



# CS 225

## Data Structures

*February 5 – Trees*

*G Carl Evans*

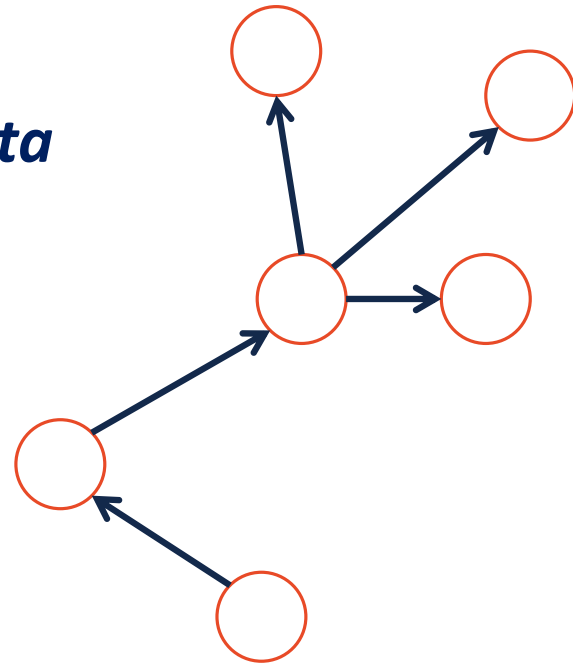
# Trees

*“The most important non-linear data structure in computer science.”*

*- David Knuth, The Art of Programming, Vol. 1*

**A tree is:**

- 
- 



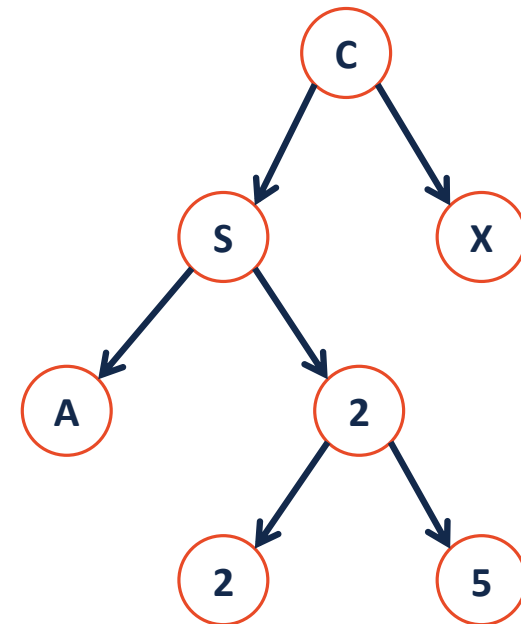
# Binary Tree – Defined

*A binary tree T is either:*

- 

**OR**

- 



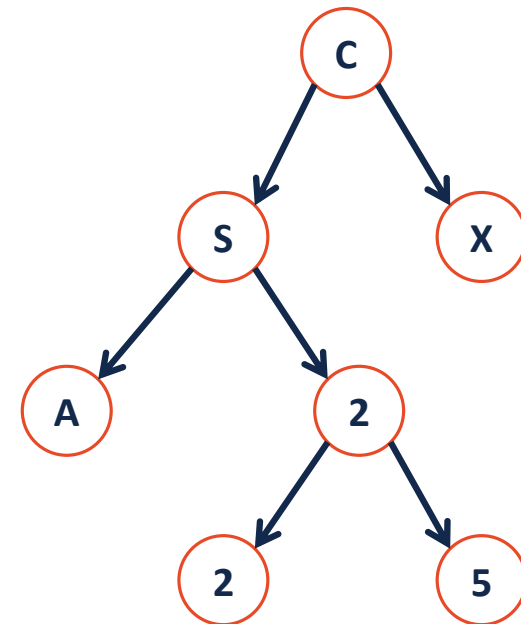
# Binary Tree – Defined

*A binary tree T is either:*

- $T = \emptyset$

OR

- $T = (r, T_L, T_R)$

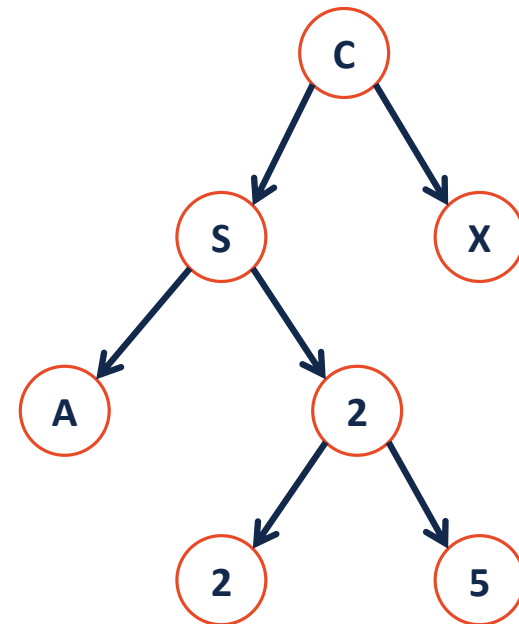


