

# CS477 Formal Software Development Methods

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Slides based in part on previous lectures by Mahesh Vishwanathan, and by Gul Agha

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## Proof Rules

Will give Sequent version of Natural Deduction rules  
 All rules from Propositional Logic included

$$\frac{\Gamma \vdash \psi[t/x]}{\Gamma \vdash \exists x.\psi} \text{ Ex I}$$

$$\frac{\Gamma \vdash \exists x.\psi \quad \Gamma \cup \{(\psi[y/x])\} \vdash \varphi}{\Gamma \vdash \varphi} \text{ Ex E}$$

provided  
 $y \notin \text{fv}(\varphi) \cup (\text{fv}(\psi) \setminus \{x\}) \cup \bigcup_{\psi' \in \Gamma} \text{fv}(\psi')$

$$\frac{\Gamma \vdash \psi[y/x]}{\Gamma \vdash \forall x.\psi} \text{ All I}$$

$$\frac{\Gamma \vdash \forall x.\psi \quad \Gamma \cup \{\psi[t/x]\} \vdash \varphi}{\Gamma \vdash \varphi} \text{ All E}$$

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## Example

Show

$$\frac{}{\{\} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)}$$

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## Example of Failure

Let's try to show

$$\frac{\frac{\{\}}{\{\}} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)}{\{\}} \text{Imp I}$$

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## Example of Failure

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## Example of Failure

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## Simple Imperative Language

- We will give rules for simple imperative language

$$\langle \text{command} \rangle ::= \langle \text{variable} \rangle := \langle \text{term} \rangle$$

$$| \langle \text{command} \rangle; \dots; \langle \text{command} \rangle$$

$$| \text{if } \langle \text{statement} \rangle \text{ then } \langle \text{command} \rangle \text{ else } \langle \text{command} \rangle$$

$$| \text{while } \langle \text{statement} \rangle \text{ do } \langle \text{command} \rangle$$

- Could add more features, like for-loops

## Substitution

- Notation:  $P[e/v]$  (sometimes  $P[v \rightarrow e]$ )
- Meaning: Replace every  $v$  in  $P$  by  $e$
- Example:

$$(x + 2)[y - 1/x] = ((y - 1) + 2)$$

## The Assingment Rule

$$\frac{}{\{P[e/x]\} x := e \{P\}}$$

Example:

$$\frac{}{\{ \quad ? \quad \} x := y \{ x = 2 \}}$$

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$$\frac{}{\{P[e/x]\} x := e \{P\}}$$

Examples:

$$\frac{}{\{y = 2\} x := y \{x = 2\}}$$

$$\frac{}{\{y = 2\} x := 2 \{y = x\}}$$

$$\frac{}{\{x + 1 = n + 1\} x := x + 1 \{x = n + 1\}}$$

$$\frac{}{\{2 = 2\} x := 2 \{x = 2\}}$$

## The Assignment Rule – Your Turn

- What is the weakest precondition of

$$x := x + y \{x + y = wx\}?$$

$$\left\{ \begin{array}{c} ? \\ x := x + y \\ \{x + y = wx\} \end{array} \right\}$$

## The Assignment Rule – Your Turn

- What is the weakest precondition of

$$x := x + y \{x + y = wx\}?$$

$$\left\{ \begin{array}{c} (x + y) + y = w(x + y) \\ x := x + y \\ \{x + y = wx\} \end{array} \right\}$$

## Precondition Strengthening

$$\frac{(P \Rightarrow P') \{P'\} C \{Q\}}{\{P\} C \{Q\}}$$

- Meaning: If we can show that  $P$  implies  $P'$  (i.e.  $P \Rightarrow P'$ ) and we can show that  $\{P'\} C \{Q\}$ , then we know that  $\{P\} C \{Q\}$
- $P$  is **stronger** than  $P'$  means  $P \Rightarrow P'$

## Precondition Strengthening

- Examples:

$$\frac{x = 3 \Rightarrow x < 7 \quad \{x < 7\} x := x + 3 \{x < 10\}}{\{x = 3\} x := x + 3 \{x < 10\}}$$

$$\frac{True \Rightarrow (2 = 2) \quad \{2 = 2\} x := 2 \{x = 2\}}{\{True\} x := 2 \{x = 2\}}$$

$$\frac{x = n \Rightarrow x + 1 = n + 1 \quad \{x + 1 = n + 1\} x := x + 1 \{x = n + 1\}}{\{x = n\} x := x + 1 \{x = n + 1\}}$$

## Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} x := x * x \{x < 25\}}{\{x = 3\} x := x * x \{x < 25\}}$$

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## Which Inferences Are Correct?

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