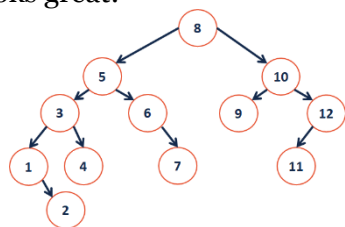


BTree Motivation

Big-O assumes uniform time for all operations, but this isn't always true.

However, seeking data from the cloud may take 100ms+.
...an $O(\lg(n))$ AVL tree no longer looks great:



Consider Instagram profile data:

How many profiles?		
How much data /profile?		
	AVL Tree	BTree
Tree Height		

BTree Motivations

Knowing that we have long seek times for data, we want to build a data structure with two (related) properties:

- 1.
- 2.

BTree_m



Goal: Build a tree that uses _____ /node!
...optimize the algorithm for your platform!

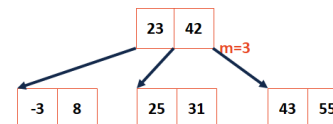
A **BTree of order m** is an m-way tree where:

1. All keys within a node are ordered.

BTree Insert, using m=5

...when a BTree node reaches **m** keys:

BTree Insert, m=3:



Great interactive visualization of BTrees:

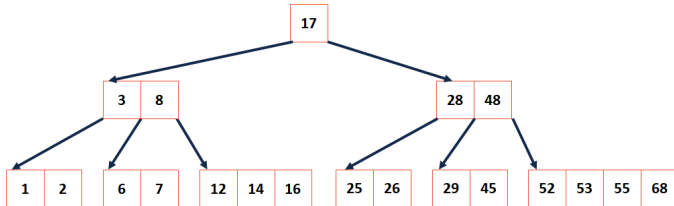
<https://www.cs.usfca.edu/~galles/visualization/BTree.html>

BTree Properties

For a BTree of order **m**:

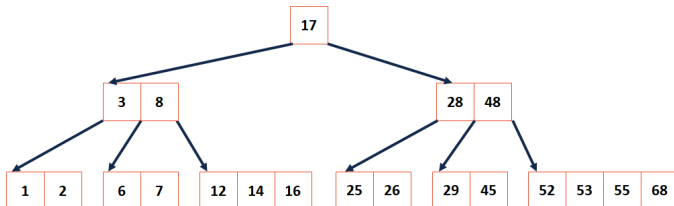
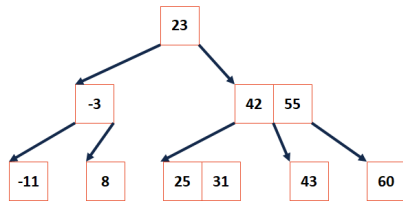
1. All keys within a node are ordered.
2. All leaves contain no more than **m-1** nodes.
3. All internal nodes have exactly **one more key than children**.
4. Root nodes can be a leaf or have **[2, m]** children.
5. All non-root, internal nodes have **[ceil(m/2), m]** children.
6. All leaves are on the same level.

Example BTree



What properties do we know about this BTree?

BTree Search



BTree .hpp

```

100 bool Btree<K, V>::_exists(BTreeNode & node, const K & key) {
101     unsigned i;
102     for (i=0; i<node.keys_ct_ && key<node.keys_[i]; i++) { }
103
104     if ( i < node.keys_ct_ && key == node.keys_[i] ) {
105         return true;
106     }
107
108     if ( node.isLeaf() ) {
109         return false;
110     } else {
111         BTreeNode nextChild = node._fetchChild(i);
112         return _exists(nextChild, key);
113     }
114 }

```

BTree Analysis

The height of the BTree determines maximum number of _____ possible in search data.

...and the height of our structure:

Therefore, the number of seeks is no more than: _____.

...suppose we want to prove this!

BTree Analysis

In our AVL Analysis, we saw finding an upper bound on the height (given **n**) is the same as finding a lower bound on the nodes (given **h**).

Goal: We want to find a relationship for BTrees between the number of keys (**n**) and the height (**h**).

CS 225 – Things To Be Doing:

1. Programming Exam 2 reservations open in CBTF
2. Mp_traversal due next Monday (Oct. 18)
3. lab_avl due Sunday
4. Daily POTDs are ongoing!