



CS 173: Discrete Structures

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Office Hour: Wed. 1-2, 2215 SC

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Agenda

- Quickly revisiting propositional logic
- What you can do with what you now know
 - Logic
 - Number Theory
- Set Theory
 - Section 2.1 : Definitions and Notation
 - Section 2.2 : Operations on Sets





Revisiting Propositional Logic

- Informally, a “predicate” $P(x)$ is a function
 - It returns a Boolean value (i.e. T or F)
 - The value depends on an input variable(s)
 - It is used with quantifiers

- Example:

“If you don't know where you are going, you will wind up somewhere else.”

$\forall x \in \text{People}, \underbrace{\neg \text{KnowsWhereGoing}(x)} \rightarrow \underbrace{\text{EndsUpElsewhere}(x)}$

- A predicate must take on a TF value
- A propositional variable must take on a TF value
 - e.g. saying $Q = \text{“somewhere else”}$ would be wrong

$P(x, y)$ capital

$p \vee q$

lower case





Things you can do with predicate logic

- Build an artificially intelligent super-brain! (maybe)
- "Intelligence is ten million rules."
- Douglas Lenat CEO of Cycorp
- Cyc™ is an automatic reasoning system consisting of
 - an inference engine
 - a knowledge-base of rules
 - E.g. "All trees are plants"
- "one of the most controversial endeavors of the artificial intelligence history" -- Wikipedia entry



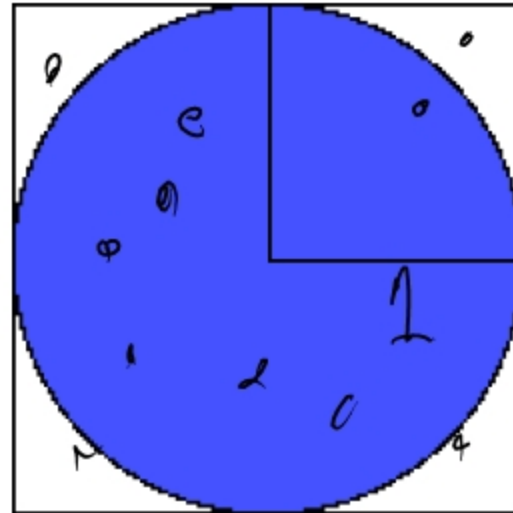


Things you can do with number theory

- Implement a Random Number Generator

- Lots of uses

- Cryptography
- Randomized algorithms
- Monte Carlo methods



- Instigating a financial system collapse (maybe)

- Can you really write a program that generates RNs?

points in circle

—————
points

~ area circ
~
—————
area sq^r

~ π
~ ———
4





Random Number Generation

- Linear Congruential generators are common
 - Most C/C++ standard library “rand()” implementations
 - “random()” and “rand48()” are usually better if available
- Introduced by D.H. Lehmer (1949)
Sequence $X_{n+1} = (aX_n + c) \bmod m$
 - $m > 0$ (usually a prime)
 - $0 < a < m$
 - $0 \leq c < m$
 - $0 \leq X_0 < m$
- *How many distinct values can be generated?*

m , they will
be in $[0, m-1]$



Random Number Generation

$$X_{n+1} = (aX_n + c) \bmod m$$

- How many distinct values can be generated? ✓

- What should we try to maximize?

"period" = how many values do you generate before repeating

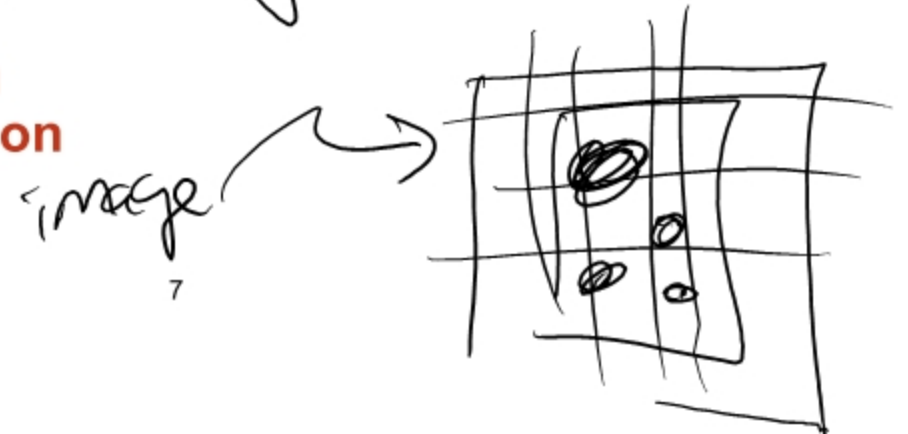
- Is choosing X_0 important ?

