

# Data Structures

## Graph Traversals

CS 225

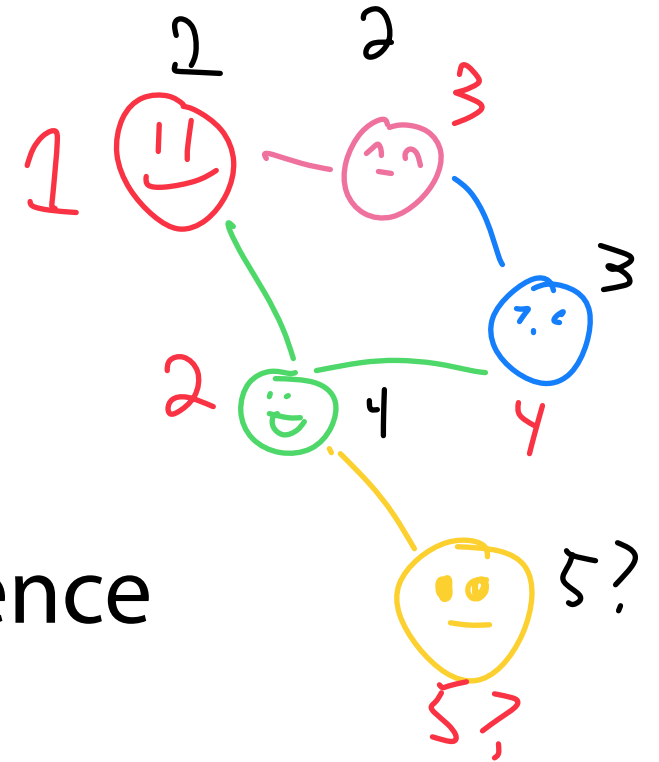
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# Learning Objectives

Discuss graph traversal algorithms



$|V| = n, |E| = m$



H.T.  $O(1)^*$

Expressed as O(f)	Edge List	Adjacency Matrix	Adjacency List
Space	$n+m$	$n^2$	$n+m$
insertVertex(v)	<del>1*</del>	$n^*$	<del>1*</del>
removeVertex(v)	<u><math>m^{**}</math></u>	$n$	$\text{deg}(v)$
insertEdge(u, v)	$1$	$1$	$1^*$
removeEdge(u, v)	$m$	$1$	$\text{min}(\text{deg}(u), \text{deg}(v))$
incidentEdges(v)	$m$	$n$	$\text{deg}(v)$
areAdjacent(u, v)	$m$	$1$	$\text{min}(\text{deg}(u), \text{deg}(v))$

amortized  $O(1)$  H.T. expected  $O(1)$

$M+n$

# Graph Traversals

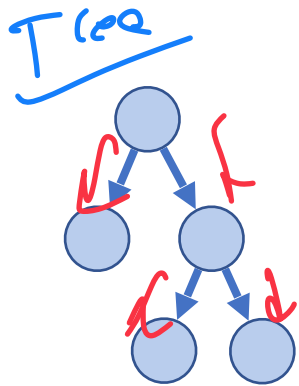
→ why + reversal?  
↳ solving a maze (specific)  
↳ shortest path

Sub-structures  
information about  
graph structure

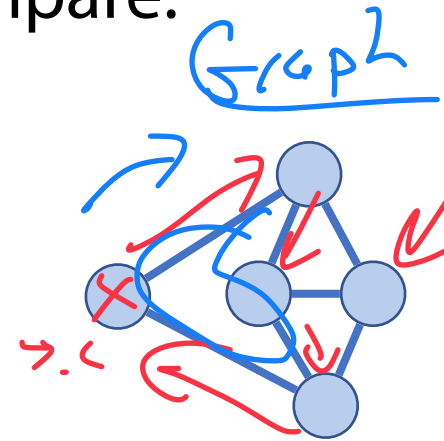
**Objective:** Visit every vertex and every edge in the graph.

How can we systematically go through a complex graph in the fewest steps?

