

CS477 Formal Software Development Methods

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Proof Rules

Will give Sequent version of Natural Deduction rules
All rules from Propositional Logic included

$$\frac{\Gamma \vdash \psi[t/x]}{\Gamma \vdash \exists x.\psi} \text{Ex I}$$

$$\frac{\Gamma \vdash \exists x.\psi \quad \Gamma \cup \{\psi[y/x]\} \vdash \varphi}{\Gamma \vdash \varphi} \text{Ex E}$$

provided

$$y \notin \text{fv}(\varphi) \cup (\text{fv}(\psi) \setminus \{x\}) \cup \bigcup_{\psi' \in \Gamma} \text{fv}(\psi')$$

$$\frac{\Gamma \vdash \psi[y/x]}{\Gamma \vdash \forall x.\psi} \text{All I}$$

$$\frac{\Gamma \vdash \forall x.\psi \quad \Gamma \cup \{\psi[t/x]\} \vdash \varphi}{\Gamma \vdash \varphi} \text{All E}$$

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$$y \notin (\text{fv}(\psi) \setminus \{x\}) \cup \bigcup_{\psi' \in \Gamma} \text{fv}(\psi')$$

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Example

Show

$$\frac{}{\{ \} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)}$$

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Example

Show

$$\frac{\{(\exists x. \forall y. x \leq y)\} \vdash \forall x. \exists y. y \leq x}{\{ \} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)} \text{Imp I}$$

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Example

Show

$$\frac{\frac{\frac{\frac{\frac{\{(\exists x. \forall y. x \leq y)\} \vdash \exists y. y \leq \textcolor{red}{x}}{\{(\exists x. \forall y. x \leq y)\} \vdash \forall x. \exists y. y \leq \textcolor{red}{x}} \text{All I}}{\{(\exists x. \forall y. x \leq y)\} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)} \text{Imp I}}{}}{}}$$

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Example

Show

$$\frac{\frac{\frac{\frac{\frac{\{(\exists x. \forall y. x \leq y)\} \vdash \exists x. \forall y. \textcolor{red}{x} \leq y}{\{(\exists x. \forall y. x \leq y)\} \vdash \exists y. \textcolor{red}{y} \leq \textcolor{red}{x}} \text{All I}}{\frac{\frac{\frac{\{(\exists x. \forall y. x \leq y)\} \vdash \exists y. y \leq \textcolor{red}{x}}{\{(\exists x. \forall y. x \leq y)\} \vdash \forall x. \exists y. y \leq \textcolor{red}{x}} \text{All I}}{\frac{\frac{\{(\exists x. \forall y. x \leq y)\} \vdash \forall x. \exists y. y \leq \textcolor{red}{x}}{\{(\exists x. \forall y. x \leq y)\} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)} \text{Imp I}}{}}{}}{}}$$

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Example of Failure

Let's try to show

$$\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{Imp I}$$

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Example of Failure

Let's try to show

$$\frac{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y}{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y} \text{Ex I}}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{Imp I}$$

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Let's try to show

$$\frac{\frac{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y}{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y} \text{All I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{Ex I}} \text{Imp I}}$$

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Example of Failure

Let's try to show

$$\frac{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall x. \exists y. y \leq x \text{ Hyp}}{\frac{\frac{\{\forall x. \exists y. y \leq x; \exists y. y \leq x\} \vdash \forall y. y \leq x}{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y} \text{All E}}{\frac{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y}{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y} \text{All I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{Ex I}} \text{Imp I}}}}$$

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Example of Failure

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Example of Failure

Let's try to show

$$\begin{array}{c}
 \frac{\text{Hyp} \quad \left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \end{array} \right\} \vdash \exists y. y \leq x}{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \\ y \leq x \end{array} \right\} \vdash y \leq x} \text{ Hyp } r \\
 \frac{}{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \end{array} \right\} \vdash \forall y. y \leq x} \text{ Ex } \\
 \frac{\text{Hyp} \quad \left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \end{array} \right\} \vdash \forall y. y \leq x}{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \\ \forall y. y \leq x \end{array} \right\} \vdash \forall y. y \leq x} \text{ All } I \\
 \frac{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \end{array} \right\} \vdash \forall y. z \leq y}{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \\ \forall y. z \leq y \end{array} \right\} \vdash \forall y. z \leq y} \text{ All } I \\
 \frac{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \end{array} \right\} \vdash \forall y. z \leq y}{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \\ \exists y. z \leq y \end{array} \right\} \vdash \exists y. z \leq y} \text{ Ex } I \\
 \frac{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \end{array} \right\} \vdash \exists y. z \leq y}{\left\{ \begin{array}{l} \forall x. \exists y. y \leq x; \\ \exists y. y \leq x \\ \exists y. z \leq y \end{array} \right\} \vdash \exists x. \forall y. x \leq y} \text{ Imp } I \\
 \frac{}{\{ \} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)}
 \end{array}$$

Floyd-Hoare Logic

- Also called **Axiomatic Semantics**
- Based on formal logic (first order predicate calculus)
- Logical system built from **axioms** and **inference rules**
- Mainly suited to simple imperative programming languages
- Ideas applicable quite broadly

Floyd-Hoare Logic

- Used to formally prove a property (**post-condition**) of the **state** (the values of the program variables) after the execution of program, assuming another property (**pre-condition**) of the state holds before execution

