

Recursion and Structural Induction Homework

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Fall 2020

Problem 1. Consider a set S defined recursively: $1 \in S$, and if $x \in S$ then $-x \in S$ and $2x \in S$.

- What is the set S ? (Give an equivalent non-recursive description, e.g. in set-builder notation or using sets we've named in the past.)
- A recursive definition for a set is called *ambiguous* if there is any element in the set that can be constructed from the given base cases and constructors in more than one way.¹ Prove the above definition for S is ambiguous.
- Is it possible to define a recursive function $f : S \rightarrow \mathbb{N}$ which counts how many constructor steps were used to produce a given element of S ? Why or why not? (For example, we might want $f(1) = 0$ because 1 is in the base case so it looks like it doesn't use any applications of the constructor step.)
- Come up with a new *unambiguous* recursive definition for the same set S .²

¹ Being ambiguous does not make a definition invalid, but it sometimes makes it less useful.

² Hint: You may want multiple base case elements instead of just 1.

Problem 2. Consider a set S defined recursively: $-2 \in S$, and if $x \in S$ then $2x \in S$ and $x^2 - 1 \in S$. You may use without proof that every element of S is an integer.

- Prove $5 \notin S$.
- Prove $\forall y \in S (|y| \geq 2)$ by structural induction.
- Prove that for every $y \in S$, if $y < 0$ then y is even.³

³ Hint: use the result from part (b).

Problem 3.

- Give a recursive definition for a function $rev : \{0,1\}^* \rightarrow \{0,1\}^*$ which reverses a binary string. For example, $rev(001111) = 111100$.
- Prove by structural induction that $rev(st) = rev(t)rev(s)$.

Problem 4. Prove by structural induction: If T is a binary tree, then for any labeling of T 's vertices using blue and orange such that the root is blue and every leaf is orange, there exists some blue vertex that has an orange child.