# Tree Property: height

***height(T)***: length of the longest path from the root to a leaf

**Given a binary tree T:**

***height(T) =***



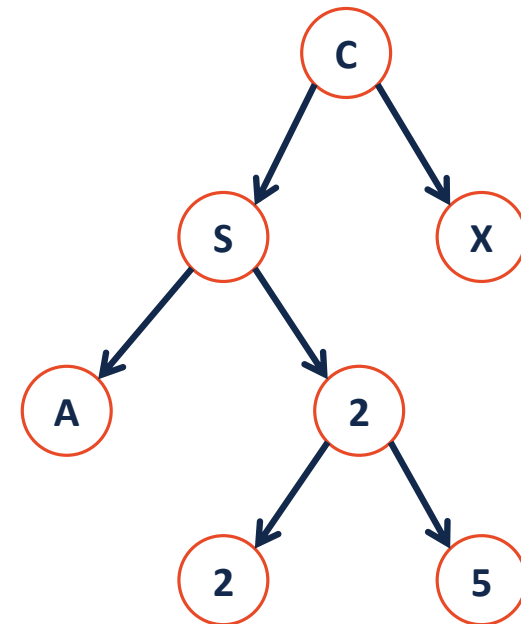
# Tree Property: height

***height(T)***: length of the longest path from the root to a leaf

**Given a binary tree T:**

$$\mathit{height}(T) = \max(\mathit{height}(T_L), \mathit{height}(T_R)) + 1$$

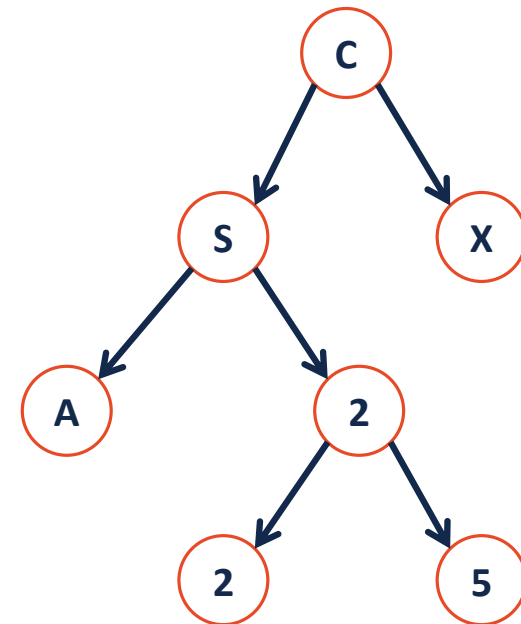
$$\mathit{height}(\emptyset) = -1$$



# Tree Property: full

A tree  $F$  is **full** if and only if:

- 1.
- 2.

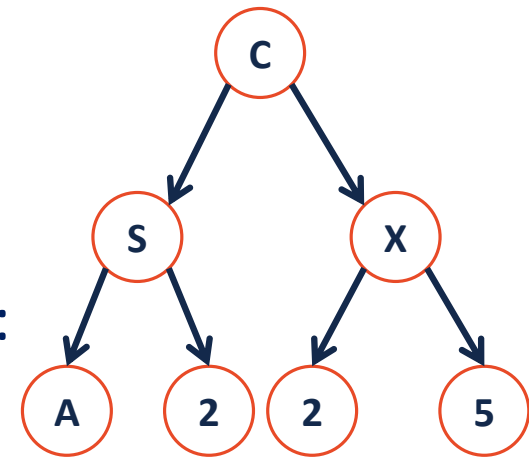


# Tree Property: perfect

A **perfect** tree  $P$  is defined in terms of the tree's height.

Let  $P_h$  be a perfect tree of height  $h$ , and:

- 1.
- 2.