Floyd-Hoare Logic

- Goal: Derive statements of form

$$\{P\} C \{Q\}$$

- P, Q logical statements about state, P precondition, C program

- Example:

$$\{x = 1\} x := x + 1 \{x = 2\}$$

Floyd-Hoare Logic

- Approach:** For each type of language statement, give an axiom or inference rule stating how to derive assertions of form

$$\{P\} C \{Q\}$$

where C is a statement of that type

- Compose axioms and inference rules to build proofs for complex programs

Partial vs Total Correctness

- An expression $\{P\} C \{Q\}$ is a **partial correctness** statement
- For **total correctness** must also prove that C terminates (i.e. doesn't run forever)
 - Written: $[P] C [Q]$
- Will only consider partial correctness here

Simple Imperative Language

- We will give rules for simple imperative language

$$\begin{aligned}\langle \text{command} \rangle ::= & \langle \text{variable} \rangle := \langle \text{term} \rangle \\ | & \langle \text{command} \rangle ; \dots ; \langle \text{command} \rangle \\ | & \text{if } \langle \text{statement} \rangle \text{ then } \langle \text{command} \rangle \text{ else } \langle \text{command} \rangle \\ | & \text{while } \langle \text{statement} \rangle \text{ do } \langle \text{command} \rangle\end{aligned}$$

- Could add more features, like for-loops

Substitution

- Notation: $P[e/v]$ (sometimes $P[v \rightarrow e]$)
- Meaning: Replace every v in P by e
- Example:

$$(x + 2)[y - 1/x] = ((y - 1) + 2)$$

The Assignment Rule

$$\overline{\{P[e/x]\}} x := e \{P\}$$

Example:

$$\overline{\{ ? \}} x := y \{ x = 2 \}$$

The Assignment Rule

$$\overline{\{P[e/x]\}} x := e \{P\}$$

Example:

$$\overline{\{ \square = 2 \}} x := y \{ \boxed{x} = 2 \}$$

The Assignment Rule

$$\overline{\{P[e/x]\}} x := e \{P\}$$

Example:

$$\overline{\{ \boxed{x} = 2 \}} x := y \{ \boxed{x} = 2 \}$$

The Assignment Rule

$$\overline{\{P[e/x]\}} x := e \{P\}$$

Examples:

$$\overline{\{y = 2\}} x := y \{x = 2\}$$

$$\overline{\{y = 2\}} x := 2 \{y = x\}$$

$$\overline{\{x + 1 = n + 1\}} x := x + 1 \{x = n + 1\}$$

$$\overline{\{2 = 2\}} x := 2 \{x = 2\}$$

The Assignment Rule – Your Turn

- What is the weakest precondition of

$$x := x + y \{x + y = wx\}?$$

$$\left\{ \begin{array}{c} ? \\ x := x + y \\ \{x + y = wx\} \end{array} \right.$$

The Assignment Rule – Your Turn

- What is the weakest precondition of

$$x := x + y \{x + y = wx\}?$$

$$\left\{ \begin{array}{c} (x + y) + y = w(x + y) \\ x := x + y \\ \{x + y = wx\} \end{array} \right.$$

Precondition Strengthening

$$\frac{(P \Rightarrow P') \{P'\} C \{Q\}}{\{P\} C \{Q\}}$$

- Meaning: If we can show that P implies P' (i.e. $(P \Rightarrow P')$) and we can show that $\{P\} C \{Q\}$, then we know that $\{P\} C \{Q\}$

• P is **stronger** than P' means $P \Rightarrow P'$

Precondition Strengthening

- Examples:

$$\frac{x = 3 \Rightarrow x < 7 \quad \{x < 7\} x := x + 3 \{x < 10\}}{\{x = 3\} x := x + 3 \{x < 10\}}$$

$$\frac{\text{True} \Rightarrow (2 = 2) \quad \{2 = 2\} x := 2 \{x = 2\}}{\{\text{True}\} x := 2 \{x = 2\}}$$

$$\frac{x = n \Rightarrow x + 1 = n + 1 \quad \{x + 1 = n + 1\} x := x + 1 \{x = n + 1\}}{\{x = n\} x := x + 1 \{x = n + 1\}}$$

Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} x := x * x \{x < 25\}}{\{x = 3\} x := x * x \{x < 25\}}$$

$$\frac{\{x = 3\} x := x * x \{x < 25\}}{\{x > 0 \wedge x < 5\} x := x * x \{x < 25\}}$$

$$\frac{\{x * x < 25\} x := x * x \{x < 25\}}{\{x > 0 \wedge x < 5\} x := x * x \{x < 25\}}$$

Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} x := x * x \{x < 25\}}{\{x = 3\} x := x * x \{x < 25\}} \text{ YES}$$

$$\frac{\{x = 3\} x := x * x \{x < 25\}}{\{x > 0 \wedge x < 5\} x := x * x \{x < 25\}}$$

$$\frac{\{x * x < 25\} x := x * x \{x < 25\}}{\{x > 0 \wedge x < 5\} x := x * x \{x < 25\}}$$

Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}{\{x = 3\} \ x := x * x \ \{x < 25\}} \text{ YES}$$

$$\frac{\{x = 3\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}} \text{ NO}$$

$$\frac{\{x * x < 25\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}$$

Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}{\{x = 3\} \ x := x * x \ \{x < 25\}} \text{ YES}$$

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