

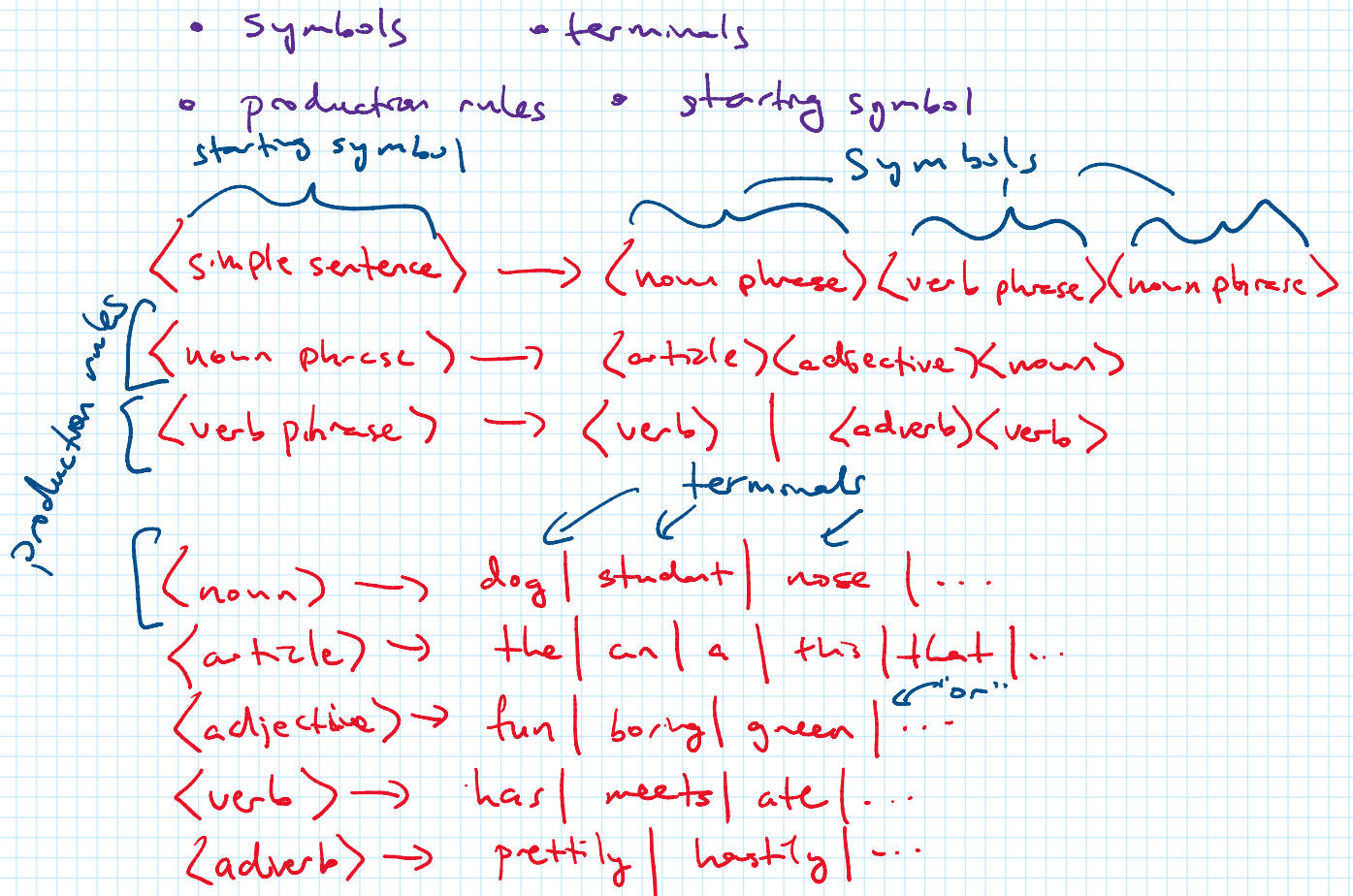
CS 173 Lecture 15a: Context-Free Grammars

Informally: a recursive way of construct strings " " see Chapter

Historically:

- Linguistics
 - attempt to model human language via constructing sentences out of recursive phrases.
- Computer Science
 - a way to construct new programming languages
 - parsing \rightarrow compiling

CFG consists of



CFG: starting w/ a starting symbol:

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 replace symbols w/ one of the options in the production rule
 then recurse until only terminals are left

<simple sentence>

→ <np><vp><np> ← choice: <v> instead of <adv><v>

→ <art><adj><n> <v> <art><adj><n>

→ the fun dog has a green nose.

← ↑ → ...
 choices

<ss>
 → <np><vp><np>

→ <art><adj><n> <adv><v> <art><adj><n>

→ This funny student prettily ate an boring homework.

valid string in the grammar.

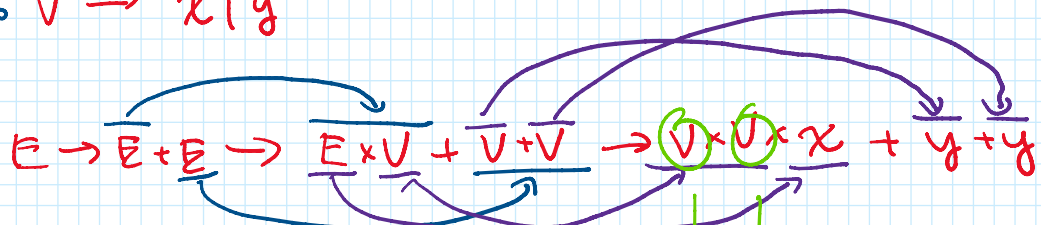
bad English "an boring"
 makes no sense

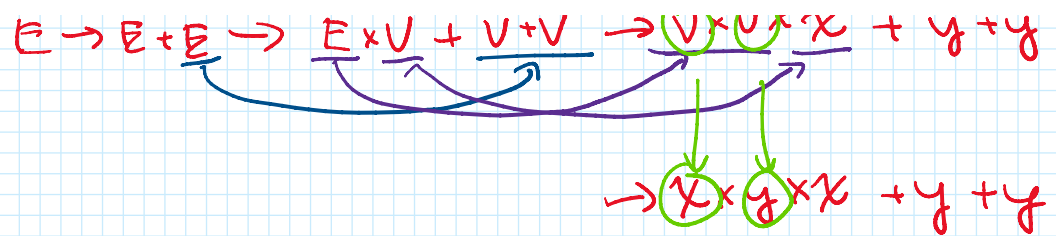
~~_____ X _____~~

Some conventions: starting symbol's prod rule written first
 symbols are often upper case letters
 terminals are lower case letters or numbers
 (not always)

expression
 starting symbol → $E \rightarrow E + E \mid E \times V \mid V + V \mid V \times V$
 terminal

variables
 $V \rightarrow x \mid y$





"CFG for 'simple' arithmetic expressions
w/ variables x, y "