

#22: kd-Trees and BTrees Intro

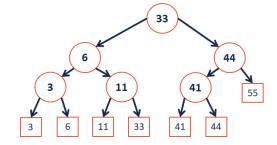
March 6, 2019 · Fagen-Ulmschneider, Zilles

Range-based Searches:

Q: Consider points in 1D: $p = \{p_1, p_2, ..., p_n\}$what points fall in [11, 42]?



Tree Construction:

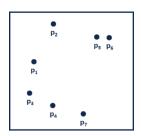


Range-based Searches:

Running Time:

Extending to k-dimensions:

Consider points in 2D: $p = \{p_1, p_2, ..., p_n\}$:

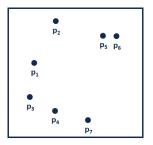


...what points are inside a range (rectangle)?

...what is the nearest point to a query point \mathbf{q} ?

kd-Tree Motivation:

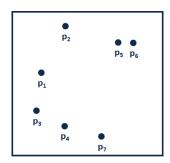
First, let's try and divide our space up:

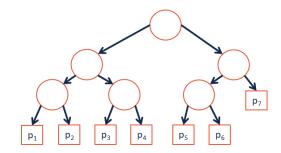


kd-Tree Construction:

How many dimensions exist in our input space?

How do we want to "order" our dimensions?





Motivation

Can we always fit our data in main memory?

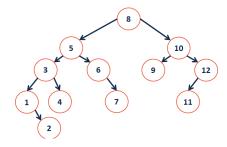
Where else do we keep our data?

-

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vs. CPU: 3 GHz == 3 billion ops / sec * ____ cores

AVL Operations on Disk:



How deep do AVL trees get?

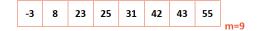
BTree Motivations

Knowing that we have long seek times for data, we want to build a data structure with two (related) properties:

1.

2.

BTree_m



Goal: Build a tree that uses ______/node! _____/node! _____/node!

A **BTree of order m** is an m-way tree where:

- 1. All keys within a node are ordered.
- 2. All leaves contain no more than m-1 nodes.

BTree Insert, using m=5

...when a BTree node reaches **m** keys:

CS 225 - Things To Be Doing:

- 1. Programming Exam B starts next Tuesday (March 12th)
- 2. MP4 extra credit ongoing (final deadline Monday, March. 11th)
- 3. lab_avl released this week; course feedback in lab this week!
- **4.** Daily POTDs are ongoing!