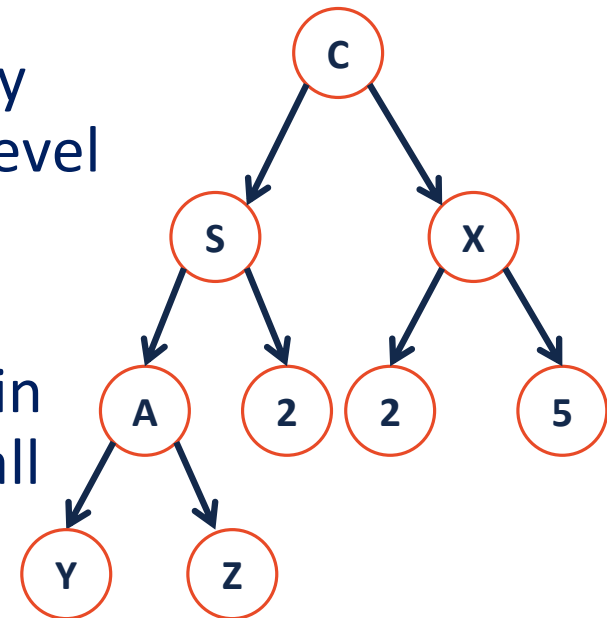




# Tree Property: complete

**Conceptually:** A perfect tree for every level except the last, where the last level is “pushed to the left”.

**Slightly more formal:** For all levels  $k$  in  $[0, h-1]$ ,  $k$  has  $2^k$  nodes. For level  $h$ , all nodes are “pushed to the left”.



# Tree Property: complete

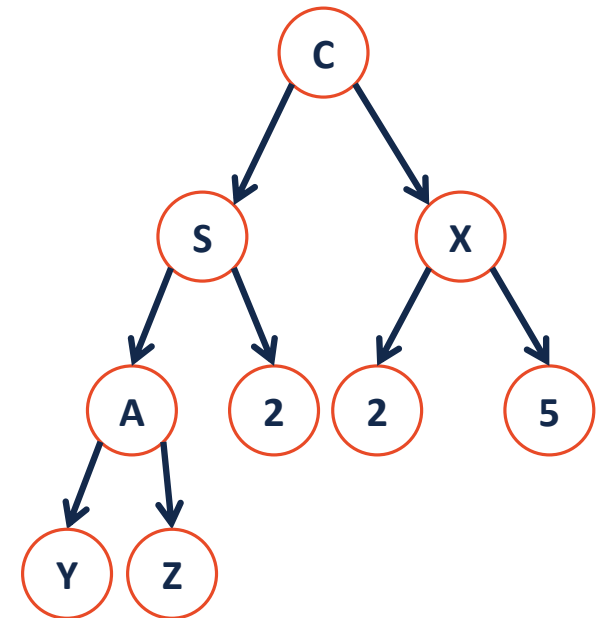
A **complete** tree  $C$  of height  $h$ ,  $C_h$ :

1.  $C_{-1} = \{\}$
2.  $C_h$  (where  $h > 0$ ) =  $\{r, T_L, T_R\}$  and either:

$T_L$  is \_\_\_\_\_ and  $T_R$  is \_\_\_\_\_

**OR**

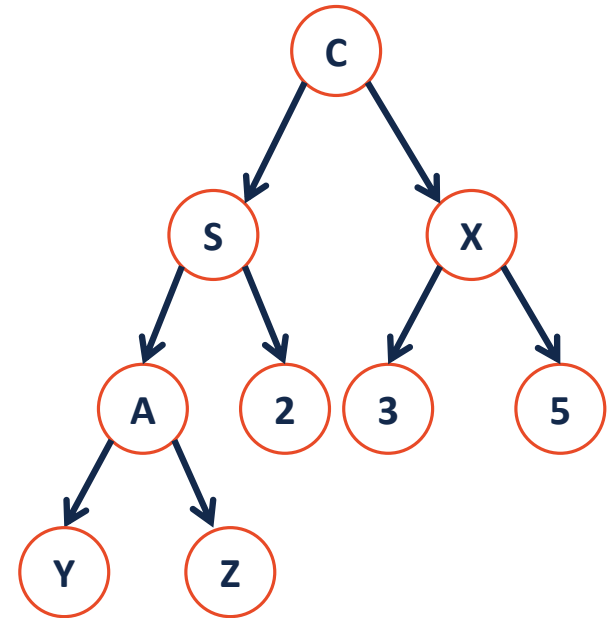
$T_L$  is \_\_\_\_\_ and  $T_R$  is \_\_\_\_\_



# Tree Property: complete

Is every **full** tree **complete**?