Tree traversals won't work — lets compare:

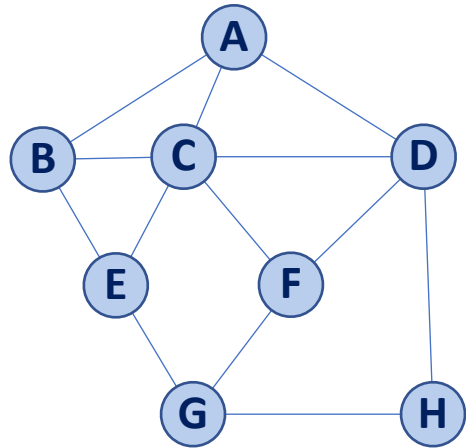


- Rooted
- Acyclic
- Notion of completeness



- No root / clear start node
- Cycles
- No notion of completeness

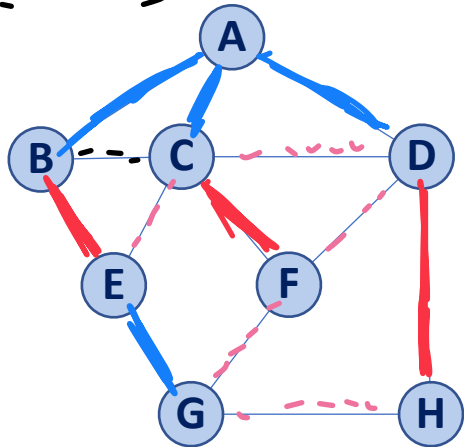
# Traversal: BFS



- 1) Starting node (Random or user input)
- 2) Track current node
- 3) Track nodes and edges we have visited

# Traversal: BFS

Discovery  
--- cross



V (Current node)  
W (my neighbors)

0) First edge = undiscovered distance labels = distance  
1) Initialize a queue (u/v)  
↳ add v to queue  
↳ storing dist / pred

2) while queue not empty  
↳ Dequeue (v)  
↳ Add all unvisited neighbors to queue  
↳ Process neighbors  
↳ Set Dist  
↳ Set Pred

↳ Predecessor

v	d	P	Adjacent Edges
A	0	null	B C D
B	1	A	<del>A</del> <del>C</del> <del>E</del>
C	1	A	<del>A</del> <del>B</del> <del>D</del> <del>E</del> <del>F</del>
D	1	A	<del>A</del> <del>C</del> <del>F</del> <del>H</del>
E	2	B	<del>B</del> <del>C</del> <del>G</del>
F	2	C	<del>C</del> <del>D</del> <del>G</del>
G	2	E	<del>E</del> <del>F</del> <del>H</del>
H	2	D	<del>D</del> <del>G</del>

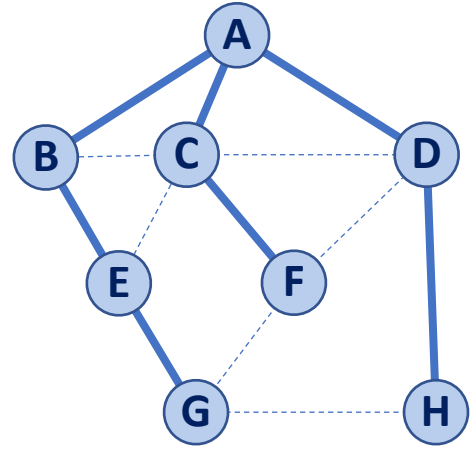


front

back



# Traversal: BFS $n = |V|, m = |E|$



v	d	P	Adjacent Edges
A	0	-	B C D
B	1	A	A C E
C	1	A	A B D E F
D	1	A	A C F H
E	2	B	B C G
F	2	C	C D G
G	3	E	E F H
H	2	D	D G

$$O(m) \rightsquigarrow 2|E| = \sum_{v \in V} \text{deg}(v)$$

Runtime?

