e-15
5	3.52137970680456779959	0.85598640	0

Conclusion

- Newton's method converges *much* more quickly than bisection
- Newton's method requires an analytical formula for $f'(x)$
- The algorithm is simple as long as $f'(x)$ is available.
- Iterations are not guaranteed to stay inside an ordinal bracket.



Divergence of Newton's Method



Since

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$

the new guess, x_{k+1} , will be far from the old guess whenever $f'(x_k) \approx 0$



Newton's Method: Convergence

Recall

Convergence of a method is said to be of order r if there is a constant C such that

$$\lim_{k \rightarrow \infty} \frac{|e_{k+1}|}{|e_k|^r} = C$$

If Newton's method converges then it is of order 2 (quadratic) when $f'(x_*) \neq 0$. (assuming f'' is continuous) For ξ_k between x_k and x_*

$$f(x_*) = f(x_k) + (x_* - x_k)f'(x_k) + \frac{1}{2}(x_* - x_k)^2 f''(\xi_k) = 0$$

So

$$\frac{f(x_k)}{f'(x_k)} + x_* - x_k + \frac{1}{2}(x_* - x_k)^2 \frac{f''(\xi_k)}{f'(x_k)} = 0$$

Then

$$x_* - x_{k+1} + \frac{1}{2}(x_* - x_k)^2 \frac{f''(\xi_k)}{f'(x_k)} = 0$$

Thus

$$\frac{|x_* - x_{k+1}|}{|x_* - x_k|^2} = \frac{1}{2} \left| \frac{f''(\xi_k)}{f'(x_k)} \right| \rightarrow \frac{1}{2} \left| \frac{f''(x_*)}{f'(x_*)} \right| \text{ as } x_k \rightarrow x_*$$



Reciprocal Approximation

- Consider the task of computing $1/q$ for some q without using division.
- We can write this as: find the root x of $f(x) = 1/(xq) - 1 = 0$.
- What is Newton's Method for this?
- $f'(x) = -1/(x^2q)$. Thus

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

or

$$x_{n+1} = x_n - \left(\frac{1/(x_n q) - 1}{-1/(x_n^2 q)} \right)$$



Reciprocal Approximation

- Consider the task of computing $1/q$ for some q without using division.
- We can write this as: find the root x of $f(x) = 1/(xq) - 1 = 0$.
- What is Newton's Method for this?
- $f'(x) = -1/(x^2q)$. Thus

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

or

$$x_{n+1} = x_n - \left(\frac{1/(x_n q) - 1}{-1/(x_n^2 q)} \right) \frac{x_n^2 q}{x_n^2 q}$$

$$x_{n+1} = x_n + x_n - x_n^2 q = 2x_n - x_n^2 q = 2x_n - x_n^2 q$$



Example: Compute $1/3 = 0.01010101\dots$ binary

- Find the bracket:

- $1/2 > 1/3 > 1/4$

① $x_1 = 1/4$

② $x_2 = 2x_1 - x_1^2q = 1/2 - 3/16 = 5/16 = 0.0101$ (*binary*)

③ $x_3 = 2 \times 5/2^4 - 3 \times 25/2^8 = (160 - 75)/2^8 = 85/2^8 = 0.01010101$ (*binary*)

④ $x_4 = 2 \times 85/2^8 - 3 \times 85^2/2^{16} = 21845/2^{16} = 0.0101010101010101$ (*binary*)

In 3 steps, computed 16 bits in $1/3$

How many binary digits are computed in the next step?



Instructor Notes

- Modification of Newton's Method for root finding when $\frac{df}{dx}(\text{root}) = 0$. Use the formula,

$$x_{n+1} = x_n - m * \frac{f(x_n)}{f'(x_n)}$$

where m is the multiplicity of the root.

- or solve $0 = g(x) = \frac{f(x)}{f'(x)}$