If every **complete** tree **full**?





## Tree ADT

**insert**, inserts an element to the tree.

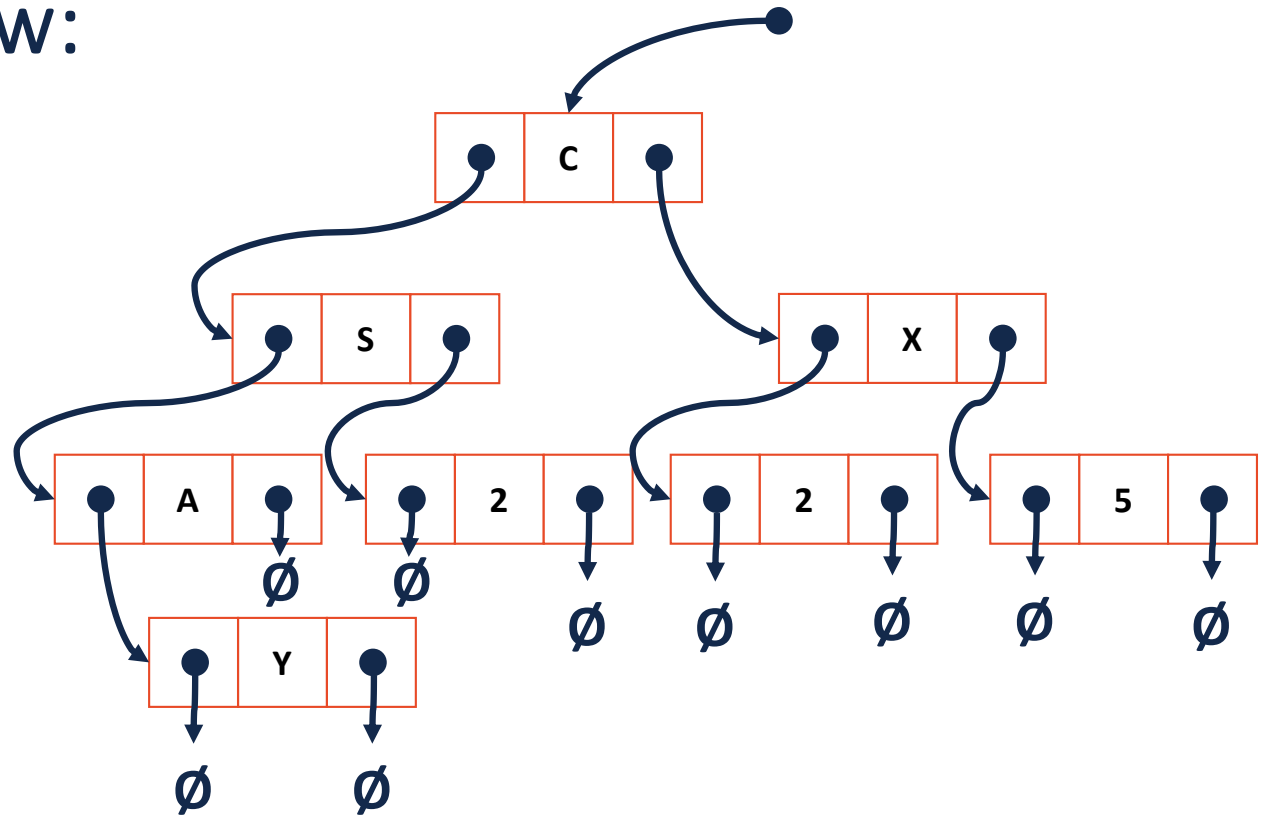
**remove**, removes an element from the tree.