X_0 = "seed"

clock time^{often} used for seed

- **US Patent 5732138 - Method for seeding a pseudo-random number generator with a cryptographic hash of a digitization of a chaotic system**

Lava lamp





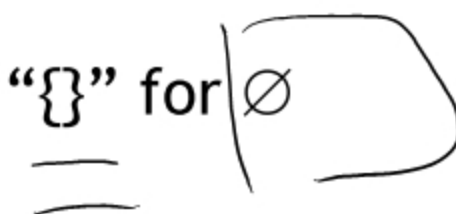
Set Theory - Definitions and notation

A set is an unordered collection of elements.

- Order and repetition are irrelevant
 $\{1, 1, 2, 3, 3\} = \{1, 2, 3\}$ since repetition is irrelevant.
 $\{1, 2, 3\} = \{3, 2, 1\}$ since sets are unordered.
- Sets can be finite or infinite
 $\{1, 2, 3, \dots\}$ is a way we denote an infinite set
- $\emptyset = \{\}$ is the empty set, or the set containing no elements.

Note: $\emptyset \neq \{\emptyset\}$

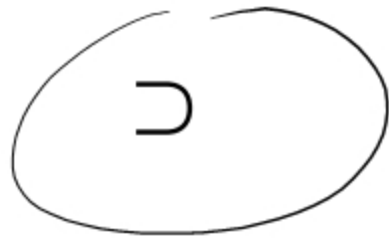
Also, don't use the notation " $\{\}$ " for



use that



Set Theory - Definitions and notation



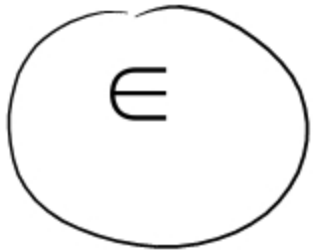
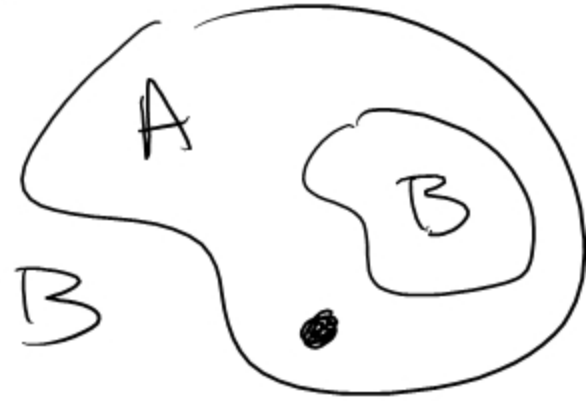
$$A \supset B$$

A proper superset of

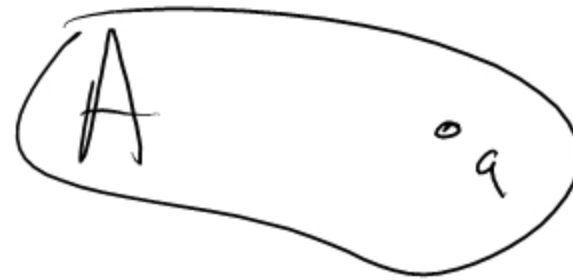
\supseteq

$$A \supseteq B$$

, A could be $\equiv B$



$$a \in A$$



: and |

→ "such that"

$$\{x : x > 0\}$$



Set Theory - Definitions and notation

Quiz time:

Is $\{x\} \subseteq \{x\}$? **Yes**

Is $\{x\} \in \{x, \{x\}\}$? **Yes**

Is $\{x\} \subseteq \{x, \{x\}\}$? **Yes**

Is $\{x\} \in \{x\}$? **No**





Set Theory - Ways to define sets

- Explicitly: {John, Paul, George, Ringo}
- Implicitly: {1,2,3,...}, or {2,3,5,7,11,13,17,...}
- Set builder: { x : x is prime }, { x | x is odd }. In general { x : $P(x)$ is true }, where $P(x)$ is some description of the set.

: and | are read
"such that" or
"where"

Ex. Let $D(x,y)$ denote "x is divisible by y."

Give another name for

$$\{ x : \forall y ((y > 1) \wedge (y < x)) \rightarrow \neg D(x,y) \}.$$

Primes

Can we use **any** predicate P to define a set

$$S = \{ x : P(x) \}?$$





Set Theory - Russell's Paradox

Can we use **any** predicate P to define a set
 $S = \{ x : P(x) \}$?

Define $S = \{ x : x \text{ is a set where } x \notin x \}$ **No!**

Then, if $S \in S$, then by defn of S , $S \notin S$.

But, if $S \notin S$, then by defn of S , $S \in S$.

There is a town with a barber who shaves all the people (and only the people) who don't shave themselves.

Who shaves the barber?