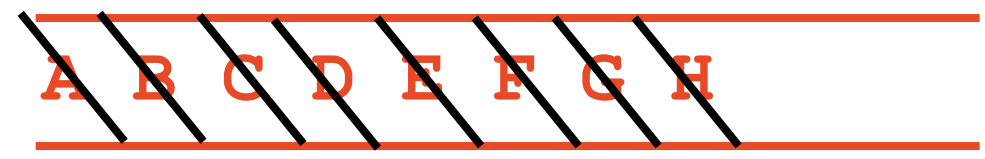
↳ # of times looked at vertex

↳ # of times processed vertex

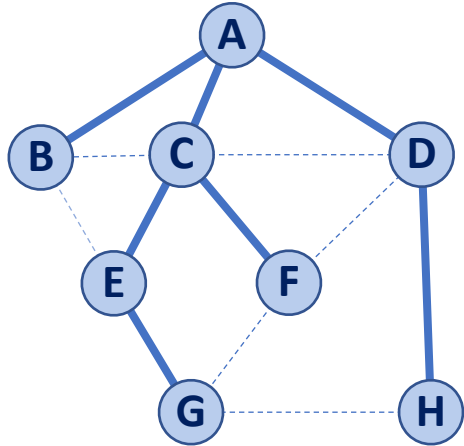
↳ # of times labeled edge

looked  $2m$

$$O(n) \longrightarrow O(m+n)$$



# Traversal: BFS



v	d	P	Adjacent Edges
A	0	-	<b>C B D</b>
B	1	A	A C E
C	1	A	A B D E F
D	1	A	A C F H
E	2	B	B C G
F	2	C	C D G
G	3	E	E F H
H	2	D	D G

Different traversals for  
 different start nodes  
 &  
 different neighbor order





**Input:** Graph, G

**Output:** A labeling of the edges in G as discovery or cross

(tree)

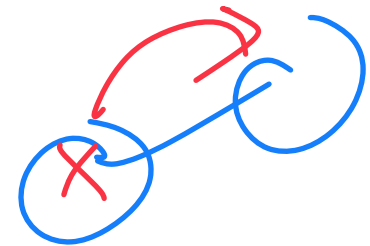
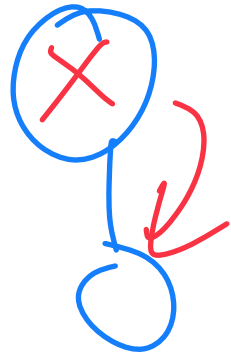
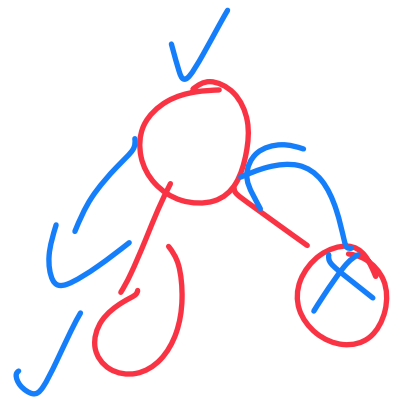
```
1 BFS(G):
2   foreach (Vertex v : G.vertices()):
3     setPred(v, NULL)
4     setDist(v, -1)
5
6   foreach (Edge e : G.edges()):
7     setLabel(e, UNEXPLORED)
8
9   foreach (Vertex v : G.vertices()):
10    if getDist(v) == -1:
11      BFS(G, v)
```

Initialize step

Set labels

Unconnected Component

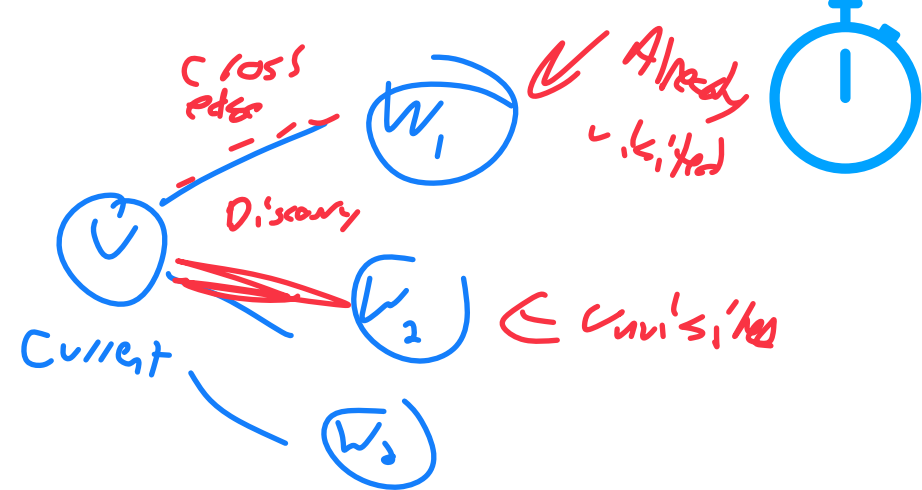
"random" start



```

1 BFS(G):
2   foreach (Vertex v : G.vertices()):
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5
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8
9   foreach (Vertex v : G.vertices()):
10    if getDist(v) == -1:
11      BFS(G, v)

```



```

12 BFS(G, v):
13   Queue q
14   setDist(v, 0)
15   q.enqueue(v)
16
17   while !q.empty():
18     v = q.dequeue()
19
20     foreach (Vertex w : G.adjacent(v)):
21       if( getDist(w) == -1):
22         setLabel((v, w), DISCOVERY)
23         setPred(w, v)
24         setDist(w, v + 1)
25         q.enqueue(w)
26       else * if edge unvisited?
27         setLabel((v, w), CROSS)

```

initialize

Look at edge

unvisited vertex / edge sets label

if edge unvisited?