**access**, access elements from the tree.

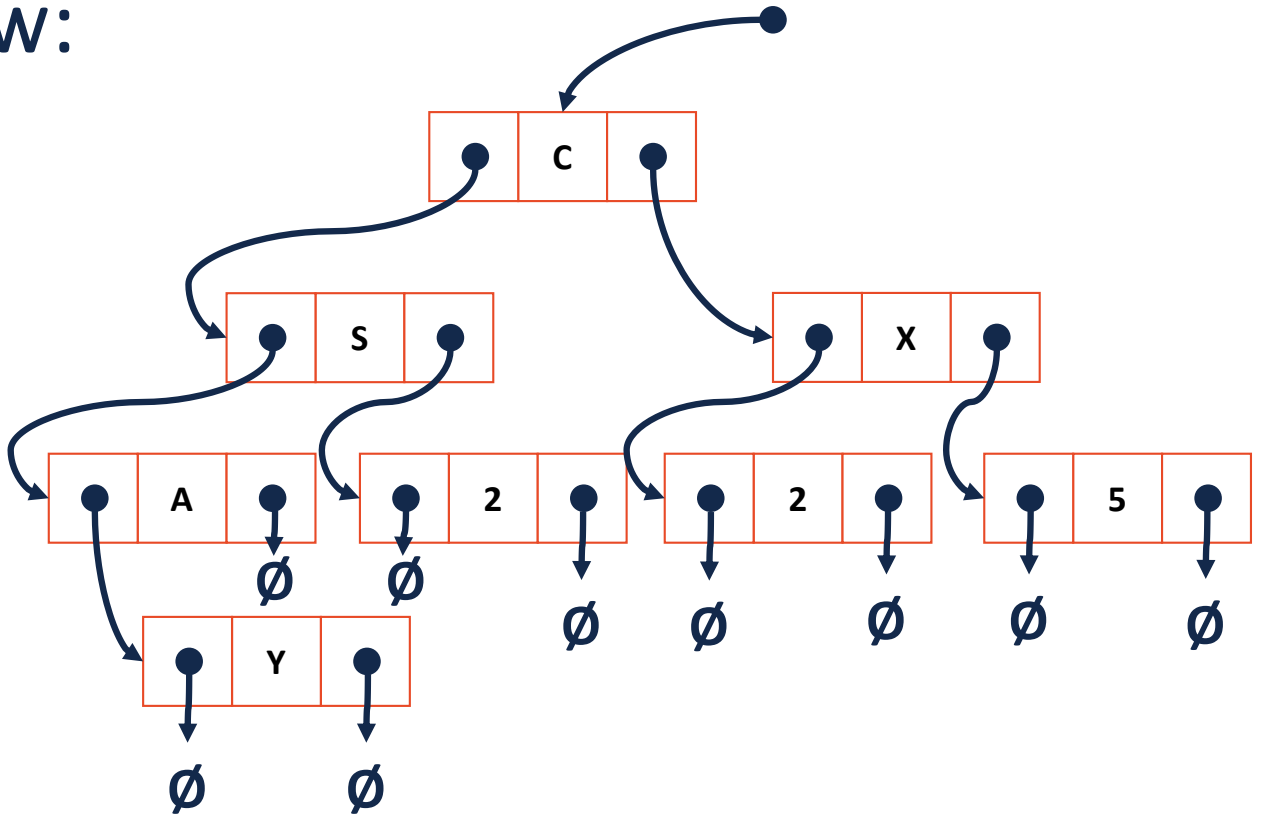
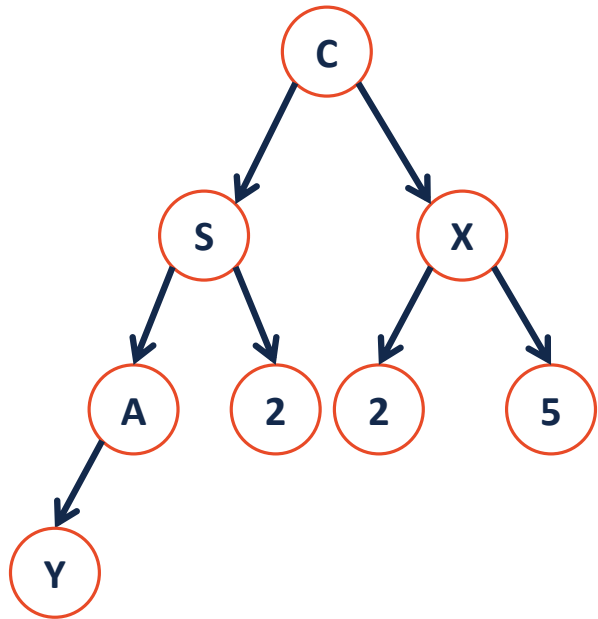
## BinaryTree.h

```
1 #pragma once
2
3 template <class T>
4 class BinaryTree {
5     public:
6         /* ... */
7
8     private:
9
10
11
12
13
14
15
16
17
18
19 };
```

Trees aren't new:



Trees aren't new:





## How many NULLs?

**Theorem:** If there are  $n$  data items in our representation of a binary tree, then there are \_\_\_\_\_ NULL pointers.





