

CS 173, Fall 2015
Examlet 11, Part A

NETID:

FIRST:

LAST:

Discussion: Thursday 2 3 4 5 Friday 9 10 11 12 1 2

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01 Chop( $a_1, \dots, a_n; b_1, \dots, b_n$ )  \\ input is 2 lists of n integers, n is a power of 2
02     if ( $n = 1$ )
03         return  $a_1 b_1$ 
04     else
05          $p = \frac{n}{2}$ 
06          $rv = \text{Chop}(a_1, \dots, a_p, b_1, \dots, b_p)$ 
07          $rv = rv + \text{Chop}(a_1, \dots, a_p, b_{p+1}, \dots, b_n)$ 
08          $rv = rv + \text{Chop}(a_{p+1}, \dots, a_n, b_{p+1}, \dots, b_n)$ 
09          $rv = rv + \text{Chop}(a_{p+1}, \dots, a_n, b_1, \dots, b_p)$ 
10     return rv

```

1. (5 points) Suppose that $T(n)$ is the running time of Chop on an input array of length n . Give a recursive definition of $T(n)$. Assume that dividing the list in half takes $O(n)$ time.

Solution:

$$T(1) = c$$

$$T(n) = 4T(n/2) + dn + f$$

2. (4 points) What is the height of the recursion tree for $T(n)$, assuming n is a power of 2?

Solution: $\log_2 n$

3. (3 points) What is the amount of work (aka sum of the values in the nodes) at level k of this tree?

Solution: There are 4^k nodes, each containing $f + dn/2^k$. So the total work is $4^k f + 2^k dn$

4. (3 points) How many leaves are in the recursion tree for $T(n)$? (Simplify your answer.)

Solution: $4^{\log_2 n} = 4^{\log_4 n \log 2^4} = n^{\log 2^4} = n^2$

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01 Crunch(k,n)  \ \ inputs are positive integers
02     if (n = 1) return k
03     else if (n = 2) return k2
04     else
05         half = ⌊n/2⌋
06         answer = Crunch(k,half)
07         answer = answer*answer
08         if (n is odd)
09             answer = answer*k
10         return answer

```

1. (5 points) Suppose $T(n)$ is the running time of Crunch. Give a recursive definition of $T(n)$.

Solution: $T(1) = c, T(2) = d$

$$T(n) = T(n/2) + f$$

2. (4 points) What is the height of the recursion tree for $T(n)$? (Assume that n is a power of 2.)

Solution: $\log_2 n - 1$

3. (3 points) How many leaves are in the recursion tree for $T(n)$?

Solution: One.

4. (3 points) What is the big-Theta running time of Crunch?

Solution: $\Theta(\log n)$

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01 Process ( $a_1, \dots, a_n$ : array of integers)
02   if ( $n = 1$ )
03     if ( $a_1 > 8$ ) return true
04     else return false
05   else if (Process( $a_1, \dots, a_{n-1}$ ) is true and Process( $a_2, \dots, a_n$ ) is true)
06     return true
07   else return false

```

1. (3 points) If Process returns true, what must be true of the values in the input array?

Solution: The values in the input array must all be greater than 8.

2. (5 points) Give a recursive definition for $T(n)$, the running time of Process on an input of length n , assuming it takes constant time to set up the recursive calls in line 05.

Solution:

$$T(1) = c$$

$$T(n) = 2T(n-1) + d$$

3. (3 points) What is the height of the recursion tree for $T(n)$?

Solution: n

4. (4 points) What is the big-theta running time of Process?

Solution: $\Theta(2^n)$

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01 Twiddle( $a_0, \dots, a_{n-1}$ )  \ \ input is an array of n integers
02     if ( $n = 2$  and  $a_0 > a_1$ )
03         swap( $a_0, a_1$ )  \ \ interchange the values at positions 0 and 1 in the array
04     else if ( $n > 2$ )
05          $p = \lfloor \frac{n}{4} \rfloor$ 
06          $q = \lfloor \frac{n}{2} \rfloor$ 
07          $r = p + q$ 
08         Twiddle( $a_0, \dots, a_q$ )  \ \ constant time to make smaller array
09         Twiddle( $a_{q+1}, \dots, a_{n-1}$ )  \ \ constant time to make smaller array
10         Twiddle( $a_p, \dots, a_r$ )  \ \ constant time to make smaller array

```

1. (5 points) Suppose that $T(n)$ is the running time of Twiddle on an input array of length n . Give a recursive definition of $T(n)$.