# Count connected components?

Yes

# of cs in dist : s  
# components

```

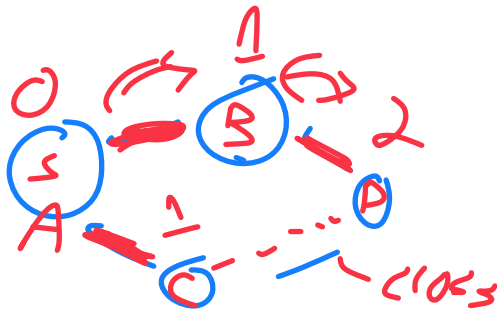
1 BFS (G) :
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7     setLabel(e, UNEXPLORED)
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9   foreach (Vertex v : G.vertices()) :
10    if getDist(v) == -1:
11      BFS (G, v)

```

count #

## Cycle Detection?

↳ cross edge is  
at least one cycle



```

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20   foreach (Vertex w : G.adjacent(v)) :
21     if( getDist(w) == -1) :
22       setLabel((v, w), DISCOVERY)
23       setPred(w, v)
24       setDist(w, v + 1)
25       q.enqueue(w)
26   else IF edge unlabeled (edge == UNEXPLORED)
27     setLabel((v, w), CROSS)

```

IF edge unlabeled (edge == UNEXPLORED)

# BFS Observations

What is the shortest path from **A** to **H**?

2

A - D - H

What is the shortest path from **E** to **H**?

↳ No idea!

BFS gives shortest path from source to every other node

v	d	P	Adjacent Edges
A	0	-	C B D
B	1	A	A C E
C	1	A	B A D E F
D	1	A	A C F H
E	2	C	B C G
F	2	C	C D G
G	3	E	E F H
H	2	D	D G

source →

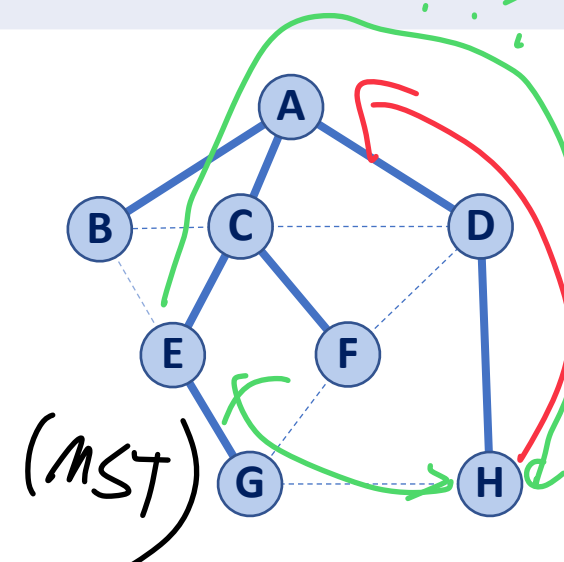
???

If my node has distance **d**, do I know anything about the nodes connected by a **cross edge**?

vertex  $u, v$  :  $|d(u) - d(v)| \leq 1$   
w/ cross edge

What structure is made from **discovery edges**?

