

Kruskal's Algorithm

```

Pseudocode for Kruskal's MST Algorithm
1  KruskalMST(G):
2    DisjointSets forest
3    foreach (Vertex v : G):
4      forest.makeSet(v)
5
6    PriorityQueue Q // min edge weight
7    foreach (Edge e : G):
8      Q.insert(e)
9
10   Graph T = (V, {})
11
12   while |T.edges()| < n-1:
13     Vertex (u, v) = Q.removeMin()
14     if forest.find(u) != forest.find(v):
15       T.addEdge(u, v)
16       forest.union( forest.find(u) ,
17                    forest.find(v) )
18
19   return T
  
```

Kruskal's Running Time Analysis

We have multiple choices on which underlying data structure to use to build the Priority Queue used in Kruskal's Algorithm:

Priority Queue Implementations:	Heap	Sorted Array
Building : 6-8		
Each removeMin : 13		

Based on our algorithm choice:

Priority Queue Implementation:	Total Running Time
Heap	
Sorted Array	

Reflections

Why would we prefer a Heap?

Why would we prefer a Sorted Array?

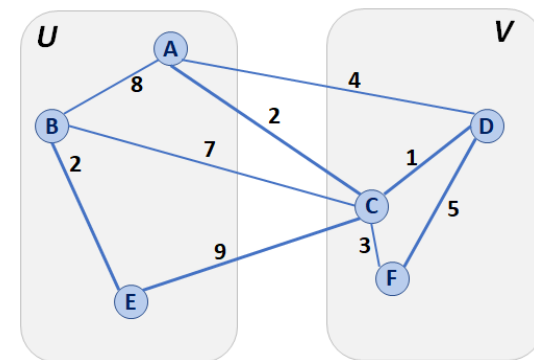
Partition Property

Consider an arbitrary partition of the vertices on G into two subsets U and V .

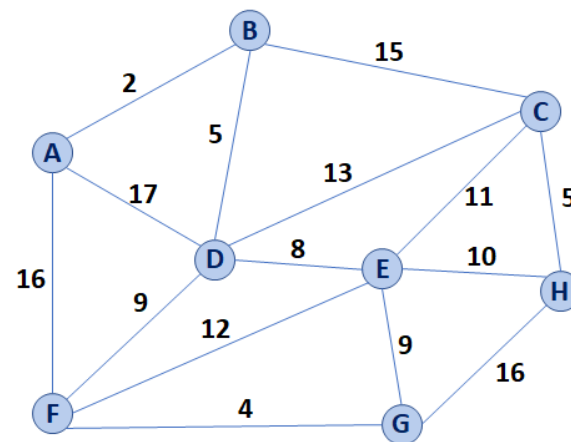
Let e be an edge of minimum weight across the partition.

Then e is part of some minimum spanning tree.

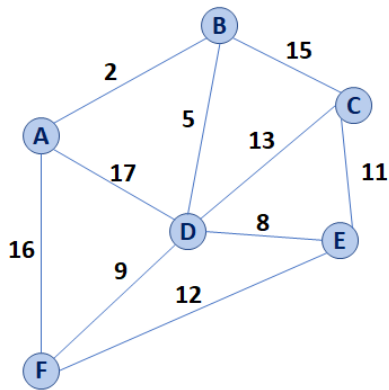
Proof in CS 374!



Partition Property Algorithm



Prim's Minimum Spanning Tree Algorithm



Pseudocode for Prim's MST Algorithm	
1	PrimMST(G, s):
2	Input: G, Graph;
3	s, vertex in G, starting vertex of algorithm
4	Output: T, a minimum spanning tree (MST) of G
5	
6	foreach (Vertex v : G):
7	d[v] = +inf
8	p[v] = NULL
9	d[s] = 0
10	
11	PriorityQueue Q // min distance, defined by d[v]
12	Q.buildHeap(G.vertices())
13	Graph T // "labeled set"
14	
15	repeat n times:
16	Vertex m = Q.removeMin()
17	T.add(m)
18	foreach (Vertex v : neighbors of m not in T):
19	if cost(v, m) < d[v]:
20	d[v] = cost(v, m)
21	p[v] = m
22	
23	return T

	Adj. Matrix	Adj. List
Heap		
Unsorted Array		

Running Time of MST Algorithms

Kruskal's Algorithm:

Prim's Algorithm:

Q: What must be true about the connectivity of a graph when running an MST algorithm?

...what does this imply about the relationship between **n** and **m**?

Kruskal's MST	Prim's MST

Q: Suppose we built a new heap that optimized the decrease-key operation, where decreasing the value of a key in a heap updates the heap in amortized constant time, or $O(1)^*$. How does that change Prim's Algorithm runtime?

Final big-O Running Times of classical MST algorithms:

Kruskal's MST	Prim's MST

CS 225 – Things To Be Doing:

- Programming Exam C is different than usual schedule:**
Exam: Monday, Dec 2 – Wednesday, Dec 4
- MP7 Released – Slightly different structure:**
Deadline on Tuesday, Dec. 3 for Part 1
- lab_ml in lab this week also due Tuesday, Dec. 3!
- Daily POTDs are ongoing for +1 point /problem but pausing over break