Solution:

$$T(1) = c, T(2) = d$$

$$T(n) = 3T(n/2) + f$$

2. (4 points) What is the height of the recursion tree for $T(n)$, assuming n is a power of 2?

Solution: $\log_2 n - 1$

3. (3 points) What is the amount of work (aka sum of the values in the nodes) at level k of this tree?

Solution: $f \cdot 3^k$

4. (3 points) How many leaves are in the recursion tree for $T(n)$? (Simplify your answer.)

Solution: $3^{\log_2 n - 1} = 1/3(3^{\log_2 n}) = 1/3(3^{\log_3 n \log_2 3}) = 1/3 \cdot n^{\log_2 3}$

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01 MyFunc( $a_1, \dots, a_n$ )  \ \ input is an array of n integers
02   for  $i := 1$  to  $n - 1$ 
03      $min := i$ 
04     for  $j := i$  to  $n$ 
05       if  $a_j < a_{min}$  then  $min := j$ 
06     swap( $a_i, a_{min}$ )  \ \ interchange the values at positions  $i$  and  $min$  in the array

```

1. (4 points) If the input is 10, 5, 2, 3, 8, what are the array values after two iterations of the outer loop?

Solution: After the second iteration, it contains 2, 3, 10, 5, 8.

2. (4 points) Let $T(n)$ be the number of times that line 5 is executed. Express $T(n)$ using summation notation, directly following the structure of the code.

Solution: The i th time through the outer loop, the inner loop runs $n - i + 1$ times. So the total number of times that line 5 executes is:

$$\sum_{i=1}^{n-1} (n - i + 1)$$

3. (4 points) Find an (exact) closed form for $T(n)$. Show your work.

Solution: If we break apart the sum and then substitute in a new index variable $p = n - i$ we get:

$$\sum_{i=1}^{n-1} (n - i + 1) = (n - 1) + \sum_{i=1}^{n-1} (n - i) = (n - 1) + \sum_{p=1}^{n-1} p = (n - 1) + \frac{n(n - 1)}{2}$$

Simplifying, we get

$$(n - 1) + \frac{n(n-1)}{2} = n - 1 + \frac{1}{2}n^2 - \frac{1}{2}n = \frac{1}{2}n^2 + \frac{1}{2}n + 1$$

4. (3 points) What is the big-theta running time of MyFunc?

Solution: $\Theta(n^2)$

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00 Kitty( $a_1, \dots, a_n$ ) : list of  $n$  positive integers,  $n \geq 2$ )
01   if ( $n = 2$ ) return  $|a_1 - a_2|$ 
02   else
03       bestval = 0
04       for  $k = 1$  to  $n$ 
05           newval = Kitty( $a_1, a_2, \dots, a_{k-1}, a_{k+1}, \dots, a_n$ )  \\ \ constant time to remove  $a_k$ 
06           if (newval > bestval) bestval = newval
07       return bestval

```

1. (3 points) Describe (in English) what Kitty computes.

Solution: Kitty computes the largest Kitty difference between two values in the list. Or, equivalently, the largest value minus the smallest value.

2. (5 points) Suppose that $T(n)$ is the running time of Kitty on an input list of length n . Give a recursive definition of $T(n)$.

Solution:

$$T(2) = c$$

$$T(n) = nT(n-1) + d$$

3. (3 points) What is the height of the recursion tree for $T(n)$?

Solution: $n - 2$

4. (4 points) How many leaf nodes are there in the recursion tree for $T(n)$?

Solution: $\frac{n!}{2}$