Minimum Spanning tree



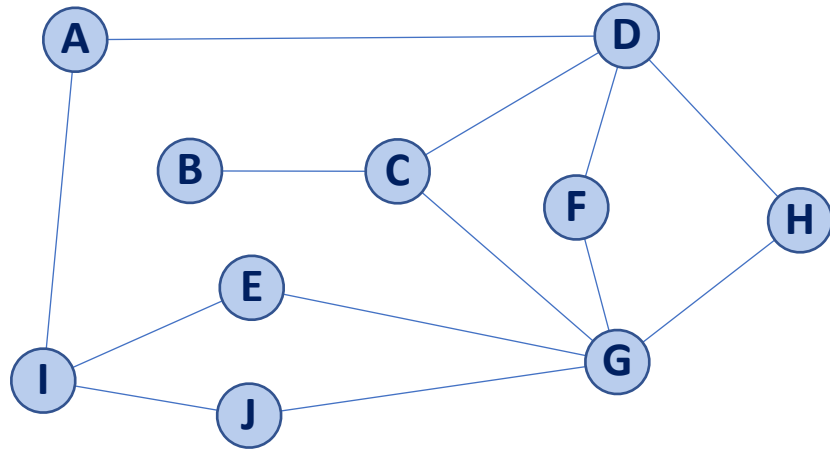
# BFS Observations

1. BFS can be used to count components
2. BFS can be used to detect cycles
3. The BFS 'distance' value is always the shortest distance from source to any vertex (and the discovery edges form a MST)
4. The endpoints of a cross edge never differ in distance by more than 1 (  $|\mathbf{d(u)} - \mathbf{d(v)}| \leq 1$  )



