## Data Structures <br> Balanced Binary Search Trees

Department of Computer Science

## Learning Objectives

Discuss the big picture problem with BSTs
Introduce the self-balancing BST

## BST Analysis

Every operation on a BST depends on the height of the tree.
... how do we relate $O(h)$ to $n$, the size of our dataset?

## BST Analysis

What is the max number of nodes in a tree of height $h$ ?

## BST Analysis

What is the $\boldsymbol{\operatorname { m i n }}$ number of nodes in a tree of height $h$ ?

## BST Analysis

A BST of $n$ nodes has a height between:
Lower-bound: $O(\log n)$


Upper-bound: $O(n)$


## Height-Balanced Tree

What tree is better?


Height balance: $b=\operatorname{height}\left(T_{R}\right)-\operatorname{height}\left(T_{L}\right)$
A tree is "balanced" if:

## BST Rotations (The AVL Tree)

We can adjust the BST structure by performing rotations.

These rotations:
1.
2.

## BST Rotations (The AVL Tree)

We can adjust the BST structure by performing rotations.


## Left Rotation



## Left Rotation



Right Rotation


Right Rotation


## Coding AVL Rotations

Two ways of visualizing:

1) Think of an arrow 'rotating' around the center
2) Recognize that there's a concrete order for rearrangements


Ex: Unbalanced at current (root) node and need to rotateLeft?
Replace current (root) node with it's right child.
Set the right child's left child to be the current node's right
Make the current node the right child's left child

AVL Rotation Practice


## AVL Rotation Practice



Somethings not quite right...

## LeftRight Rotation



## LeftRight Rotation



RightLeft Rotation

AVL Rotations

## AVL Rotations

Four kinds of rotations: (L, R, LR, RL)

1. All rotations are local (subtrees are not impacted)
2. The running time of rotations are constant
3. The rotations maintain BST property

## Goal:

AVL Rotation Practice


## AVL vs BST ADT

The AVL tree is a modified binary search tree that rotates when necessary

How does the constraint on balance affect the core functions?
Find

Insert

Remove

AVL Find


AVL Insertion

```
struct TreeNode {
    T key;
    unsigned height;
    TreeNode *left;
    TreeNode *right;
};
```



## AVL Insertion

## Insert (pseudo code):

1: Insert at proper place
2: Check for imbalance
3: Rotate, if necessary
4: Update height

```
struct TreeNode {
    T key;
    unsigned height;
    TreeNode *left
    TreeNode *right;
};
```



## Rebalancing on insert



## Theorem:

If an insertion occurred in subtrees $\mathrm{t}_{\mathbf{3}}$ or $\mathbf{t}_{4}$ and an imbalance was first detected at $\mathbf{t}$, then a $\qquad$ rotation about $\mathbf{t}$ restores the balance of the tree.

We gauge this by noting the balance factor of $t$ is ___ and the balance factor of $t$->right is $\qquad$ .

## Rebalancing on insert

## Theorem:

If an insertion occurred in subtrees $\mathbf{t}_{\mathbf{1}}$ or $\mathbf{t}_{\mathbf{2}}$ and an imbalance was first detected at $\mathbf{t}$, then a $\qquad$ rotation about $\mathbf{t}$ restores the balance of the tree.

We gauge this by noting the balance factor of $t$ is ___ and the balance factor of $t$->left is $\qquad$ .

## Rebalancing on insert



## Theorem:

If an insertion occurred in subtrees $\mathrm{t}_{2}$ or $\mathbf{t}_{\mathbf{3}}$ and an imbalance was first detected at $\mathbf{t}$, then a $\qquad$ rotation about $\mathbf{t}$ restores the balance of the tree.

We gauge this by noting the balance factor of $t$ is $\qquad$ and the balance factor of $t$->right is $\qquad$ .

## Rebalancing on insert



## Theorem:

If an insertion occurred in subtrees $\mathrm{t}_{2}$ or $t_{3}$ and an imbalance was first detected at $\mathbf{t}$, then a $\qquad$ rotation about $\mathbf{t}$ restores the balance of the tree.

We gauge this by noting the balance factor of $t$ is $\qquad$ and the balance factor of $t$->left is $\qquad$ .

Rebalancing on insert


AVL Insertion Practice