Set Theory - Cardinality

If S is finite, then the *cardinality* of S , $|S|$, is the number of distinct elements in S .

- If $S = \{1, 2, 3\}$, $|S| = 3$. $|S| = 3$
- If $S = \{3, 3, 3, 3, 3\}$, $|S| = 1$. $|S| = 1$
- If $S = \emptyset$, $|S| = 0$. $|S| = 0$
- If $S = \{\emptyset, \{\emptyset\}, \{\emptyset, \{\emptyset\}\}$, $|S| = 3$. $|S| = 3$
- If $S = \{0, 1, 2, 3, \dots\}$, $|S|$ is infinite. (more on this later)





Set Theory - Power sets

If S is a set, then the *power set* of S is
 $2^S = \{x : x \subseteq S\}$.

aka $P(S)$

We say, "P(S) is the set of all subsets of S."

If $S = \{a\}$, $2^S = \{\emptyset, \{a\}\}$.

If $S = \{a, b\}$, $2^S = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$.

If $S = \emptyset$, $2^S = \{\emptyset\}$.

If $S = \{\emptyset, \{\emptyset\}\}$, $2^S = \{\emptyset, \{\emptyset\}, \{\{\emptyset\}\}, \{\emptyset, \{\emptyset\}\}\}$.

$2^S = \{\emptyset, \{\emptyset\}, \{\{\emptyset\}\}, \{\emptyset, \{\emptyset\}\}\}$
Fact: if S is finite, $|2^S| = 2^{|S|}$. (if $|S| = n$, $|2^S| = 2^n$)

$$|2^S| = 2^{|S|}$$

power set S
 2^S or $P(S)$

$$S = \{a\}$$

$$2^S = \{\emptyset, \{a\}\}$$

$$S = \{a, b\}$$

$$2^S = \{\emptyset, \{a\}, \{b\},$$

$$\{a, b\}\}$$

$$\text{if } S = \emptyset, 2^S = \{\emptyset\}$$





Set Theory - Cartesian Product

The *Cartesian Product* of two sets A and B is:

$$A \times B = \{ \langle a, b \rangle : a \in A \wedge b \in B \}$$

ordered pairs

If $A = \{ \text{Charlie, Lucy, Linus} \}$, and
 $B = \{ \text{Brown, VanPelt} \}$, then

$$A \times B = \{ \langle \text{Charlie, Brown} \rangle, \langle \text{Lucy, Brown} \rangle, \langle \text{Linus, Brown} \rangle, \langle \text{Charlie, VanPelt} \rangle, \langle \text{Lucy, VanPelt} \rangle, \langle \text{Linus, VanPelt} \rangle \}$$

n-tuple

$$A_1 \times A_2 \times \dots \times A_n = \{ \langle a_1, a_2, \dots, a_n \rangle : a_1 \in A_1, a_2 \in A_2, \dots, a_n \in A_n \}$$

A, B finite $\rightarrow |A \times B| = ?$

- a) $A \times B$
- b) $|A| + |B|$
- c) $|A + B|$
- d) $|A| |B|$

$$|A \times B| = |A| |B|$$

$$A = \{ 1, 2 \}$$

$$B = \{ a, b \}$$

$$A \times B =$$

$$\{ \langle 1, a \rangle, \langle 1, b \rangle, \langle 2, a \rangle, \langle 2, b \rangle \}$$

$$C = \{ 9, 10 \}$$

$$A \times B \times C =$$

$$\{ \langle 1, a, 9 \rangle, \langle 1, a, 10 \rangle, \dots \}$$



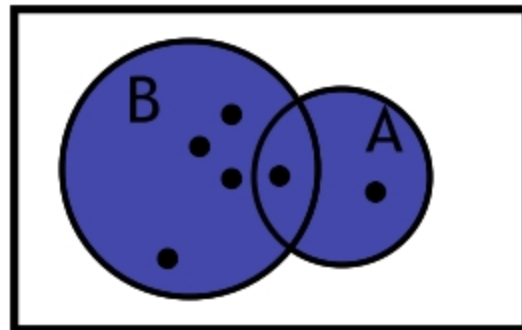
Set Theory - Operators

The *union* of two sets A and B is:

$$A \cup B = \{x : x \in A \vee x \in B\}$$

If $A = \{\text{Charlie, Lucy, Linus}\}$, and
 $B = \{\text{Lucy, Desi}\}$, then

$$A \cup B = \{\text{Charlie, Lucy, Linus, Desi}\}$$





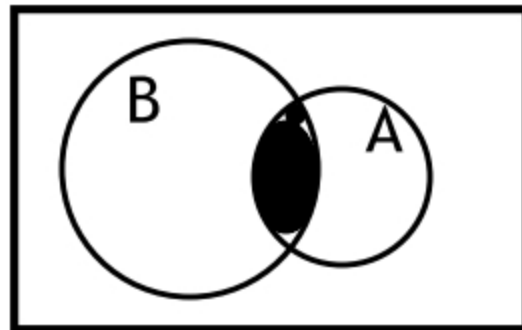
Set Theory - Operators

The *intersection* of two sets A and B is:

$$A \cap B = \{x : x \in A \wedge x \in B\}$$

If $A = \{\text{Charlie, Lucy, Linus}\}$, and
 $B = \{\text{Lucy, Desi}\}$, then

$$A \cap B = \{\text{Lucy}\}$$





Set Theory - Operators

The *intersection* of two sets A and B is:

$$A \cap B = \{x : x \in A \wedge x \in B\}$$

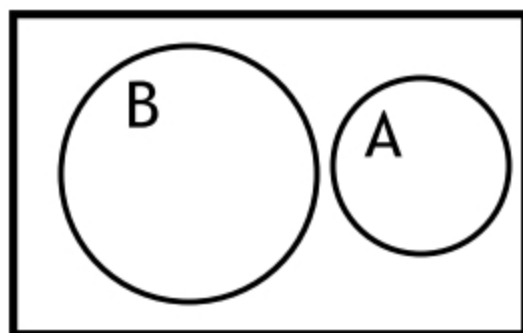
If $A = \{x : x \text{ is a US president}\}$, and

$B = \{x : x \text{ is in this room}\}$, then

$$A \cap B = \{x : x \text{ is a US president in this room}\} = \emptyset$$



A and B are disjoint



Sets whose intersection is empty are called *disjoint sets*



Set Theory - Operators

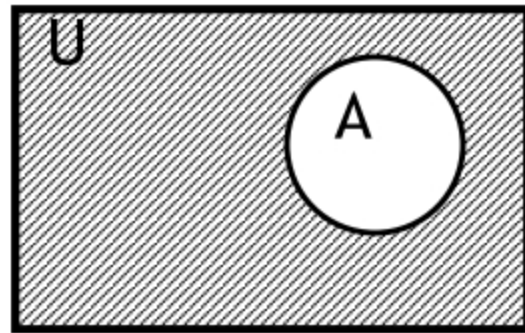
The complement of a set A is:

$$\overline{A} = \{x : x \notin A\}$$



If $A = \{x : x \text{ is bored}\}$, then

$$\overline{A} = \{x : x \text{ is not bored}\} = \emptyset$$



$$\begin{aligned} \overline{\emptyset} &= U \\ \text{and} \\ \overline{U} &= \emptyset \end{aligned}$$

U = "universe"

$$\text{Let } U = \mathbb{N}$$

$$A = \{0, 1, 2\}$$

$$\overline{A} = \{3, 4, \dots\}$$

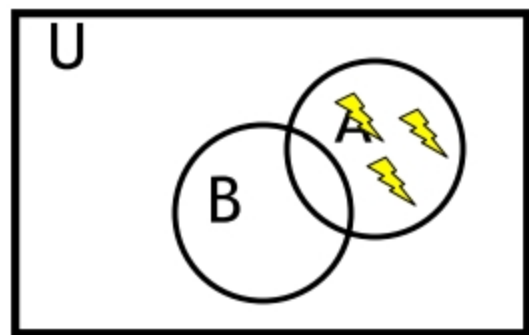
$$\begin{aligned} \overline{\emptyset} &= U \\ \overline{U} &= \emptyset \end{aligned}$$





Set Theory - Operators

The set difference, $A - B$, is:



$$A - B = \{x : x \in A \wedge x \notin B\}$$

$$A - B = \underbrace{A \cap \bar{B}}$$

$$A = \{1, 2, 3\}$$

$$B = \{2\}$$

$$A - B = \{1, 3\}$$

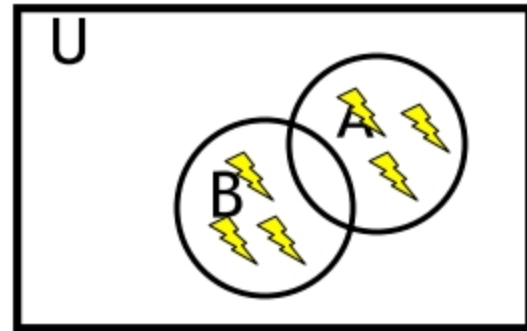


Set Theory - Operators

The *symmetric difference*, $A \oplus B$, is:

$$A \oplus B = \{x : (x \in A \wedge x \notin B) \vee (x \in B \wedge x \notin A)\}$$

$$= (A - B) \cup (B - A)$$



$$A = \{1, 2, 3\}$$

$$B = \{3, 4\}$$

$$A \oplus B =$$

$$\{1, 2, 4\}$$