CS 473: Algorithms, Fall 2010
HW 5 (due Tuesday, October 12)

This homework contains four problems. **Read the instructions for submitting homework on the course webpage**. In particular, *make sure* that you write the solutions for the problems on separate sheets of paper; the sheets for each problem should be stapled together. Write your name and netid on each sheet.

Collaboration Policy: For this home work, Problems 1-2 can be worked in groups of up to 3 students each.

Problem 0 should be answered in Compass as part of the assessment HW5-Online and should be done individually.

0. (10 pts) HW5-Online.

1. • (30 pts) In a word processor the goal of "pretty-printing" is to take text with a ragged right margin, like this,

```
Call me Ishmael.
Some year ago,
never mind how long precisely,
having little or no money in my purse,
and nothing particular to interest me on shore,
I though I would sail about a little
and see the watery part of the world.
```

and turn it into text whose right margin is as "even" as possible, like this

```
Call me Ishmael. Some year ago, never
mind how long precisely, having little
or no money in my purse, and nothing
particular to interest me on shore, I
though I would sail about a little
and see the watery part of the world.
```

To make this precise enough for us to start thinking about how to write a pretty-printer for text, we need to figure out what it means for the right margins to be "even". So suppose our text consists of a sequence of *words*, $W = \{w_1, w_2, \ldots, w_n\}$, where w_i consists of c_i characters. We have a maximum line length of L. We will assume we have a fixed-width font and ignore issues of punctuation or hyphenation.

A *formatting* of W consists of a partition of the words in W into *lines*. In the words assigned to a single line, there should be a space after each word except the last; and so if $w_j, w_{j+1}, \ldots, w_k$ are assigned to one line, then we should have

$$\Big[\sum_{i=j}^{k-1}(c_i + 1)\Big] + c_k \leq L.$$

We will call an assignment of words to a line *valid* if it satisfies this inequality. The difference between the left-hand side and the right-hand side will be called the *slack* of the line-that is, the number of spaces left at the right margin.

Given a partition of a set of words W, the *penalty* of the formatting is the sum of the *squares* of the slacks of all lines (including the last line). Give an efficient algorithm to find a partition of a set of words W into valid lines, so that the penalty of the formatting becomes minimized.

- (20 pts) Implement your iterative algorithm in C or C++ or Java. The input to your program will be L n c_1 c_2 $\ldots$ c_n on a single line: L is the number of characters per line allowed, n is the number of words in the text and each c_i is the length of a word w_i in our text, $W = \{w_1, \ldots, w_n\}$, in order. The output should be the penalty of an optimal formatting as well as the lengths of all lines in an optimal formatting of the words. You need to submit a print out of the code along with the output of your code on a set of inputs that we will provide on the website.

2. (40 pts) Let $G = (V, E)$ be an undirected graph. A subset $S \subseteq V$ of nodes in G is called a *dominating set* if for all $v \in V$, we have $v \in S$ or there is some node $u \in S$ such that $(u, v) \in E$. In other words every node in $V \setminus S$ is connected by an edge to some node in S. Given non-negative weights $w(v)$ on the nodes of V the goal is to find a minimum-weight dominating set in G. This problem is known to be NP-Hard in general graphs. Describe a polynomial time algorithm for this problem when G is a tree.