# How many NULLs?

**Base Cases:**

**$n = 0$ :**

**$n = 1$ :**

**$n = 2$ :**



# How many NULLs?

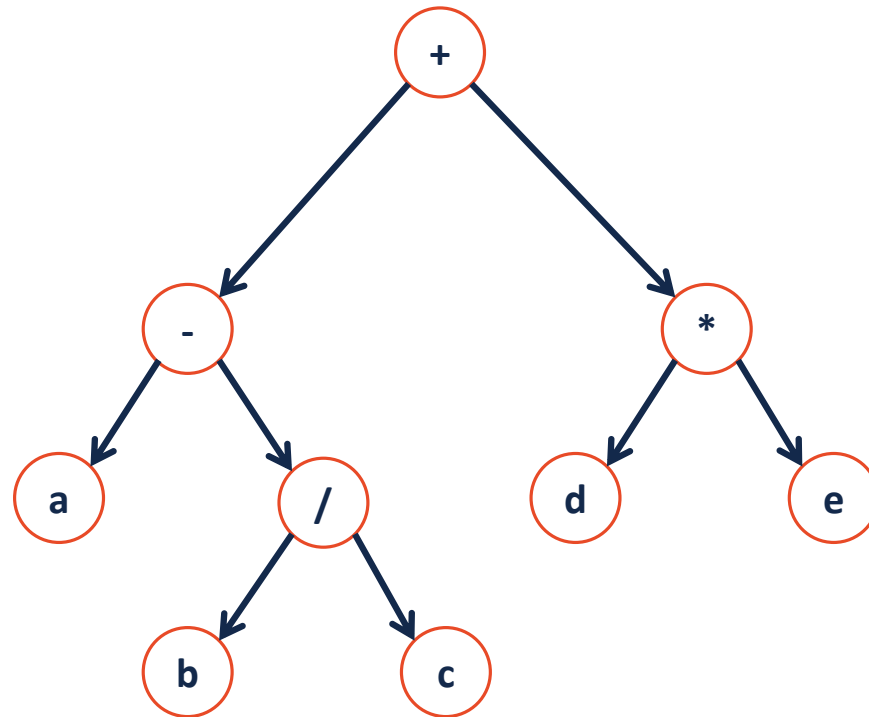
**Induction Hypothesis:**



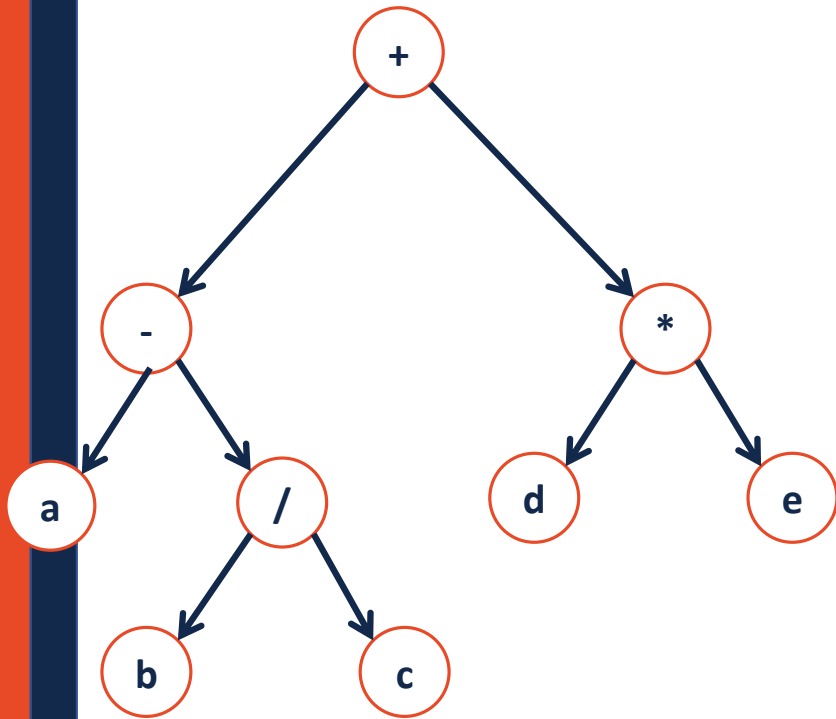
## How many NULLs?

