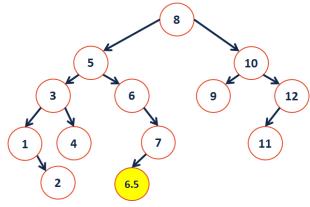
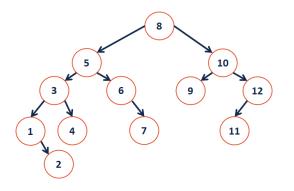


**CS 2**<br/>**2 5**#20: AVL Analysis<br/>March 1, 2019 · Wade Fagen-Ulmschneider

#### **AVL Insertion**



## **AVL Removal**



# **Running Times:**

	AVL Tree
find	
insert	
remove	

# **Motivation:**

Big-O is defined as:

Let **f(n)** describe the height of an AVL tree in terms of the number of nodes in the tree (**n**). Visually, we can represent the big-O relation:

 $f(n) \le c \times g(n)$ : Provides an upper bound:

The height of the tree, **f(n)**, will always be less than **c** × **g(n)** for all values where  $\mathbf{n} > \mathbf{k}$ .

 $f^{-1}(h) \ge c \times g^{-1}(h)$ : Provides a lower bound:

The number of nodes in the tree, **f**<sup>1</sup>(**h**), will always be <u>greater</u> <u>than</u>  $\mathbf{c} \times \mathbf{g}^{-1}(\mathbf{h})$  for all values where  $\mathbf{n} > \mathbf{k}$ .

# **Plan of Action:**

Goal: Find a function that defines the lower bound on **n** given **h**.

Given the goal, we begin by defining a function that describes the smallest number of nodes in an AVL of height **h**:

Proving our IH:

**V.** Using a proof by induction, we have shown that:

...and by inverting our finding:

Theorem:

An AVL tree of height **h** has at least \_\_\_\_\_\_.

I. Consider an AVL tree and let **h** denote its height.

**II.** Case: \_\_\_\_\_

**III.** Case: \_\_\_\_\_

### **Summary of Balanced BSTs:**

Advantages	Disadvantages

**IV.** Case: \_\_\_\_\_

Inductive hypothesis (IH):

# CS 225 – Things To Be Doing:

- **1.** Theory Exam 2 is ongoing
- **2.** mp\_traversals is out, EC+7 this coming Monday
- **3.** lab\_huffman is due on Sunday
- **4.** Daily POTDs are ongoing, +1 each completed POTD! :)