# Traversal: DFS

↪ DFS is BFS but a stack  
and not a queue

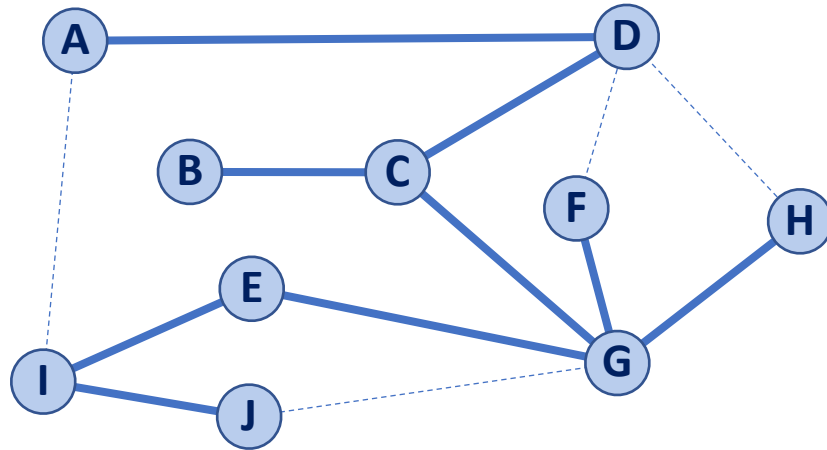




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12 DFS (G, v) :
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15     if( getDist(w) == -1):
16       setLabel((v, w), DISCOVERY)
17       setPred(w, v)
18       setDist(w, v + 1)
19       DFS (G, w)
20   else:
21     setLabel((v, w), BACK)
22
23
24
25
26
27
```

# Traversal: DFS



Does distance have meaning here?

Do our edge labels have meaning here?

————— Discovery Edge

----- Back Edge



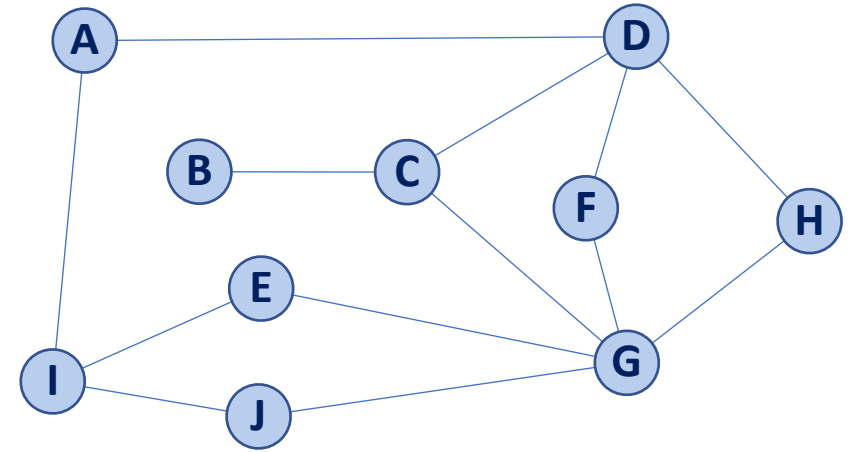
# Running time of DFS

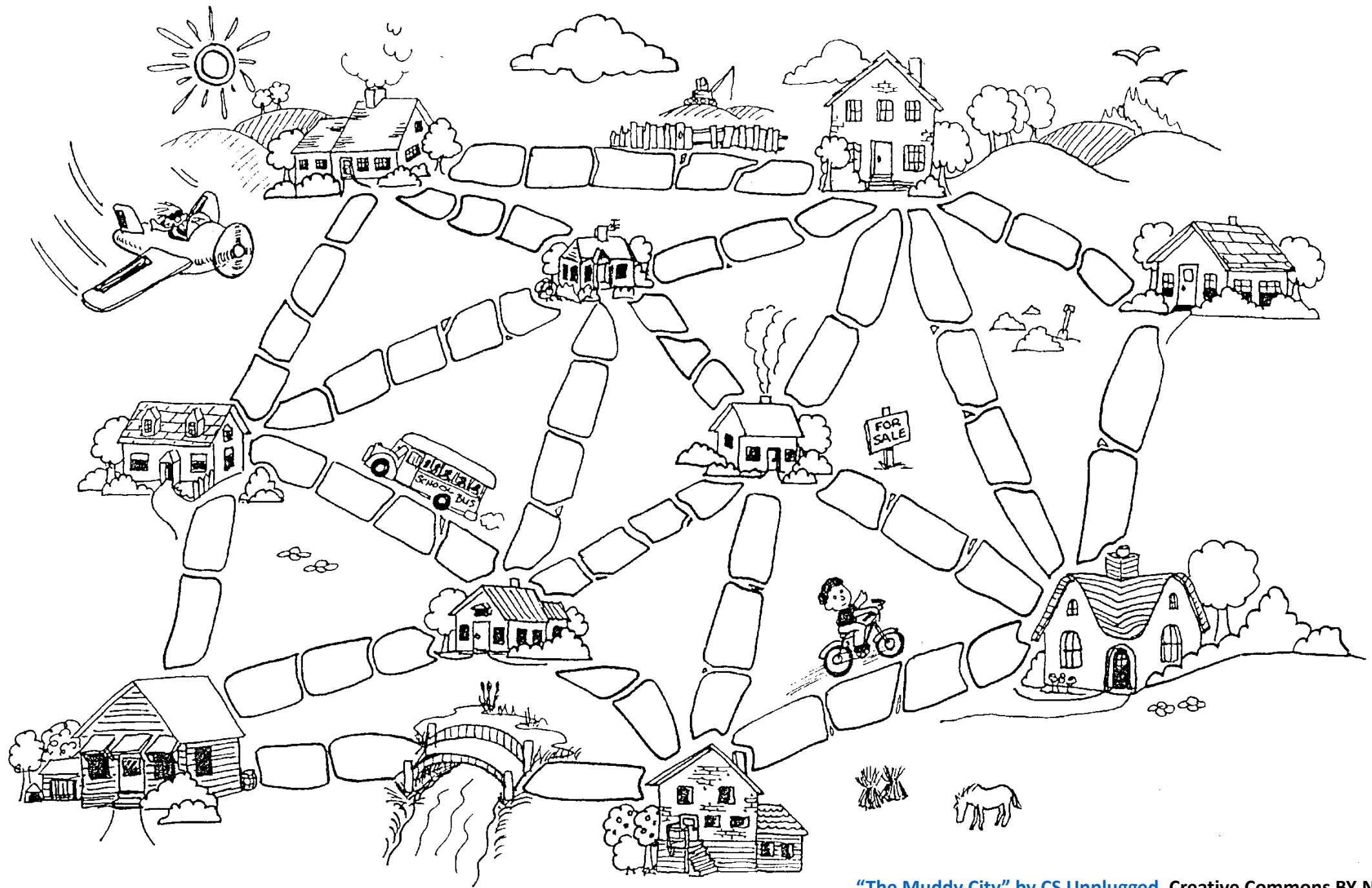
## Labeling:

- Vertex:
- Edge:

## Queries:

- Vertex:
- Edge:





# Minimum Spanning Tree Algorithms

**Input:** Connected, undirected graph  $G$  with edge weights (unconstrained, but must be additive)

**Output:** A graph  $G'$  with the following properties:

- $G'$  is a spanning graph of  $G$
- $G'$  is a tree (connected, acyclic)
- $G'$  has a minimal total weight among all spanning trees