Consider an arbitrary tree **T** containing **n** data elements:

# Access All the Nodes - Traversals

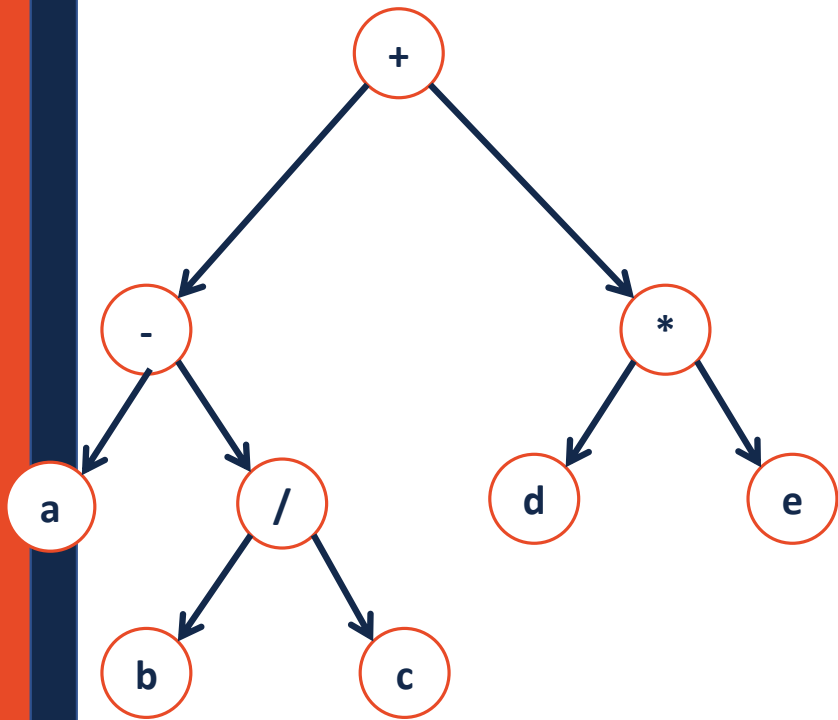


# Traversals



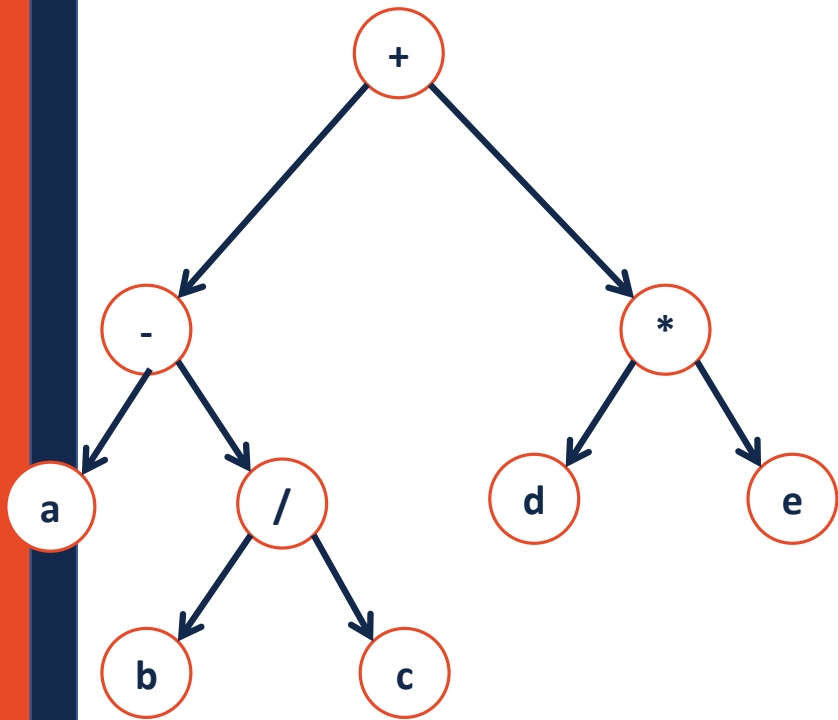
```
49 template<class T>
50 void BinaryTree<T>::__Order(TreeNode * cur)
51 {
52
53
54
55
56
57
58 }
```

# Traversals



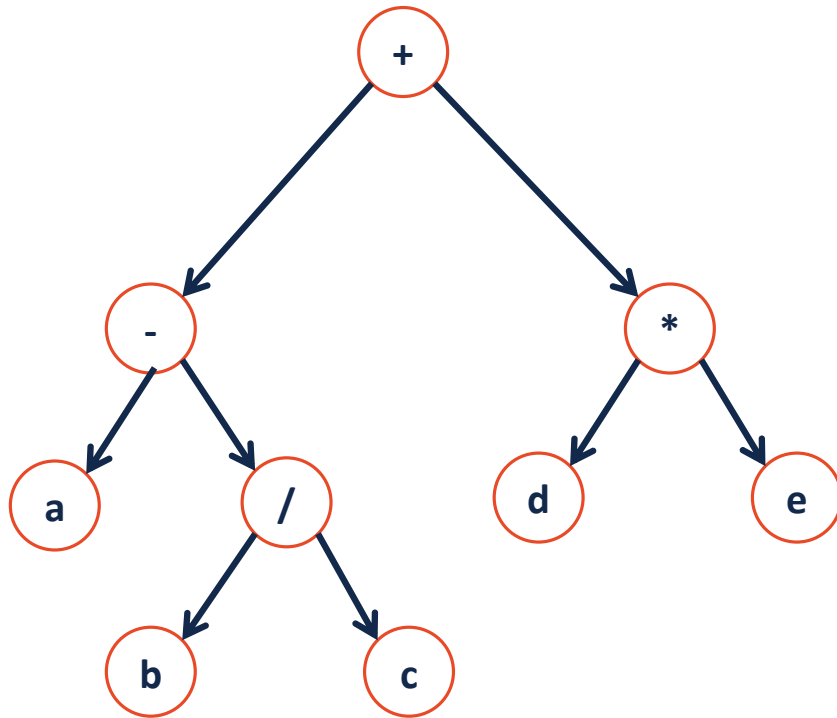
```
49 template<class T>
50 void BinaryTree<T>::__Order(TreeNode * cur) {
51     if (cur != NULL) {
52         _____;
53         __Order(cur->left);
54         _____;
55         __Order(cur->right);
56         _____;
57     }
58 }
```

# Traversals



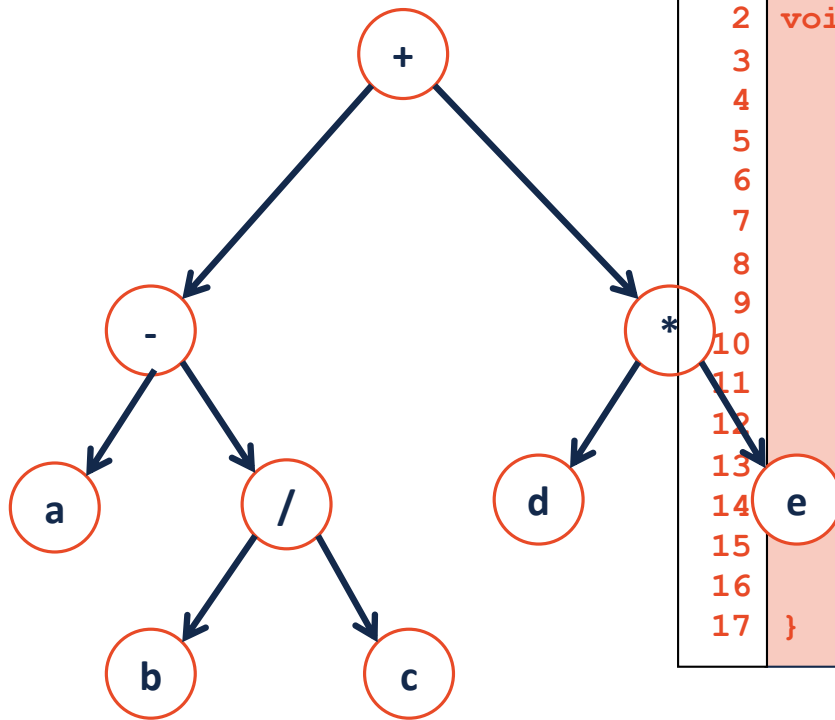
```
49 template<class T>
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54         _____;
55         __Order(cur->right);
56         _____;
57     }
58 }
```

# A Different Type of Traversal





# A Different Type of Traversal



```
1  template<class T>
2  void BinaryTree<T>::levelOrder(TreeNode * root) {
3
4
5
6
7
8
9
10
11
12
13
14  e
15
16
17 }
```



# Traversal vs. Search

**Traversal**

**Search**



# Search: Breadth First vs. Depth First

**Strategy: Breadth First Search (BFS)**

**Strategy: Depth First Search (DFS)**